## **Sophisticated Quantum Search Without Entanglement**

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Grover's quantum search algorithm has recently been implemented without entanglement, by replacing multiple particles with a single particle having exponentially many states. We recall that *all* physical resources must be accounted for to quantify algorithm complexity, and that this scheme typically incurs exponential costs in some other resource(s). In particular, we demonstrate that a recent experimental realization requires exponentially increasing precision. There is, however, a quantum algorithm which searches a "sophisticated" database (not unlike a Web search engine) with a single query, but which we show does not require entanglement even for multiparticle implementations.

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Quantum algorithms must exploit some physical resource unavailable to classical computers in order to solve problems in fewer steps [1–5]. Entanglement, which seems the "spookiest" [6] to many people, has been argued to be the crucial quantum mechanical resource [7]. This belief inspires, for example, the criticism that NMR experiments performed to date [8] have not actually realized quantum algorithms because at each time step the state of the system can be described as a probabilistic ensemble of unentangled quantum states [9]. Lloyd [10] and Ahn *et al.* [11] have recently suggested, however, that entanglement is *not* necessary for Grover's quantum search algorithm [4]. In this Letter we clarify the situation by demonstrating that, contrary to their claims, the experimental realization of Ahn *et al.* [11] requires an exponentially increasing amount of a resource—precision—replacing entanglement. But we do not conclude from this that entanglement (or some replacement resource) is required. Rather, we make the new and surprising observation that efficient quantum search of a "sophisticated" database (not unlike a Web search engine) requires no entanglement at any time step: a quantum-over-classical reduction in the number of queries is achieved using only interference, not entanglement, within the usual model of quantum computation.

The problem which forms the context for our discussion is database search–identifying a specific record in a large database. Formally, we label the records  $\{0, 1, \ldots, N-1\}$ , where, for convenience when we write the numbers in binary, we take  $N = 2^n$  for *n* a positive integer. Grover considered databases which, when queried about a specific number, respond only that the guess is correct or not [4]. On a classical reversible computer we can implement a query by a pair of registers  $(x, b)$ , where x is an *n*-bit string representing the guess, and *b* is a single bit which the database will use to respond to the query. If the guess is correct, the database responds by adding 1 (mod2) to *b*; if it is incorrect, it adds 0 to *b*. That is, the response of the database is the operation:  $(x, b) \rightarrow (x, b \oplus f_a(x))$ , where  $\Theta$  denotes addition mod2 and  $f_a(x) = 1$  when  $x = a$ , and 0 otherwise. Thus if *b* changes, we know that the guess is correct. Classically, it takes  $N - 1$  queries to solve this problem with probability 1.

Quantum algorithms work by supposing that they will be realized in a quantum system, such as those described by Lloyd [10] and Ahn *et al.* [11], which can be in a superposition of "classical" states. These states form a basis for the Hilbert space whose elements represent states of the quantum system. The simplest such system is a *qubit* [12], which can be in a superposition of the states of a classical bit, i.e., 0 and 1. More generally, Grover's algorithm works with quantum queries which are linear combinations  $\sum c_{x,b} |x, b\rangle$ , where  $c_{x,b}$  are complex numbers satisfying  $\sum |c_{x,b}|^2 = 1$  and  $|x, b$ is Dirac notation [13] for the quantum state which represents the classical state  $(x, b)$  of the two registers. The operations in quantum algorithms are unitary transformations, the quantum mechanical generalization of reversible classical operations. Thus the operation of the database that Grover considered is implemented on superpositions of queries by a unitary transformation ( *f<sub>a</sub>*-controlled-NOT), which takes  $|x, b\rangle$  to  $|x, b \oplus f_a(x)\rangle$ . Figure 1 illustrates a quantum circuit implementation [14] of Grover's algorithm [4]. By using  $\lfloor \frac{\pi}{4} \rfloor$  $\frac{\pi}{4} \sqrt{N}$ quantum queries, it identifies the answer with probability close to 1: The final vectors for the *N* possible answers *a* are nearly orthogonal.

This quantum circuit acts on an initial state  $\psi_0 = |0\rangle \cdots$  $|0\rangle|0\rangle = |0...0,0\rangle$ . The first set of gates transforms the state to  $\psi_1 = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \otimes \cdots \otimes \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = \sum_{x=0}^{N-1} |x\rangle (|0\rangle - |1\rangle) / \sqrt{2N}$ . Both of these states,  $\psi_0$  and  $\psi_1$ , are tensor products of the states of the individual qubits, so they are *unentangled* [16]. This is no longer true for subsequent states of the system (except when  $N = 2$ ). The last qubit, however, is never entangled with the others—after the first time step it remains in state  $\frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$ . Lloyd's observation [10], which is exploited by Ahn *et al.* [11], is that the absence



FIG. 1. A (schematic) quantum circuit implementing Grover's algorithm. Each horizontal line represents a single qubit, which is initialized (on the left) in state  $\ket{0}$ . The portion of the circuit enclosed in the grey square is repeated  $\left[\frac{\pi}{4} \sqrt{N}\right]$  times and then the top  $n$  qubits are measured.  $H$  is the Hadamard transformation  $\begin{pmatrix} 1 \\ 1 \end{pmatrix}$ qubits are measured. *H* is the Hadamard transforma-<br> $\binom{1}{1}/\sqrt{2}$ , *X* is the Pauli matrix  $\binom{0}{1}$ , and the "gates" acting on all  $n + 1$  qubits are  $f_a$ - and  $f_0$ -controlled-NOT transformations, respectively. As was first noted by Brassard and Høyer, and subsequently by Grover, the  $H^{n\otimes}(\otimes I_2)$  transformation conjugating the  $f_0$ -controlled-NOT gate can be replaced by almost any unitary transformation [15]. Our discussion is independent of this choice.

of entanglement in the  $N = 2$  case of Grover's algorithm (for which the guess register consists of a single qubit) generalizes to arbitrary *N* if the guess register is realized by one *N* state particle rather than by *n* qubits. In fact, Jozsa and Ekert [7] made exactly this observation several years ago: They wrote, "The state of  $n$  qubits is a  $2^n$ dimensional space and can be isomorphically viewed as the state space of a *single* particle with 2*<sup>n</sup>* levels. Thus we simply view certain states of a single  $2^n$  level particle as 'entangled' via their correspondence under a chosen isomorphism between  $\mathscr{A}_1\mathscr{H}_2$  and  $\mathscr{H}_{2^n}$  (where  $\mathcal{H}_k$  denotes a Hilbert space of dimension *k*.)." So, despite the implication of the Lloyd [10] and Ahn *et al.* [11] papers, there is nothing special about Grover's algorithm: Reformulating *any* quantum algorithm this way, i.e., disregarding the tensor product structure of Hilbert space implicit in the use of qubits, removes entanglement from the system by *definition.* Nevertheless, one might hope that if a quantum algorithm—such as Grover's—can be implemented naturally with a single particle, as Lloyd suggests [10] and as Ahn *et al.* realized experimentally with *N* Rydberg levels of a cesium atom [11], there is some physical advantage to be gained.

But Jozsa and Ekert [7] continue, "However the physical implementation of this correspondence appears always to involve an exponential overhead in some physical resource so that the isomorphism is *not* a valid correspondence for considerations of complexity.", again anticipating Lloyd's discussion [10]. Although their data indicate that increasing *N* requires more repetitions of the experiment to extract the answer [11], Ahn *et al.* neglect the exponential overhead required for measurement and for realization of  $N \times N$  unitary transformations: They claim that extrapolation from their  $N = 8$  experiments to  $N = 20$  is straightforward and suggest that ultrafast shaped terahertz pulses [17] might realize more general unitary transformations than those used in their implementation of Grover's algorithm. But because the difference (detuning) between adjacent Rydberg energy levels converges to 0 polynomially in 1*N* (for *N* labeling the energy levels) [13,17], both the laser pulses and the final measurements must be specified with exponentially increasing precision in *n*, the size of the problem [18]. This should be contrasted with the standard model for quantum computation using *polylocal* transformations implemented by polynomially many bounded size gates on Hilbert spaces with a tensor product decomposition [19,20]; these require specification of only polynomially many nontrivial amplitudes with constant precision. As Bernstein and Vazirani [2] and Shor [21] already emphasized in their original analyses of quantum models for computing, *all* physical resources must be accounted for to quantify algorithm complexity; it is a mistake to ignore some because the requirements for them do not overwhelm small *N* experiments.

Having identified an exponential cost associated with the Ahn *et al.* realization [11] of Lloyd's suggestion for entanglement removal [10], we are now ready to demonstrate that it is also a mistake to infer, as Lloyd's presentation might lead one to [22], that quantum algorithms *require* entanglement—or an exponential amount of some resource replacing it—to improve on classical algorithms. Rather than Grover's "naive" database, let us consider a sophisticated database which, when queried about a specific number, responds with information about how close the guess is to the answer. This kind of response is more like that returned by, for example, Web search engines, which typically order pages by relevance [23]. A simple measure of relevance comes from the vector space model of information retrieval [24]: The records in the database and the guess are represented by vectors; then (the cosine of) the angle between a guess and any record measures their similarity and can be computed from the dot product of their vectors. In our setting the sophisticated database acts on a query  $(x, b)$  by computing the dot product of the *n* dimensional binary vectors  $x \cdot a$  and adding it to *b* (mod2). Thus  $(x, b) \rightarrow (x, b \oplus g_a(x))$ , where  $g_a(x) = x \cdot a$ . Classically, *n* queries suffice to identify *a* with probability 1.

Quantum mechanically, an underappreciated algorithm of Bernstein and Vazirani [2], rediscovered by Terhal and Smolin [5], searches this sophisticated database with only a *single* quantum query [25]. The operation of the database is implemented by the unitary transformation ( $g_a$ -controlled-NOT) which takes  $|x, b\rangle$  to  $|x, b \otimes g_a(x)\rangle$ . A quantum circuit for their algorithm (slightly improved [14]) is shown in Fig. 2. The first set of gates is the same as in Fig. 1, and takes  $\psi_0 =$ <br> $|0...0,0\rangle$  to  $\psi_1 = \sum_{x=0}^{N-1} |x\rangle (|0\rangle - |1\rangle)/\sqrt{2N}$ . After the database responds to this quantum query, the state is  $\psi_2 = \sum_{x=0}^{N-1} (-1)^{x+a} |x\rangle (|0\rangle - |1\rangle)/\sqrt{2N}$ . The final set of gates database responds to this quantum query, the state is  $\psi_2$  = outputs  $\psi_3 = |a\rangle \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$ , whereupon measuring



FIG. 2. A (schematic) quantum circuit implementing Bernstein and Vazirani's algorithm. Each horizontal line again represents a qubit, which is initialized (on the left) in state j0. *H* and *X* are the same as in Fig. 1 and the gate acting on all  $n + 1$  qubits is the *ga*-controlled-NOT transformation of the sophisticated database. The top *n* qubits are measured at the end of the circuit.

the first (*n*-qubit) register identifies *a* with probability 1 (the output states for different *a*'s are orthogonal). Comparing with Grover's algorithm, we recognize that the last qubit still remains unentangled with the first register, so that we could again implement the latter with a single 2*<sup>n</sup>* state particle and have no entanglement at any time step. But this would be redundant: *There is no entanglement in Bernstein and Vazirani's algorithm.* To see this, one observes that, just as in Grover's algorithm, there is no entanglement in  $\psi_0$  or  $\psi_1$ , and there is none in  $\psi_3$ , since  $|a\rangle$  is simply a tensor product of qubits each in state  $|0\rangle$ or  $|1\rangle$ . However,  $\psi_3$  was obtained from  $\psi_2$  by a unitary transformation acting on each of the  $n + 1$  qubits separately. Such a unitary transformation cannot change the entanglement of a state, so  $\psi_2$  must also be unentangled.

To summarize, any quantum algorithm in the usual polylocal model for quantum computing can be rewritten to have no entanglement at any time step, simply by disregarding the tensor product structure of the Hilbert space. Doing so physically incurs some exponential cost: in energy, in measurement precision, or in the specification of the required unitary transformations. One should not, however, conclude that entanglement is *required* for quantum-over-classical complexity reduction. Without entanglement at any time step, Bernstein and Vazirani's quantum algorithm for sophisticated database search does not just reduce the number of queries required classically by a square root factor but does so all the way from *n* to 1. Futhermore, we have shown for the first time that quantum interference alone suffices to reduce the query complexity of a problem within the standard model for quantum computation. Since implementing the *ga*-controlled-NOT "gate" with a subcircuit of *local* gates would introduce entanglement at intermediate time steps, however, one might conclude that counting queries (or, more generally, nonlocal function calls) is a poor way to study the power of quantum algorithms [26]. But it was Simon's algorithm [3] (which exponentially reduces the number of nonlocal evaluations required to determine the period of a function) that led to Shor's quantum factoring algorithm [21], so

it seems more productive to understand the quantum search of a sophisticated database as demonstrating the importance of interference and orthogonality, rather than entanglement, in quantum algorithms. This perspective may contribute to discovering the new algorithms necessary for quantum computing to become more generally useful.

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