

Knottedness in Ring Polymers

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We represent ring polymers in free space with the rod-bead model and show through unbiased computer simulations that the probability of observing a trivial self-entanglement (P) has a decreasing exponential dependence on the contour length (N) of the polymer, or that $P = \exp(-N/N_0)$. The characteristic length (N_0) varies by many orders of magnitude depending on chain flexibility and solvent quality. We also suggest that sufficiently large knots are always composite, not prime.

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Both chemists and biochemists have synthesized polymeric rings. Once a ring is formed, its topological state is uniquely defined and invariant, at least until a bond is broken. Any nondestructive manipulation—like stretching, bending, or folding—will not impact a ring's topological state. While natural (circular) DNA is unknotted, synthetic rings can exhibit very complex knotted configurations.^{1,2}

A great deal of theoretical effort has been expended in the study of entangled systems, since both the dynamic and the thermodynamic properties of such systems are functions of their topological state.³⁻¹⁰ The simplest system that exhibits entanglement is a single-ring polymer. More complicated polymeric systems that are of interest are a ring linked to a wire (this problem bears an interesting relation to the Aharonov-Bohm effect¹⁰), rings linked to each other, and networks.

A most fundamental question about self-entanglement has been posed by Delbrück:¹ How does the probability of observing self-entanglement ($1 - P$) in a ring depend on the contour length (N)? Delbrück's question has so far proven resistant to all analytical attacks, so several numerical efforts have been attempted by various researchers.¹¹⁻¹⁶

Much of the early numerical work^{11,12} was done on a lattice (instead of in free space); biased methods were used to generate ring polymers, and systems were small. Later work¹⁴ fixed many of the aforementioned problems, where the polymers were modeled off lattice and properly formed, but computer technology at the time was not sufficient for researchers to study systems with very large sizes and good statistics. The most well received work to date seems to be that of Michels and Wiegel,^{10,15} which suggests that a power law relates P to N for Gaussian chains. (A Gaussian chain has $r=0$ for the model that we describe shortly.)

In this work, we study unbiased rings in free space, which seem to be the largest analyzed to date (by almost an order of magnitude), and for the first time attempt a methodical study of the effects of the polymers's exclud-

ed volume and flexibility on knotting. We believe that a decreasing exponential relates P to N , or that $P = \exp(-N/N_0)$, where the characteristic length (N_0) can vary by several orders of magnitude depending on the flexibility of the rings.

We represent ring polymers with the free-space rod-bead model.¹⁷ Two parameters are required to create a model instance: the contour length (N) and the bead radius (r). A ring is composed of a closed backbone of N unit bonds. At the N vertices where two bond ends abut, we imagine beads, each of radius r . If any two beads intersect, the excluded-volume condition is said to be violated, and the entire instance must be rejected. A small r represents a flexible polymer in an ideal (Θ) solvent, while a large r represents a swollen (and less flexible) polymer in a good solvent.

We produce ring-polymer model instances with the dimerization method of Chen.¹⁴ This method is preferable to alternatives¹⁷ for three reasons: It is unbiased, it is very efficient in terms of computer time, and it is well suited to parallel architectures.

The rings produced by dimerization each have a unique topological state, and can be characterized by a knot invariant known as the Alexander polynomial $\Delta(s)$.^{18,19} Other more elaborate invariants exist, like the Jones polynomial,²⁰ but their complexity makes them difficult to efficiently calculate, so we follow the lead of other researchers and use $\Delta(s)$.

The Alexander polynomial is a true invariant, in that a given knot will always be represented by the same $\Delta(s)$, no matter how much we bend, fold, or stretch it. However, there are cases where aliasing exists; that is, two knots in different topological states can lamentably share the same $\Delta(s)$. For example, Fig. 1 demonstrates that $\Delta[5_1] = \Delta[10_{132}]$ and that $\Delta[8_{10}] = \Delta[(3_1)(3_1)(3_1)]$. It is particularly unfortunate that some common combinations of simple prime knots [like a granny knot, $(3_1)(3_1)$] can be mistaken for more exotic prime knots (like 8_{20}).

Given a knot, there is a simple algorithm to determine

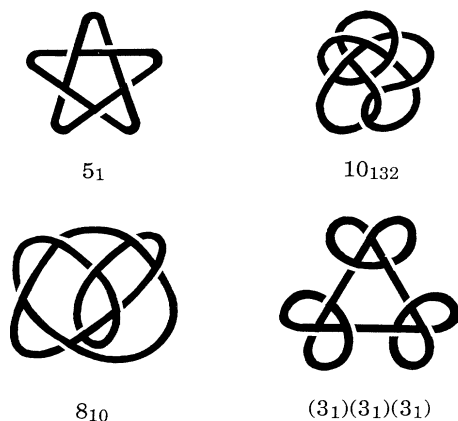


FIG. 1. Knots 5_1 and 10_{132} share the same Alexander polynomial, $\Delta(s) = s^4 - s^3 + s^2 - s + 1$, as do knots 8_{10} and $(3_1)(3_1)(3_1)$, which have $\Delta(s) = s^6 - 3s^5 + 6s^4 - 7s^3 + 6s^2 - 3s + 1$.

its Alexander polynomial.^{10,18,19} However, one step requires taking the determinant of an M by M matrix of polynomials, where $M+1$ is the number of bond-bond crossings observed when the knot is properly projected onto a plane. Even using Gaussian elimination to take the determinant requires time M^3 , which can be prohibitive in the study of large systems; it is easy to get many hundreds or even thousands of crossings.

It is advantageous to preprocess a ring in such a way as to reduce the number of crossing ($M+1$) that it displays when projected onto a plane. We have developed a very simple smoothing operation that properly maintains any ring's topological state while significantly reducing $M+1$; if the triangle defined by any three sequential beads is not crossed by any bond, the middle bead is removed and the two survivors are directly connected. This operation does distort the ring, but the underlying topological state is properly maintained. This operation can be done very efficiently, in time $N \ln(N)$, frequently to the point where the resulting ring has three beads left, with zero crossings. We feel that many entanglement problems are impractical unless one strives to remove detail that is irrelevant to the underlying topological features.

After smoothing a large ring, it is not efficient to immediately start the expensive Gaussian elimination. It often becomes apparent that one has a composite knot that is composed of several factors.²¹ If a knot is com-

$$\Delta(\boxed{A} \boxed{B}) = \Delta(\boxed{A}) \cdot \Delta(\boxed{B}).$$

FIG. 2. If a knot can be "factored" into two parts as indicated, the $\Delta(s)$ of the entire composite knot is the product of the Alexander polynomials of the constituents.

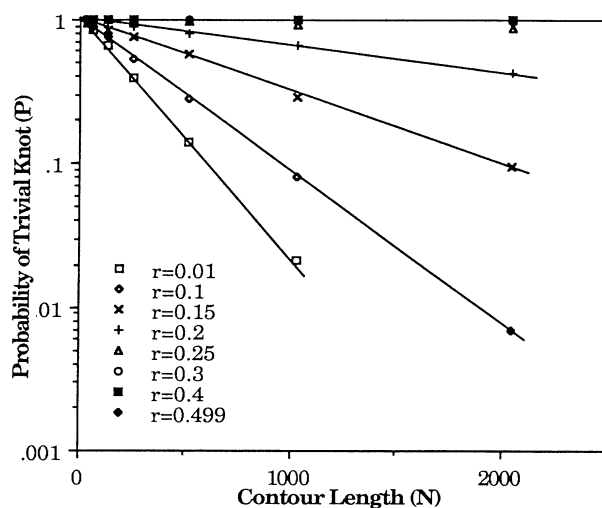


FIG. 3. A semilogarithmic plot of the probability of observing a trivial knot (P) as a function of the number of beads in a ring (N), for several values of bead radius (r).

posed of parts A and B , as in Fig. 2, $\Delta[AB] = \Delta[A]\Delta[B]$.

Given that the identification time goes as M^3 , "halving" the problem results in a fourfold speedup. [A second benefit is that it is sometimes possible to identify a $(3_1)(3_1)$ without concern that it might be an 8_{20} , for example.] We calculate determinants by using Gaussian elimination to solve a system of polynomials, utilizing symbolic methods.²² It is possible to avoid symbolic computation altogether [by precompiling tables of $\Delta(s = -1)$ and $\Delta(s = -2)$ for known Alexander polynomials], but one gives up the ability of identifying new large knots.^{11,21} We have indeed found knots with more than ten crossings,²¹ the size of the knot table that we have used,¹⁸ but this happens rarely.

Figure 3 is a semilogarithmic plot of the probability of observing a trivial knot (open loop), P , as a function of the number of beads, N , for various values of bead radius r . Each point represents 5000 rings. The error in P is 2%, given a 95% confidence interval. We see a linear fit, and therefore propose the empirical relation $P(N)$

TABLE I. The characteristic length (N_0) and radius-of-gyration exponent (ν_{eff}) as functions of bead radius (r).

Radius	N_0	ν_{eff}
0.01	2.6×10^2	0.51
0.1	4.0×10^2	0.54
0.15	8.7×10^2	0.57
0.2	2.4×10^3	0.59
0.25	1.5×10^4	0.60
0.3	1.1×10^5	0.60
0.4	3×10^5	0.61
0.499	8×10^5	0.61

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