Can the Channel Capacity of a Light-Wave Communication System Be Increased by the Use of Photon-Number-Squeezed Light?

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The channel capacity C of a light-wave communication system based on photoevent point-process observations *cannot* be increased by use of photon-number-squeezed light. Under a constraint of maximum photon rate λ_{max} , the channel capacity $C = \lambda_{max}/e$ is achieved with a Poisson process. On the other hand, the channel capacity of a communication system based on photon counting *can* be increased by use of photon-number-squeezed light. The improvement vanishes in the limit of very small mean counts.

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All light-wave communication systems that have been developed to date make use of Poisson (or super-Poisson) light.¹ For Poisson light, the variance of the photon number is identically equal to its mean for all values of the counting time T. Photon-number-squeezed light, on the other hand, has a photon-number variance that is less than its mean for all or some values of T.^{2,3} Such light is intrinsically nonclassical in nature. The earliest source of unconditionally photon-number-squeezed (also called sub-Poisson) light exhibited only a slight reduction of the variance.⁴ Stronger photon-number squeezing has been produced in more recent experiments^{5,6} and continuing developments⁷⁻⁹ promise further improvement in the future. It is therefore of interest to investigate whether there might be an advantage to using photon-numbersqueezed light in a direct-detection light-wave communication system.

There are two classes of mechanisms by means of which unconditionally photon-number-squeezed light may be generated. In the first class, squeezed photons are produced from a beam of *initially Poisson* (or super-Poisson) *photons*. This can be achieved in a number of ways, e.g., by the use of correlated photon beams or quantum nondemolition (QND) measurements.¹⁰ An experiment of this kind was recently carried out by Rarity, Tapster, and Jakemen.¹¹ Sub-Poisson photons were generated from the pair of correlated photon beams produced in parametric down-conversion; one of the twin beams was used to gate the other beam selectively via dead-time control.

The second class of mechanisms relies on the direct generation of squeezed photons from a beam of *initially sub-Poisson excitations* (e.g., electrons).^{3,12} This technique was first used by Teich and Saleh in a space-charge-limited version of the Franck-Hertz experiment.⁴

Perhaps the simplest implementation of this principle is achieved by the driving of a light-emitting diode with a sub-Poisson electron current.^{6,8}

In this Letter we show that the channel capacity of a light-wave communication system based on the observation of the *photoevent point process cannot* in principle be increased by the use of photon-number-squeezed light. On the other hand, the channel capacity of a *photon-counting* system *can* be increased by the use of photon-number-squeezed light. The channel capacity is the maximum rate of information that can be transmitted through a channel without error.¹³ We also provide an example in which the use of photon-number-squeezed light produced from Poisson light either degrades or enhances the *error performance* of a simple binary on-off keying photon-counting system, depending on where the average power constraint is placed.

Consider the transformation of a Poisson beam of photons (represented by a Poisson point process¹⁰ N_t of rate μ_t) into a sub-Poisson beam of photons represented by a point process M_t of rate λ_t . The events of the initial process N_t are assumed to be observable (e.g., by the use of correlated photon beams or a QND measurement) and their registrations used to operate a mechanism which, in accordance with a specified rule, leads to the events of the transformed photon process M_t . The rate λ_t of the process M_t is thereby rendered a function of the realizations of the initial point process N_t at prior times, i.e., $\lambda_t = \lambda_t (N_t; t' \le t)$.

Several examples of transformations of this kind that have been suggested for use in quantum optics are illustrated in Fig. 1 and discussed below. It is assumed for simplicity (but without loss of generality) that the various conversions can be achieved in an ideal manner.

(i) Dead-time deletion: Delete all photons within a



FIG. 1. Several transformations of Poisson photons into sub-Poisson photons that have been suggested for use in quantum optics.

prescribed fixed (nonparalyzable) dead time τ following the registration of a photon.¹⁴ Rarity, Tapster, and Jakeman¹¹ generated photon-number-squeezed light by using one of the twin beams produced in parametric down-conversion to gate photons selectively from the other beam via dead-time control. Dead-time deletion could also be used with correlated photon beams produced in other ways.

(ii) Coincidence decimation: Remove all pairs of photons separated by a time shorter than a prescribed time interval τ' . This is achieved, for example, in secondharmonic generation; two photons closer than the intermediate-state lifetime of the second-harmonicgeneration process are exchanged for a third photon (which is at twice the frequency and therefore easily eliminated).¹⁵

(iii) Decimation: Select every rth photon (r=2, 3, ...) of an initially Poisson photon process, deleting all intermediate photons. Saleh and Teich¹⁶ suggested using correlated photon beams to implement this technique. In cascaded atomic emissions from ⁴⁰Ca, for example, sequences of correlated photon pairs (green and violet) are emitted. The green photons can be detected and used to operate a gate that passes every rth violet photon. Decimation control could also be used in conjunction with parametric-down-conversion photon twins.

(iv) Overflow count deletion: The number of photons occurring in preselected time intervals $[0, T_0]$, $[T_0, 2T_0], \ldots$, is counted, retaining the first n_0 photons in each time interval (without changing their occurrence times) and deleting the remainder. If the average number of photons in $[0, T_0]$ of the initial process is $\gg n_0$, then the transformed process will almost always contain n_0 photons within this time interval. As an example, Mandel¹⁷ suggested that if a collection of n_0 atoms in the ground state are subjected to a brief, intense, incoherent excitation pulse, all n_0 atoms will become excited with high probability; the radiated optical field would then be describable, to good approximation, by an n_0 -photon state. Related schemes have been proposed by Yuen¹⁸ and by Stoler and Yurke¹⁹ for use with parametric processes.

We proceed to illustrate that none of these modifications can increase the channel capacity of a communication system based on photoevent point-process observations.

If a constraint is placed on the rate of the initial Poisson process $\mu_t \leq \mu_{max}$, then it is obvious that C cannot be increased by the modification $N_t \rightarrow M_t$. This is simply a consequence of the definition of channel capacity: It is the rate of information carried by the system without error, maximized over all coding, modulation, and *modification* schemes. Can the modification $N_t \rightarrow M_t$ increase the channel capacity if the constraint is instead placed on the rate of the modified process λ_t (i.e., $\lambda_t \leq \lambda_{max}$)?

We address this question for an arbitrary self-exciting point process M_t of rate $\lambda_t(M_t; t' \leq t)$. This is a process that contains an inherent feedback mechanism in which present event occurrences are affected by the previous event occurrences of the same point process. Of course, the modified Poisson processes $N_t \rightarrow M_t$ introduced above are special cases of self-exciting point processes.

An example of a system that generates a self-exciting point process is that of rate compensation (by linear feedback) of a source which, without feedback, would produce a Poisson process. Let each photon registration at time t_i cause the rate of the process to be modulated by a factor $h(t-t_i)$ (which vanishes for $t < t_i$). In linear negative feedback the rate is $\lambda_t = \lambda_0 - \sum_i h(t - t_i)$, where λ_0 is a constant. If the instantaneous photon registration rate happens to be above the average then it is reduced, and vice versa. This process is schematically illustrated in Fig. 1(f) for two adjacent subintervals T_1 and T_2 . Yamamoto, Imoto, and Machida²⁰ suggested the use of rate compensation in conjunction with a QND measurement (using the optical Kerr effect), but it could be used just as well, for example, with correlated photon pairs. Dead-time deletion can be viewed as a special case of rate compensation in which the occurrence of an event zeros the rate of the process for a specified time period τ after the registration.¹⁰

Now consider a communication system that uses a point process $M_t(X)$ whose rate $\lambda_t(X)$ is modulated by a signal X_t . The process $M_t(X)$ can be an arbitrary self-exciting point process (e.g., it can be photon-number squeezed) which includes processes obtained by the feed-forward or feedback modification of an otherwise Poisson process. Neither feedforward nor feedback transformations can increase the capacity of this channel, as provided by Kabanov's theorem.²¹

Kabanov's theorem.— The capacity of the pointprocess channel cannot be increased by feedback. Under the constraint $\lambda_t \leq \lambda_{max}$, the channel capacity C is

$$C = \lambda_{\max}/e. \tag{1}$$

When the capacity is achieved, the output of the pointprocess channel is a Poisson process with rate $\lambda_t = \lambda_{max}/e$ (the base *e* has been used for simplicity). The channel capacity has also been determined under added constraints on the minimum rate (dark events) and on the mean rate.²¹ A coding theorem has also been proved.

In summary, *no* increase in the channel capacity of a point-process light-wave communication system may be achieved by the use of photons that are first generated with Poisson statistics and subsequently converted into sub-Poisson statistics regardless of whether the power constraint is placed at the Poisson photon source or at the output of the conversion process. Nor may an increase in channel capacity be achieved by the use of feedback to generate a self-exciting point process.

These conclusions are valid only when there are no restrictions on the receiver structure. The conclusion is different if the receiver is operated in the *photon*counting regime, in which information is carried by a random variable n representing the number of photoevents registered in time intervals of prescribed duration T (rather than by the photon occurrence times).

The capacity of the photon-counting channel is given by 13

$$C = B[\bar{n}\ln(1+1/\bar{n}) + \ln(1+\bar{n})], \qquad (2)$$

where \bar{n} is the mean number of counts in T and B = 1/T is the bandwidth. Two limiting expressions emerge:

$$C = \begin{cases} B\bar{n}\ln(1/\bar{n}), & \bar{n} \ll 1, \\ B\ln(\bar{n}), & \bar{n} \gg 1. \end{cases}$$
(3)

If an added constraint is applied to the photon counts, such that they must obey the Poisson counting distribution, the capacity is further reduced. In that case, the limiting results analogous to Eq. (3) are

$$C = \begin{cases} B\bar{n}\ln(1/\bar{n}), & \bar{n} \ll 1, \\ \frac{1}{2}B\ln(\bar{n}), & \bar{n} \gg 1. \end{cases}$$
(4)

In the case of photon counting, therefore, an increase in the channel capacity *can* in principle be realized by the use of photon-number-squeezed light. However, in the small mean-count limit $\bar{n} \ll 1$ (very short T), the capacity of the Poisson counting channel approaches that of the unrestricted counting channel, and the advantage of photon-number squeezing disappears. This is not unexpected in view of the result obtained from Kabanov's theorem for the point-process channel.

The channel capacity provides a limit on the maximum rate of error-free information transmission for all codes, modulation formats, and receiver structures.¹³ As such, it does not specify the performance (error probability) achievable by a communication system with prescribed coding, modulation, and receiver structure.

It is therefore of interest to examine the performance of a system with specified structure. We consider a binary on-off keying photon-counting system.¹ The information is transmitted by the selection of one of two values for the photon rate λ_t , in time slots of (binarydigit) duration T. The receiver operates by counting the number of photons received during the time interval Tand then deciding which rate was transmitted in accordance with a likelihood-ratio decision rule (threshold test). For simplicity, it is assumed that background light, dark noise, and thermal noise are absent so that photon registrations are not permitted when the keying is OFF (i.e., false alarms are not possible). Furthermore, the detector quantum efficiency is taken to be unity so that system performance is limited only by the quantum fluctuations of the light.

A measure of performance for a digital system such as this is the error probability P_e . In the simplified system described above, errors are possibly only when the keying is ON and 0 photons are received (a miss). For a Poisson transmitter, P_e is¹

$$P_e(\text{Poisson}) = \frac{1}{2} \exp(-\bar{n}), \tag{5}$$

where \bar{n} denotes the mean number of emitted photons. To minimize P_e , \bar{n} is made equal to its maximum allowed value \bar{n}_{max} . This result is now compared with those obtained for photon-number-squeezed light derived from an initially Poisson source. The outcome will depend on where the mean photon-number constraint is placed. Two transformations are explicitly considered: deadtime deletion and decimation.

(i) Dead-time deletion: For a nonparalyzable deadtime modifier that is always blocked for a dead-time period τ at the beginning of the counting interval T, the passage of 0 photons arises from the emission of 0 photons in the time $T - \tau$, independent of the number of emissions during τ . The error probability for this system is therefore

$$P_e(\text{dead time}) = \frac{1}{2} \exp[-\bar{n}(1 - \tau/T)].$$
 (6)

To minimize error under the constraint $\bar{n} \leq \bar{n}_{max}$, we take $\bar{n} = \bar{n}_{max}$. The error is obviously larger than that for the Poisson channel [Eq. (5)] and so no performance enhancement can be achieved by use of this modifier. If, instead, the dead-time modifier is always *unblocked* at the beginning of each bit then the passage of 0 photons arises from the emission of 0 photons in the time *T*, and the dead time has no effect on the error rate in this simple system. Calculations for the unblocked counter in the presence of false alarms, however, demonstrate that the presence of deal time always does, in fact, degrade system performance with such a constraint.²² Although the detailed calculations were carried out for electrical dead time, the results are also applicable for optical dead time when the photon detection efficiency $\eta = 1$. On the

other hand, if the constraint is placed on the mean photon count \overline{m} after dead-time modification ($\overline{m} \leq \overline{m}_{max}$), it can be shown that there exists a value of \overline{m}_{max} below which performance is degraded and above which performance is improved, relative to the Poisson channel.

(ii) Decimation: We assume that the decimation parameter r=2 (i.e., every other photon of a Poisson sequence of events is selected) and that the decimation process is reset at the beginning of each bit (i.e., the first photon in each bit is not selected). The error probability is then

$$P_e(\text{decimation}) = \frac{1}{2} (1 + \bar{n}) \exp(-\bar{n}), \tag{7}$$

which again represents a degradation of performance in comparison with the Poisson channel (under a constraint $\bar{n} \leq \bar{n}_{max}$). In this case, the error rate is increased because there are two ways for the passage of 0 photons to arise in the time T: from the emission of 0 photons or from the emission of 1 photon. However, if the constraint is placed on the modified process, then, once again, there exists a value of \bar{m}_{max} below which performance is degraded and above which it is improved, relative to the Poisson channel.

It is evident from these examples that system performance can be enhanced by the use of photonnumber-squeezed light, provided that the average power constraint is applied to the squeezed light. No enhancement of system performance emerges in converting Poisson photons into squeezed photons when the average power constraint is at the Poisson source.

Losses have been ignored in the considerations above. It is important to keep in mind the inevitable random photon deletion that results from absorption, scattering, and the finite quantum efficiency of the detector. It is well known that these deletions will transform a deterministic photon number into a binomial photon-number distribution,^{17,23} which always remains sub-Poisson but approaches the Poisson boundary as the random deletion increases.²³ Mandel¹⁷ has shown that the information rate per symbol carried by such a counting channel will be greater than that for the Poisson channel, but will approach the latter as the photon-detection probability η approaches 0. We have shown elsewhere²⁴ that the performance of a simple binary on-off-keying photoncounting system, of the kind considered earlier, is also superior for a binomial source than for a Poisson source, approaching the latter as η approaches 0.

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