

Most probable paths in temporal weighted networks: An application to ocean transportEnrico Ser-Giacomi,¹ Ruggero Vasile,^{1,2} Emilio Hernández-García,¹ and Cristóbal López¹¹*IFISC, Instituto de Física Interdisciplinar y Sistemas Complejos (CSIC-UIB), E-07122 Palma de Mallorca, Spain*²*Ambrosys GmbH, Albert-Einstein-Strasse 1-5, 14473 Potsdam, Germany*

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We consider paths in weighted and directed temporal networks, introducing tools to compute sets of paths of high probability. We quantify the relative importance of the most probable path between two nodes with respect to the whole set of paths and to a subset of highly probable paths that incorporate most of the connection probability. These concepts are used to provide alternative definitions of betweenness centrality. We apply our formalism to a transport network describing surface flow in the Mediterranean sea. Despite the full transport dynamics is described by a very large number of paths we find that, for realistic time scales, only a very small subset of high probability paths (or even a single most probable one) is enough to characterize global connectivity properties of the network.

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I. INTRODUCTION

Many real-world systems can be studied by using the network paradigm [1,2]. Though in many of these situations the connections between network nodes are constant in time and it is sufficient to adopt a static network description, this approach may not be always suitable and a temporal network description is required [3]. Relevant examples are epidemic spreading, human communication, or transportation networks, i.e., systems where the strong time variability plays a crucial role in determining connections and interactions [4–6]. Also, approaching continuous dynamical systems from a network perspective [7,8] requires a spatiotemporal discretization that, in the nonautonomous case, often determines a marked time-dependence of the resulting networks. This feature is clearly recognizable, in particular, in climate networks [9–12] and in networks describing connectivity by fluid flows [13,14].

Prominent connectivity patterns in networks can be revealed by introducing the concept of paths [1,2]. While its definition is simple and intuitive for static or aggregated networks, for the temporal case paths between nodes can suddenly appear and disappear in time [3,15–22]. Thus, there is recently a focus toward the definition and characterization of paths in temporal networks. The concept of shortest path in static network analysis has been generalized to include information on the time necessary to establish a space-time connection between nodes. This was the motivation behind the development of fastest path analysis, specially for unweighted and undirected networks [3,20]. Although it is relevant to study the time required to build a path among two nodes, it is equally crucial to understand how to quantify the importance and the distribution of such paths. This issue becomes essential when one tries to exploit this information in order to define and evaluate global network properties, for example, betweenness centrality.

In this paper we extend the concept of most probable path (MPP) [23–27] to the case of temporal, weighted, and directed networks. We quantify the relative importance of the MPP with respect to the whole set of paths and to the subset of highly probable paths (HPPs) incorporating most of the connection probability. Using such sets of paths we are able to define a betweenness centrality measure for temporal weighted networks. The approach presented here is applied to a flow

network describing ocean water transport by surface currents in the Mediterranean Sea. In this example we demonstrate that information contained in MPPs (or in small subsets of HPPs) suffices to describe all major transport properties, despite the number of such paths being just a very small fraction of the full set. The MPPs correspond to the main carriers of ocean mass transport (showing connectivity patterns) and MPP betweenness to a measure for a clear identification of the main avenues of water transport in the Mediterranean Sea.

II. MOST PROBABLE PATHS

The analysis is restricted to a time interval $[t_0, t_M]$ in which $M + 1$ snapshots of the state of the network are taken at times $t_l = t_0 + l\tau$, $l = 0, 1, \dots, M$, with τ the time between them. We consider a temporal, directed, and weighted network of N nodes. Its time-dependent connectivity is described by a set of weighted adjacency matrices $\mathbf{A}^{(l)}$, ($l = 1, \dots, M$), in which the matrix element $\mathbf{A}_{IJ}^{(l)} \geq 0$ specifies the strength of connectivity from I to J during the time interval $[t_{l-1}, t_l]$. A convenient way to analyze the system is using *time-ordered graphs* (TOGs) [19]. Formally, the TOG can be considered a static network of $N \times (M + 1)$ nodes with directed and causal links. For each snapshot l , a group $V(t_l)$ of N nodes replicating the nodes of the original network can be defined. Links are then established only from nodes at successive times, i.e., from $i_{l-1} \in V(t_{l-1})$ to $j_l \in V(t_l)$ with the weights given by those in the original temporal network: $\mathbf{A}_{i_{l-1}j_l}^{(l)}$.

We now consider a flow or transport process by releasing independent random walkers in each node of the network. Their motion is assumed to be Markovian and is defined by single-step transition probabilities proportional to the entries in the adjacency matrices. Specifically the probability of reaching node k_l at time t_l under the condition of being at k_{l-1} at time t_{l-1} is

$$\mathbf{T}_{k_{l-1}k_l}^{(l)} \equiv \frac{\mathbf{A}_{k_{l-1}k_l}^{(l)}}{s_{\text{out}}^{(l)}(k_{l-1})}. \quad (1)$$

Here $s_{\text{out}}^{(l)}(k) = \sum_j \mathbf{A}_{kj}^{(l)}$ is the *out-strength* of node k during time step l , so that $\sum_j \mathbf{T}_{kj}^{(l)} = 1$. A generic M -step path μ

between two nodes I and J is defined as a $(M + 1)$ -uplet $\mu \equiv \{I, k_1, \dots, k_{M-1}, J\}$ providing a sequence of nodes crossed to reach J at time t_M from I at time t_0 . Under the Markovian assumption the probability for a random walker to take the path μ under the condition of starting at I is

$$(P_{IJ}^M)_\mu = \mathbf{T}_{Ik_1}^{(1)} \left[\prod_{l=2}^{M-1} \mathbf{T}_{k_{l-1}k_l}^{(l)} \right] \mathbf{T}_{k_{M-1}J}^{(M)}. \quad (2)$$

The MPP η_{IJ}^M is the path that maximizes Eq. (2) with respect to the intermediate nodes k_2, \dots, k_{M-1} . Its probability is denoted by $P_{IJ}^M = \max_\mu \{(P_{IJ}^M)_\mu\}$. The exact maximization of Eq. (2) can be obtained iteratively by noting that in the first step the maximum probability to reach a given node k_1 is simply $P_{Ik_1}^1 = \mathbf{T}_{Ik_1}^{(1)}$ and then using the recurrence

$$P_{Ik_{l+1}}^{l+1} = \max_{k_l} (P_{Ik_l}^l \mathbf{T}_{k_l k_{l+1}}^{(l+1)}) \quad (3)$$

for $l = 1, 2, \dots, M - 1$ until reaching $k_M = J$. This type of iterative optimization is similar (taking logarithms of the probabilities involved) to the one used to find optimal configurations of directed polymers in random media [28] and can be considered as an adaptation of the classical Dijkstra algorithm [29] to the layered and directed structure of the TOG. The computational cost of the maximization is strongly reduced by calculating first *accessibility matrices* [21] and restricting the maximization search to the set of nodes that are *accessible* from I and for which J results *accessible* as well.

To assess whether the MPP alone is a good representation of the transport dynamics we introduce the quantity $\lambda_{IJ}^M \equiv P_{IJ}^M / \sum_\mu (P_{IJ}^M)_\mu$. It corresponds to the fraction of probability carried by the MPP between I and J with respect to the sum of the probabilities of all paths connecting these nodes after M steps. Note that the sum in the denominator can be efficiently computed as the entry (I, J) in the matrix product $\prod_{l=1}^M \mathbf{T}^{(l)}$. Depending on the network under investigation, the MPP can actually carry a significative fraction of the total connection probability. When this is not the case we can relax the definition of MPP and define a subset of HPPs, which carry most of the probability. In particular we want to identify paths characterized by individually carrying a probability larger than a fraction ϵ of the MPP probability, i.e., larger than ϵP_{IJ}^M , with $0 \leq \epsilon \leq 1$. Exhaustively searching for all such paths becomes computationally prohibitive except for the smallest N and M values. Here we compute the set \mathcal{Q}_{IJ}^M of all paths of M steps between I and J that are constructed by joining the MPP from I to an intermediate k_l and the MPP from k_l to J . In principle there would be $(M - 2) \times N$ such paths, one for every choice of the intermediate k_l , but this number is in fact much smaller when considering that k_l should be accessible from I [21], that J should be accessible from k_l itself, and that many of the resulting paths turn out to be repeated. Out of these we consider the subset $\mathcal{K}_{IJ}^M(\epsilon) = \{\mu \in \mathcal{Q}_{IJ}^M | (P_{IJ}^M)_\mu > \epsilon P_{IJ}^M\}$. This set contains the MPP, and although it may miss some of the paths with probability larger than ϵP_{IJ}^M we expect it to contain a sufficiently representative sample of them. This can be checked by calculating $\lambda_{IJ}^M(\epsilon) \equiv \sum_{\nu \in \mathcal{K}_{IJ}^M(\epsilon)} (P_{IJ}^M)_\nu / \sum_\mu (P_{IJ}^M)_\mu$, the fraction of probability carried by this set of HPPs.

III. MPP BETWEENNESS

Equipped with the above definitions we can now characterize network properties that are dependent on optimal paths in different ways. One of these is the concept of *betweenness centrality*, which is generally defined as the proportion of shortest paths passing through a node. We introduce here a definition based on the number of *most probable paths* crossing a node. Specifically we define the betweenness of node K after M steps as $\mathcal{B}_K^M = \sum_{IJ} g_{IJ;K}^M / N_M$, where the sum is over all pairs of initial nodes I and final accessible nodes J , N_M is the total number of connected pairs of nodes at time step M (computable from accessibility matrices [21]), and $g_{IJ;K}^M$ is the number of times the node K appears in the most probable path connecting I and J . Fixing the time interval M corresponds to considering paths with the same temporal duration. In this way we ignore connections that are occurring at shorter or longer times [19] and that can be significantly more probable. It is possible to overcome this limitation by performing a multistep analysis: we can look at all MPPs with M in a given interval $[M_{\min}, M_{\max}]$ and choose the MPP, $\eta_{IJ}^{[M_{\min}, M_{\max}]}$, with the highest probability. The multistep analysis leads to an alternative definition of betweenness, i.e., a multistep MPP betweenness $\mathcal{B}_K^{[M_{\min}, M_{\max}]}$, which is calculated considering the multistep MPPs instead of the fixed- M one.

IV. MEDITERRANEAN FLOW NETWORK

We apply the previous concepts to a specific network describing surface water transport in the Mediterranean Sea [13,14]. Fluid transport is typically studied from a Lagrangian perspective, by following trajectories of particles released in the flow. Recent works have approached the problem from a discretized point of view [13,14,30–32]. Most of these studies are focused on the analysis and identification of coherent structures like vortices, barriers, or regions where trajectories of fluid parcels tend to be confined. Less insight is available on the pathways followed by fluid masses during transport and in the resulting connectivity patterns. In this regard, our approach complements the standard Lagrangian toolbox as will be illustrated with this oceanic flow example.

Velocity data have been collected from the Mediterranean forecasting system model, physics reanalysis component [33]. We use the horizontal surface velocity daily field of the whole Mediterranean basin at a resolution of $1/16$ degrees during 10 years of simulation (2002–2011). For the first months of each year a temporal flow network has been constructed by partitioning the surface of the Mediterranean sea in $N = 3270$ two-dimensional square boxes (approximate lateral size of 28 km), each of which is identified with a network node. Link's weights are assigned from the effective mass transport driven by ocean currents between two boxes during a given time interval [13,14]. To build the adjacency matrices at each time t_l ($l = 1 \dots M$) we homogeneously initialize $n = 500$ ideal fluid particles in each node and integrate the surface velocity for a given time τ . The matrix element $\mathbf{A}_{IJ}^{(l)}$ is the number of particles starting from node I at time t_{l-1} and arriving to node J at t_l . The normalized matrices $\mathbf{T}^{(l)}$ define a flow network.

We perform the calculations using a time step of $\tau = 10$ days, and considering M from 6 to 9 steps (i.e., in between 2

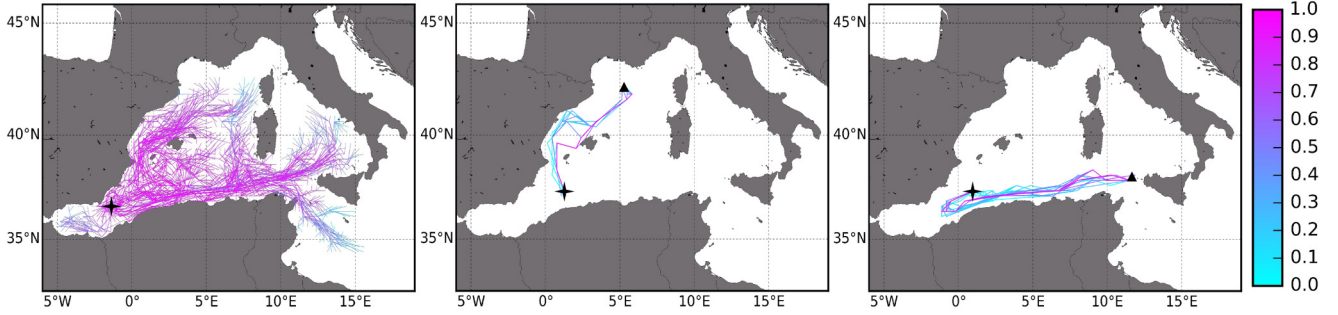


FIG. 1. (Color online) Paths of $M = 9$ steps (three months) in the Mediterranean flow network with starting date January 1, 2011, represented as straight segments joining the path nodes. Left: MPPs originating from a single node (black star) and ending in all accessible nodes. Color gives the P_{IJ}^M value of the paths in a normalized log-scale between the minimum value (10^{-15} , light turquoise) and the maximum (10^{-5} , dark pink). Center and right: all the paths in the $\mathcal{K}_{IJ}^M(\epsilon)$ set with $\epsilon = 0.1$, initial point marked by a cross and final point marked by a triangle. The center panel shows the 18 HPPs, out of a total of 54 276 paths between the two sites. The MPP, with $P_{IJ}^M = 3 \times 10^{-9}$, is displayed in dark pink, whereas the other paths are colored with a normalized logarithmic scale according to their $(P_{IJ}^M)_\mu$ values in $[\epsilon P_{IJ}^M, P_{IJ}^M]$. Right panel shows the 39 HPPs out of a total of 61×10^6 , in a similar logarithmic scale normalized in $[\epsilon P_{IJ}^M, P_{IJ}^M]$ with $P_{IJ}^M = 1.4 \times 10^{-6}$.

and 3 months, a time interval during which the horizontal-flow assumption remains a good approximation). We build ten temporal networks, each one having t_0 as January 1 of each of the years available in the simulation database (2002–2011).

V. RESULTS

In Fig. 1 (left panel) we show on map the set of all the MPPs of $M = 9$ time steps starting from a given node in one of our temporal networks (the one corresponding to 2011), and we notice how many different connections are possible from a single starting node. The P_{IJ}^M values span several orders of magnitude and this behavior is typical for the distribution of probability across MPPs. We stress here that MPPs do not coincide in general with fastest paths: the fastest connection among two nodes is not always the most probable one (see Appendix), stressing the importance of a weighted description of the network.

To assess how representative of the whole dynamics are MPPs such as the ones shown in Fig. 1, we show in Fig. 2 the distribution of λ_{IJ}^M and $\lambda_{IJ}^M(\epsilon)$ for two values of M . The distributions are collected from the λ values of the whole set of accessible pairs (I, J) in our ten temporal networks. For small M most of the MPP have significant λ values, but as M increases the peak in the distribution of λ_{IJ}^M shifts toward zero (we have checked that exponentially) as a consequence of the dramatic increase with M of the number of available paths between two nodes. Then, it becomes important to consider larger sets of paths such as $\mathcal{K}_{IJ}^M(\epsilon)$. For the cases plotted, i.e., $M = 5, 6$ and $\epsilon = 0.1$, the mean values of $\lambda_{IJ}^M(\epsilon)$ are around 0.5. This means that, despite $\mathcal{K}_{IJ}^M(\epsilon)$ may not contain the full set of paths that individually carry a probability larger than ϵP_{IJ}^M , it is large enough so that, for most of the (I, J) pairs, it contains globally over 50% of the connection probability. This result further gains meaning if we consider that the number of paths in $\mathcal{K}_{IJ}^M(\epsilon)$ is on average well below 1% of the total number of paths of $M = 5$ and 6 steps. Hence, despite the strong particle dispersion characterizing our flow networks it is true that only a small subset of paths contribute significantly to the main transport features. This conclusion will also show up when

studying global network properties, such as the betweenness centrality measure.

In Fig. 3 we show the multistep MPP betweenness $\mathcal{B}_K^{[6,9]}$, averaged over our ten networks. We have noticed that the distribution of the betweenness decreases exponentially at large betweenness values so that there are not strong hubs in the network. Spatial patterns determined by the transport dynamics of the flow are clearly evident in the figure where high betweenness areas are organized in one-dimensional-like structures corresponding to the main corridors of transport, i.e., narrow pathways that connect different regions of the ocean. Main paths of the Mediterranean Sea are observed like Cyprus and Rhodes Gyres, surrounding the Ionian basin, the Algerian current, and those along the Sicily strait, etc. Note that because of the ten-years average, individual short-lived mesoscale features (eddies and fronts) are averaged out. We note, however, that despite this interyear average, MPP

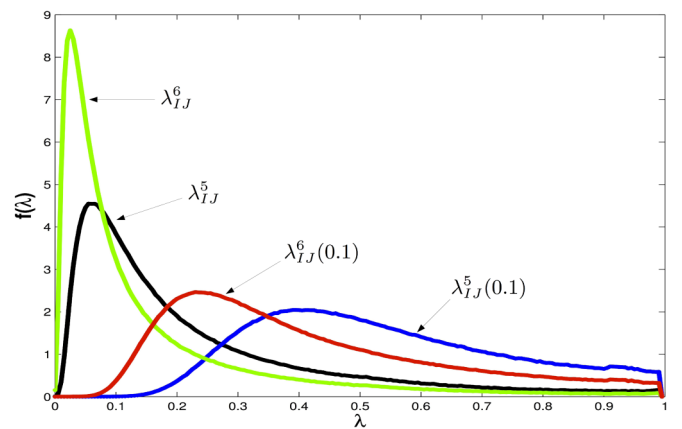


FIG. 2. (Color online) Normalized histogram $f(\lambda)$ of coefficients λ_{IJ}^M for $M = 5$ (black curve) and $M = 6$ (green curve), and $\lambda_{IJ}^M(\epsilon)$, with $\epsilon = 0.1$ for $M = 5$ (blue curve) and $M = 6$ (red curve). The statistics are compiled from all connected pairs of nodes (I, J) and the ten temporal flow networks corresponding to the first months of the ten years of velocity data. The mean values are: $\langle \lambda_{IJ}^6 \rangle = 0.24$; $\langle \lambda_{IJ}^5 \rangle = 0.16$; $\langle \lambda_{IJ}^6(0.1) \rangle = 0.52$; $\langle \lambda_{IJ}^5(0.1) \rangle = 0.42$.

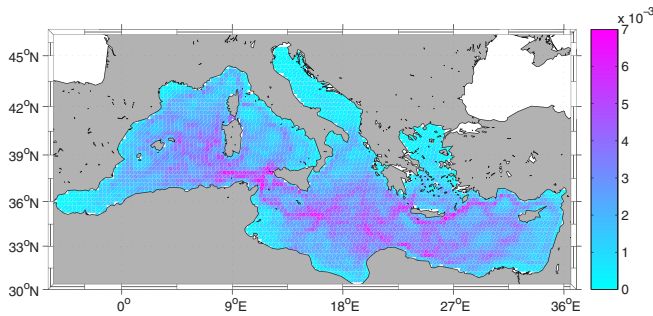


FIG. 3. (Color online) Multistep MPP betweenness $\mathcal{B}_K^{[6,9]}$ at each geographical node K , computed for each of our ten (2002–2011) temporal networks and then averaged over them.

betweenness maps remain dependent on the starting date t_0 . These observations highlight the seasonal and interannual variability of the flow and justifies further the need of a time-dependent network approach as opposed to a fully averaged static network description.

The robustness of our methodology has been tested by checking the stability of our results on MPP betweenness under different conditions. First, dealing with temporal networks it is important to understand how much the results are affected by the choice of the time-step duration τ [34]. We checked this issue considering different M and τ values but the same total duration. Results for the MPP betweenness remain nearly unchanged when keeping a constant total duration $M\tau$, confirming robustness under variations of the temporal resolution. Second, we noted that the MPP betweenness does not significantly change when computed with just the 50% of MPPs with larger values of λ (i.e., when using a threshold to retain only the most significant MPPs). Finally, to support our interpretation of most probable paths as main carriers of connectivity, we considered also betweenness calculated from HPP subsets so that $g_{I,J;K}^M$ is now the number of times node K appears in the set $\mathcal{K}_{I,J}^M(\epsilon)$ of HPPs between I and J , with $\epsilon = 0.1$. We did not appreciate relevant differences between this calculation and the one involving only MPPs. Indeed the Pearson correlation coefficient between the two betweennesses is larger than 0.9. Hence, despite the MPPs represent a small portion of the paths in the $\mathcal{K}_{I,J}^M(\epsilon)$ subsets (between 3 and 10% for $\epsilon = 0.1$, depending on the value of M), which is itself a very small fraction of the full set of paths in the network, they seem to be representative of the main spatiotemporal structures describing the global dynamics. Indeed, center and right panels of Fig. 1 show that most of the relevant paths remain spatially close to the MPP. This observation is confirmed by calculations of the spatial dispersion between paths in $\mathcal{K}_{I,J}^M(\epsilon)$, whose average turns out to be of the order of the size of the boxes defining the nodes.

VI. CONCLUSIONS

In this paper we introduced tools to compute highly probable paths in weighted temporal networks and to evaluate their relative importance. Betweenness centrality measures based on them have also been introduced. We applied this approach to characterize connectivity in the Mediterranean Sea from

a network-theory perspective. Here, MPPs correspond also to the set of paths that maximize the fraction of transported mass, giving therefore a clear physical interpretation of connection probabilities. Despite MPPs representing only a small fraction of the whole set of paths, we found that they suffice to highlight the main transport pathways across our network, since most of the HPPs remain geographically close to them. This means that paths followed by fluid masses connecting two regions are organized in elongated narrow tubes centered on the MPP.

We believe that the study of fluid transport as a network will provide new tools and insights complementing standard Lagrangian methods. Indeed most of these are devoted to the identification of barriers to transport or coherent regions with small fluid exchange with the surroundings. Here we are instead addressing the opposite question: how to detect regions and pathways that maximize fluid interchange across the network. Even though in principle pathways are simply given by trajectories, it is almost impossible to extract clear and significant patterns from the complex sets of trajectories that arise in all, except the most simple, time-dependent flows. Our approach allows us to quantify explicitly transport among two subregions of the domain, highlighting the optimal path connecting them. In this sense MPP analysis differs from simply studying the evolution in time of tracer concentrations seeded in a given region.

Beyond the fluid dynamics context, MPPs and the MPP-betweenness measure here introduced could be easily transferred to other kinds of weighted temporal networks. This could be relevant, for instance, in defining vulnerability metrics in disease spreading processes, individuate critical nodes in biological and ecological networks, or in detecting bottlenecks of reaction chains in metabolic networks.

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APPENDIX: COMPARING FASTEST AND MOST PROBABLE PATHS

In the study of temporal networks, the concept of *fastest* path has been put forward as a natural extension of the *shortest path* of static networks. In our work we define and analyze a different type of relevant path, which is the MPP. It is important to address the differences between most probable and fastest paths, and we do so in this Appendix.

The MPP refers to the path transporting the maximum fraction of water (or of probability) between two nodes, and the fastest path to the pathway linking the two nodes in the shortest time. This second concept cannot be implemented when the number of time steps is fixed. However, we can reclaim the concept of fastest path within a multistep approach, i.e., by looking at a time window specified by a range of values for the number

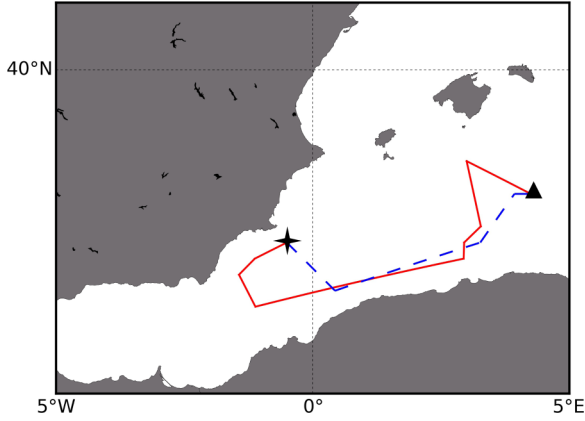


FIG. 4. (Color online) We show the *fastest* MPP (dashed blue line) and the *absolute* MPP (continuous red line), between an origin node I (black star) and a destination node J (black triangle). The considered full set of paths ranges from $M = 3$ to $M = 9$ steps of ten days. The *fastest* MPP reaches the destination node in four steps of $\tau = 10$ days while the *absolute* MPP needs eight steps, i.e., 40 days more. The probability associated to the *fastest* MPP is 5.9×10^{-7} and the probability of the *absolute* MPP is 6.7×10^{-6} .

of time steps M . We can then define the set $\mathcal{M}_{IJ}^{[M_{\min}, M_{\max}]}$ of $(M_{\max} - M_{\min} + 1)$ MPPs for the pair I, J for $M \in [M_{\min}, M_{\max}]$, and the *fastest* MPP as the MPP in $\mathcal{M}_{IJ}^{[M_{\min}, M_{\max}]}$ corresponding to the smaller M . On the other side we can also define an *absolute* MPP, i.e., the MPP in $\mathcal{M}_{IJ}^{[M_{\min}, M_{\max}]}$ having the highest probability. By comparing the set of *absolute* MPPs with the set of *fastest* MPPs we can address the question: is the fastest path necessarily the most probable?

In Fig. 4 we show that the *fastest* MPP among two nodes of the network is different from the *absolute* MPP and that its probability, in several cases, can be orders of magnitude smaller. We considered for this example paths ranging from $M = 3$ to $M = 9$ steps of ten days (i.e., a time scale of 1–3 months) with starting date January 1, 2011. The results show the importance of distinguishing between the connections realized in the shortest time and the connections that carry most of the transported mass (the most probable).

To display in a more systematic way the differences between *fastest* and *absolute* MPPs across the network, we study the rank plot of the whole set of paths during ten years (2002–2011) ranging from three to nine steps of ten days starting at January 1 of each year (see Fig. 5). The rank plot displays the probabilities of each path in one of the sets sorted in decreasing order. We see a gap in probabilities between the two sets of about one order of magnitude in most of the range displayed. The *fastest* MPPs have probabilities significantly smaller than *absolute* MPPs.

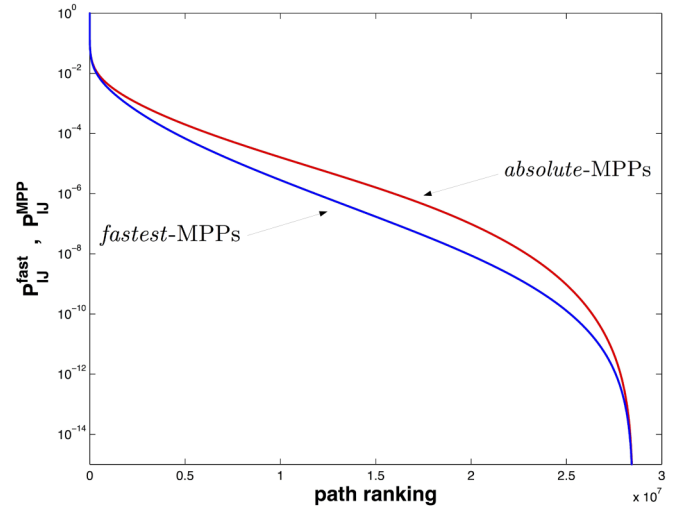


FIG. 5. (Color online) Ranking plot in which P_{IJ}^{fast} (blue line) correspond to the probability of *fastest* MPPs and P_{IJ}^{MPP} (red line) correspond to the probability of *absolute* MPPs. The range of probability values can be read from the vertical axis (logarithmic scale). The total number of optimal paths can also be read off from the horizontal axis.

Finally, we also evaluate how these differences are reflected in the betweenness measures. We define the relative difference among the betweenness computed using the set of *fastest* MPPs and *absolute* MPPs for the node K as

$$\Delta_K = 2 \frac{\mathcal{B}_K^{\text{abs}} - \mathcal{B}_K^{\text{fast}}}{\mathcal{B}_K^{\text{abs}} + \mathcal{B}_K^{\text{fast}}}, \quad (\text{A1})$$

where $\mathcal{B}_K^{\text{abs}}$ is the betweenness computed using *absolute* MPPs and $\mathcal{B}_K^{\text{fast}}$ the betweenness computed using *fastest* MPPs. We consider once more paths ranging from three to nine steps of ten days with starting date January 1, 2011, and we compute the spatial average (i.e., the average over nodes K) for the absolute value of the relative difference finding $\langle |\Delta_K| \rangle_K = 0.32$. This means that, on average, the difference between the two measures is about 30%.

We stress that all the comparisons above are among paths that are already MPPs linking a pair of nodes. Considering still fastest paths (for example, the one by which the very first particle from one node reaches the other) will lead to much stronger differences. In summary, the results show the importance of distinguishing between the connection realized in the shortest time and the connection that carries most of the transported mass. This gains even more relevance when considering possible applications, such as rescue operations, pollutant-spreading, or biological connectivity.

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