Biased and greedy random walks on two-dimensional lattices with quenched randomness: The greedy ant within a disordered environment

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The main characteristics of biased greedy random walks (BGRWs) on two-dimensional lattices with real-valued quenched disorder on the lattice edges are studied. Here the disorder allows for negative edge weights. In previous studies, considering the negative-weight percolation (NWP) problem, this was shown to change the universality class of the existing, static percolation transition. In the presented study, four different types of BGRWs and an algorithm based on the ant colony optimization heuristic were considered. Regarding the BGRWs, the precise configurations of the lattice walks constructed during the numerical simulations were influenced by two parameters: a disorder parameter ρ that controls the amount of negative edge weights on the lattice and a bias strength B that governs the drift of the walkers along a certain lattice direction. The random walks are "greedy" in the sense that the local optimal choice of the walker is to preferentially traverse edges with a negative weight (associated with a net gain of "energy" for the walker). Here, the pivotal observable is the probability that, after termination, a lattice walk exhibits a total negative weight, which is here considered as percolating. The behavior of this observable as function of ρ for different bias strengths B is put under scrutiny. Upon tuning ρ , the probability to find such a feasible lattice walk increases from zero to 1. This is the key feature of the percolation transition in the NWP model. Here, we address the question how well the transition point ρ_c , resulting from numerically exact and "static" simulations in terms of the NWP model, can be resolved using simple dynamic algorithms that have only local information available, one of the basic questions in the physics of glassy systems.

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I. INTRODUCTION

The purely stochastic motion displayed by a simple random walk (RW) and its various extensions and modifications that, e.g., go under the popular synonyms "the ant in a labyrinth" [1,2] and "the biased ant in a labyrinth" [3], give rise to intriguing effects such as (anomalous) diffusion, drift, and trapping. These, and further related topics, were actively studied since the beginning of the last century, first by analytic means and later using extensive numerical simulations. A subset of those studies, relevant in the context of the presented article, is reviewed in Sec. III. In all of these studies, a random walker is considered that traverses a given (disordered) environment subject to different dynamics-governing rules.

Here, we also consider a random walker on a disordered "energy" landscape, where, in contrast to most previous studies, the underlying energy values might have a positive or a negative sign. Regarding the dynamics of the random walker, we study different types of algorithms that are discussed in detail in Sec. II. The basic principle underlying all these algorithms is that in the course of exploring the lattice, the random walker evaluates the energy values associated with all the edges incident to the site it currently resides on in order to determine its next stepping direction. More precise, we consider "greedy" RWs for which the locally optimal choice is to preferentially traverse edges with a large negative weight. Here, this is associated with a net gain of "energy" for the walker, directly proportional to the weight of the traversed edge (correspondingly, traversing an edge with positive energy is associated with a decrease of the energy resources of the walker). We further introduce precise rules for performing a "biased" RW, causing the walker to choose a stepping direction that might not coincide with the best possible immediate choice. Note that the existence of the negative energies fundamentally changes the behavior of the model defined below and will be destroyed by lifting all energies by the same amount to a positive value. The reason for this fundamental difference is that, in the presence of a suitable amount of (nonreplenishable) negative energies, the walker might lower its energy more and more by performing longer and longer walks, in contrast to the presence of only positive energies. The energy values are drawn from specified distributions wherein a disorder parameter ρ allows us to alter the fraction of negative energies in the environment (see the discussion below).

Regarding such a disordered environment, we have previously introduced [4] and investigated [5–9] the negativeweight percolation (NWP) problem by numeric means. Regardless of the spacial dimension of the underlying (hypercubic) lattice graph, the observables in the NWP problem are *simple* lattice paths with an intrinsic dimension of d = 1. Therein the term simple means that, along such a lattice path, no node is repeated. As observables in the NWP problem one considers, e.g., minimum-weight paths (in the presence of a possibly empty set of negative-weighted loops) that span the underlying lattice between a pair of specified boundaries. The problem of finding these paths numerically can be cast into a minimum-weight perfect matching problem, e.g., discussed in Refs. [6,8] in more detail (in the latter reference, a

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comprehensive description and illustration of the respective algorithm can be found). A pivotal observation is that, as a function of the disorder parameter ρ , the NWP model features a disorder-driven phase transition, at which the actual minimumweight path energy becomes negative. In this question, note that the aforementioned transition in the path energy is a precursor to a further "geometric" phase transition at which the paths exhibit a roughness of the order of the systems size and where loops might appear that span the system along at least one direction [4]. In the limit of large system sizes, there is a particular value of the disorder parameter, signified as ρ_c , at which paths with a negative energy appear for the first time. For example, for the square (2D) lattice, considering a bimodal disorder distribution, the average path energy turns negative above $\rho_c = 0.0869(2)$ [4]. Therein, the finite-size scaling behavior is described by the critical exponent v = 1.47(6), consistent with the one that describes the geometric transition of the loops-only setup. At $\rho = 0.1032(5)$, slightly above this critical threshold, the fractal dimension of the paths was estimated to be $d_f = 1.268(1)$. Further, at this critical point, the loops that one finds in addition to the minimum-weight path are rather small and isolated. Hence, they (presumably) do not

affect the precise configuration of the minimum-weight paths. At this point, note that the NWP model is a prototypical disordered model system of intrinsic interest. As an example that makes use of the NWP problem statement, one might imagine an agent that travels along the edges of a graph. While traversing an edge, the agent either needs to pay some resource (whenever the weight of the traversed edge is positive) or he is able, once per edge, to harvest some resource (whenever the weight of the traversed edge is negative). Now, if the intention of the agent is to gain as much resources as possible, paths or loops obtained in the context of the NWP problem can serve as a guide to find routes along which the agent might move so as to optimize his yield. Further, the 2D variant of the NWP problem is interesting from a technical and algorithmic point of view. As regards this, the problem of finding ground-state spin configurations for the 2D random-bond Ising model, including the canonical Ising spin glass, on a planar triangular lattice can be mapped to the 2D NWP problem on a honeycomb lattice [10]. Further, paths including edges with possibly negative weights also appear in the context of domain wall excitations in 2D random bond Ising systems [11,12]. Thus, the NWP problem might serve to gain insight concerning the behavior of more realistic disordered systems.

As mentioned above, NWP was previously studied numerically exactly by employing a mapping to an auxiliary minimum-weighted perfect matching problem, which relies on an involved optimization algorithm specifically tailored for the problem at hand. Instead of using such a complex algorithmic approach that finds the optimum of a given optimization problem using *global* information, we here address the question how well the aforementioned transition point ρ_c (resulting from numerically exact simulations that involve a high degree of optimization) can be resolved using simple dynamic algorithms that have only *local* information about the assignment of edge weights available. Therein, so as to (locally) conform to the NWP optimization problem, we consider random walks which are "greedy" in the sense that the walker preferentially traverses edges with a negative weight. This is a more physical viewpoint on the problem, where, typically, the movement of particles is mostly determined by local information. Hence, a comparison between global static and dynamical physical viewpoints is possible. These two perspectives are often taken for the study of disordered and glassy systems [13–15]. Therefore we consider different biased random-walk approaches and study their ability to find system spanning paths with negative path weight. Averaging over different realizations of the disorder, we compute the probability to find a negative-weighted path as a function of the model parameter ρ for different simple dynamics. The algorithms we consider here all mimic moving particles (also referred to as agents or walkers) in a disordered environment. The simplest dynamics we study stems from a biased random walk on the lattice and the most intricate dynamics is implemented in terms of a particular heuristic algorithm known as "ant colony optimization." Subsequently, we will refer to $\rho_{\rm c}$ as the "dynamic" transition points, as opposed to the result from the "static" (global) simulations quoted above. One might allude to this study as being dedicated to the behavior of "the greedy ant in a disordered environment."

The remainder of the presented article is organized as follows. In Sec. II, we introduce the model in more detail and we outline the different dynamic algorithms used to perform the lattice walks. In Sec. III we review some of the related literature that alludes to effects that we also expect to observe within our simulations. In Sec. IV, we list the results of our numerical simulations and in Sec. V we conclude with a summary.

II. MODEL AND ALGORITHMS

In the present article we consider lattice graphs G = (V, E)with a 2D square lattice geometry, having side length Land periodic (free) boundary conditions (BCs) in the vertical (horizontal) direction. The considered graphs have N = |V| = L^2 nodes (plus two extra "outer" nodes) $i \in V$ and M = |E| =L(2L - 1) undirected edges $\{i, j\} \in E$ that join adjacent nodes $i, j \in V$ on the lattice, see Fig. 1. In addition, there are 2L extra edges that attach the L nodes on the left and right boundary to the respective outer nodes. We further assign a weight ω_{ij} to each $\{i, j\} \in E$. These weights represent quenched random variables that introduce disorder to the lattice. Here we consider independent identically distributed weights drawn from (i) the *bimodal* distribution,

$$P_{\rm bi}(\omega) = (1-\rho)\delta(\omega-1) + \rho\delta(\omega+1), \tag{1}$$

where ρ signifies the fraction of negative edge weights $\omega = -1$, or (ii) the *semicontinuous* distribution,

$$P_{\rm sc}(\omega) = (1 - \rho)\delta(\omega - 1) + \rho u(-1, 0), \tag{2}$$

wherein u(-1,0) signifies a random number uniformly drawn from the interval (-1,0].

Subsequently, we will study lattice walks on instances of weighted 2D lattice graphs with side length up to L = 800. The considered walks are confined to start at an outer node attached to the left lattice boundary and end as soon as the walker reaches an outer node attached to the right boundary; see Fig. 1. More precisely, to construct the lattice walks, two



FIG. 1. Example of a lattice walk on a disordered 2D square lattice of side length L = 4 and open boundaries in the vertical direction (note that in the actual simulations, periodic boundary conditions in this direction are considered). The solid (dashed) lines indicate edge weights 1 (-1) and the edges that connect the outer nodes (open circles) to the nodes on their respective boundaries carry zero weight. The fraction of negative edge weights it thus $\rho \approx 0.29$. The bold gray line indicates a possible exemplary lattice walk of weight $\omega_{\text{path}} = -2$.

classes of algorithms were used: *biased and greedy RWs* and an *ant colony optimization* (ACO) heuristic.

A. Computing paths via biased and greedy random walks (BGRWs)

In brief, the biased RWs explore the lattice graph while greedily choosing the "best" direction at each discrete time step. This might be the lowest weighted edge connected to the node the walker currently resides on or, if there is a draw, one of the lowest weighted edges picked at random. By "biased RW" it is meant that following a certain rule, a preferred direction is chosen for the next step, even if it is not the best possible immediate choice. The rule employed during our simulations reads as follows: Say the preferred direction is given by the positive x direction. First, the walker determines the direction corresponding to the best local edge choice. If this direction is already in the positive x direction, the walker goes there. Otherwise, if the best local choice does not lead towards the positive x direction, the walker accepts the choice only with probability 1 - B, where $B \in [0, 1]$ (termed "bias probability" or short "bias"), that is, if its not the best local choice, the walker goes towards the positive x direction with probability B.

Thus, a walker experiences a drift that, depending on the bias direction, effectively "guides" him or her towards a particular direction (see discussion in Sec. III). Here, we will refer to this kind of RW in a disordered environment as "biased and greedy random walk" (BGRW). While traversing the graph, the walker keeps track of the edge weights it encounters. Consequently, a lattice walk is considered successful if the total sum of the edges traversed by the walker is smaller or equal to zero. Clearly, the walker will be able to find negative-weight paths only if they exist, which can be detected by the previously mentioned exact algorithms. We anticipate that there will be an intermediate range, where negative-weight paths exist but are so rare that they cannot be detected by the locally acting walker.

We also study a variant of such lattice walks, where the walker interacts with the graph by modifying the weights of the traversed edges. This is to some extent motivated by studies, as, e.g., the one reported in Ref. [16], aimed at mimicking the foraging behavior of social monkeys. Here, we chose the rules that the walker

(i) replaces both positive and negative edge weights by weight $\omega = 0$ (aimed at modeling a finite nonreplenishable amount of both resources and costs) or

(ii) only replaces negative edge weights with the standard positive value of $\omega = 1$ (aimed at modeling nonreplenishable finite resources, only).

These types of modifications were chosen to also prevent the walker from getting stuck in regions of the lattice with a high density of negative edge weights, as for RWs in disordered energy landscapes at low temperatures, reviewed in Sec. III above. If no such interaction with the lattice is implemented, the walker could, in principle, collect unlimited amounts of negative edge weights.

Four distinct variants of BGRWs, termed RWs of type A, B, C, or D, were considered to evaluate which one performs better in finding paths with an overall negative weight. Here, the term "better" means "with higher probability for a given disorder parameter value ρ ." Also bear in mind that, ultimately, the aim is to quantify to which extent the transition point ρ_c (known from numerical simulations that involve finding the exact global optimum [4]) can be resolved using simple dynamic algorithms that have only *local* information about the assignment of edge weights available. The four variant of biased RWs, along with a brief description of the dynamic governing rules, are listed below. The variant A is based on standard RWs, while the other three variants are based on (partially) loop erased RWs.

1. BGRWs of type A

This type of walk is given by a BGRW that replaces negative edge weights [rule (ii) above] and maintains a running sum of the edge weights encountered while traversing the graph.

For completeness, note that also the variant wherein the walker replaces edge weights according to rule (i) was considered; see Fig. 2. In the figure, type A walks on a square lattice of size L = 100 for different values of the disorder parameter ρ are shown, considering the bimodal disorder distribution $P_{\rm bi}(\omega)$ [see Eq. (1)] for different fractions of negative edge weights ρ .

While traversing the lattice graph, the walker modifies positive and negative edge weights, replacing them by edge weights with value zero. As a consequence, for a small fraction of negative edge weights [as for $\rho = 0.1$ in Fig. 2(a)] it can be observed that the walker backtracks along subpaths of considerable length. This is due to the fact that, after the walker modified the respective edge weights, they are actually "cheaper" than the +1 edge weights that characterize the majority of the edges. The bias is designed to induce a drift along the positive *x* direction; hence, one finds a pronounced "channeling effect" along the horizontal lattice axis. Once such a region of zero-weighted edges develops, for small values



FIG. 2. Examples of type A lattice walks on a disordered 2D square lattice (considering bimodal disorder) of size L = 100 for different values of the disorder parameter ρ . In this example, while traversing the lattice graph under a bias parameter B = 0.05, the walker modifies positive and negative edge weights, replacing them by edge weights with value zero [i.e., rule (i) described in the text]. The subfigures correspond to (a) $\rho = 0.1$, (b) $\rho = 0.5$, and (c) $\rho = 0.9$.

of ρ , the walker has a hard time to escape from it. As is evident from Figs. 2(b) and 2(c), the channeling effect is less pronounced as the fraction of negative edge weights increases. Note that this effect is conceptually similar to the behavior of RWs in energetically disordered lattices, reviewed in Sec. III (see Refs. [17,18]), where initially an effective activation energy must be overcome in order to escape from a local minimum in the energy landscape. The general characteristics of such a lattice walk, basically confined to take place on edges of zero weight in an environment that mostly features positive edge weights, further bears some resemblance to the trapping behavior observed for RWs on percolation clusters discussed previously.

2. BGRWs of type B

This type of walk is given by a loop-erased BGRW, traversing the graph without modifying the edge weights (which is not necessary, due to the erasure of the loops; some details on loop-erased random walks (LERWs) are given in Sec. III; see also Ref. [19]). The sum of edge weights along the loop-erased walk is computed as soon as the BGRW has finished. Albeit the information processed by the walker is local in space, loop erasure can imply a long-term temporal memory of the walker, which might be interpreted as being nonlocal in time.

3. BGRWs of type C

This type of walk combines an interaction with the environment (as for type A walks) with loop erasure (as for type B walks). The lattice walk is constructed via the following two-step procedure: In the first step, the walker performs a BGRW and modifies the edge weights related to the traversed edges according to rule (i) above. At each visited node it further stores the sum of edge weights accumulated so far, the previous node it has visited, and the original weight (i.e., the value of the edge weight before it was modified by the walker) of the edge it has just traversed. In the second step, after the BGRW has terminated, the loop erasure and subsequent path evaluation is performed. Therefore, the lattice path is traced back, starting from the final node, by selecting at each intermediate node the step that yields the lowest sum of the edge weights until that point. This selection appears each time when a loop is encountered while backtracking. This is a heuristic which leads to lower path weights but is not able take combinatorial cases like a loops inside loops inside loops, etc., into account, since this would contradict the idea of using fast and local heuristic algorithms for computing the walks. Finally, the total weight of the path is computed by summing up the edges. For completeness, note that also the variant wherein the walker replaces edge weights according to rule (i) was considered.

4. BGRWs of type D

This type of walk might be referred to as a "partially" loop erased BGRW, wherein only unfavorable loops along the lattice walk are erased. More precisely, the walker modifies the weights of the traversed edges and checks each time if the node it currently resides on was visited previously. If so, the walker checks whether the created loop has a negative or positive weight. The loop is kept as part of the lattice walk if it has a negative weight, or, if it has a positive weight, the loop is discarded and the original edge weights are restored along it.

Regarding these dynamics, only the variant where the walker replaces edge weights according to rule (ii), i.e., the modification of negative edge weights only, was considered. This was done since the modification of positive and negative weights led to the walker being stuck for a very long time in a certain part of the graph, especially when using a low value for the bias and a small fraction of negative weights. The cause for this was that the walker created areas of zero weight in the graph that it could not escape efficiently if the surrounding edges were all positive and the bias was too low (somehow similar to the effective activation energy that must be overcome in the initial phase of RWs in energetically disordered lattices reviewed in Sec. III). A similar horizontal "channeling" effect can be seen in Fig. 2 for random walks of type A BGRWs, where the walker is constrained to move on an almost straight line because of the "repulsive" effect of the positive edge weights.

B. Computing paths via ant colony optimization (ACO)

Previously we considered algorithms that took into account *local* information only, that is, the *random walk* approaches we considered so far made a decision on where to go next by only considering the weights of the edges adjacent to the walkers current location.

In this subsection we will consider a nature-inspired, population-based heuristic for the solution of optimization problems called *ant colony optimization* (ACO) [20,21]. Designed to mimic the foraging behavior of actual ants, the ACO heuristic is a particular example of swarm intelligent

systems, which has proven to be valuable in solving a wide range of optimization problems that can be cast into the form of optimal-path problems. As regards this, the ACO heuristic has already been applied to different optimization problems such as, e.g., the traveling salesperson problem [22], which is a hard optimization problem, and the minimum-weight spanning tree problem [23], which is polynomially solvable, i.e., easy.

The ant *colony* represents a population of, say, *M* individual agents that are all able to construct solutions to the given problem by considering local information only (these need not necessarily be solutions of high quality). Albeit the *M* agents do not interact directly, they are able to interact indirectly through the deposition of *pheromone* on the graph edges (see discussion below). These interactions should lead to lower-energy solutions as compared to the single random-walk algorithms discussed above.

Here, the problem is to find a loopless minimum weight path from a specified source node *s* to a specified target node *t* in a weighted undirected graph G = (V, E), as discussed earlier. The ACO heuristic comprises an iterative algorithm, where in the beginning, there is no pheromone on the edges $\{i, j\} \in E$. Hence, the only information associated with a particular edge is its weight ω_{ij} . A basic variant of the ACO algorithm for the above problem can be cast into the following steps:

(i) Preprocessing of the edge weights: Transform the set of edge weights ω_{ij} to a set of (initial) transition probabilities $\tau_{i \rightarrow j} = Z_i^{-1} \exp\{-\omega_{ij}\}$, where $Z_i = \sum_{j \in \text{Nb}(i)} \exp\{-\omega_{ij}\}$ with Nb(*i*) specifying the set of nodes adjacent to node *i*. In this way it is assured that an edge weight ω_{ij} with a comparatively large negative weight will result in a comparatively high probability for the agent to take the step $i \rightarrow j$. These initial transition probabilities can also be thought of as an initial distribution of pheromone which affect the behavior of the individual agents.

(ii) Generate solution to the problem: The ability of an individual agent to construct loopless paths as solution to the problem at hand can be implemented in various ways. For a more clear illustration, let us consider one agent only, i.e., M = 1. Here, starting at a specified source node *s*, we let the agent perform a LERW, guided by the transition probabilities $\tau_{i \rightarrow j}$. During the loop erased walk, the agent sums up the edge weights along the LERW edges. As soon as the agent reaches a specified target node *t*, the guided LERW is completed and a solution to the problem, i.e., a loopless (*s*,*t*) path, is obtained. Note that this is not necessarily a solution with a good quality regarding the given optimization criterion.

(iii) Evaluate quality of the solution and deposit pheromone: In order to quantify the quality of the LERW obtained in step (ii) and so as to reflect the optimization criterion of the problem at hand, we compute the "fitness" parameter $q = \alpha(1 - \omega_p/\ell_p)$, where $\alpha > 0$ is a tunable parameter, while ω_p and ℓ_p are the weight and length of the LERW, respectively. Note that for edge weights drawn from a bimodal distribution, it holds that $q \ge 0$. Further, note that the more negative the sum of the edge weights, the larger the value of q. Then modify the transition probabilities along the LERW to $\tau_{i \rightarrow j} \rightarrow \tau_{i \rightarrow j} + q$ and normalize them as in step (i) above. This is called a delayed pheromone update, since the local information on the edges is updated after the LERW is obtained. The parameter α can be tuned to alter the influence of the quality measure q on the subsequent recomputation of the transition probabilities. Here, we optimize the ratio ω/ℓ , effectively measuring the average edge weight that contributes to the path. This quantity was chosen since it naturally normalizes the value of q to the range $q \in [0, 2 \cdot \alpha]$ and thus allows us to control the influence of the current local solution on the transition probabilities in an easy way by using the parameter α (in contrast, note that in the static NWP problem we optimize ω instead).

(iv) Evaporate pheromone: While pheromone is deposited only on the edges of particular loopless paths, it can be useful to also evaporate some of the pheromone from all edges. For example., a simple evaporation heuristic is to modify the transition probabilities along all edges to $\tau_{i \rightarrow j} \rightarrow \beta \cdot \tau_{i \rightarrow j}$, where $\beta \in [0,1]$, and to normalize them as in step (i) above.

Step (i) is a preprocessing step to generate transition probabilities that reflect the weight assignment on G. Steps (ii) and (iii) complete the "life cycle" of an individual agent: It generates a (nonoptimal) solution to the optimization problem and evaluates a fitness for the solution. The fitness is then used to alter the transition probabilities which are associated with the edges. Instead of using just one agent, steps (ii) and (iii) straightforwardly generalize to M individual agents. The construction and evaluation of M LERWs and the subsequent pheromone evaporation, i.e., steps (ii)-(iv) comprise one sweep. Consequently, the ACO algorithm is iterated for a number of, say, n sweeps. The transition probabilities change from sweep to sweep. Thereby, edges that belong to (s,t)paths with a high fitness get equipped with a larger transition probability. Albeit this induces a positive feedback that allows us to distinguish "good" paths, the parameter α can be used to limit the extent to which the transition probabilities are modified from sweep to sweep (during our simulation we used $\alpha = 1/M$). Also note that the parameter β serves to limit the "long-term memory" of the search process in order to efficiently explore a large variety of (s,t) paths (during most of our simulations we used $\beta = 0.98$). For example, setting $\alpha = 0$ and $\beta = 1$ completely suppresses the deposition and evaporation of pheromone. For that parameter choice the individual agents would construct LERWs which are guided only by the transition probabilities induced by the edge weights on the lattice [following step (i)] and they would not be able to indirectly communicate through the deposition of pheromone. Also note that in the extreme case where $\alpha \to \infty$ and $\beta = 1$, the edges $\{i, j\}$ that belong to a (s, t) path found in step (ii) would get a transition probability $\tau_{i \rightarrow j} \approx 1$ and all transition probabilities that relate to those edges that depart from that particular path would effectively be zero. Hence, once an agent crosses that path (during a later sweep), it is immediately "confined" to that path and is unlikely to escape from it. During a simulation run for a particular realization of the edge weights we store the best path found so far. This path is returned by the ACO algorithm after it terminates.

III. REVIEW OF RELATED LITERATURE

A typical observable for the characterization of ordinary RWs is given by their mean-square displacement (MSD) $\langle R(t)^2 \rangle$, which, as a function of time *t*, is expected to scale as $\langle R(t)^2 \rangle = Dt$. Therein, *D* signifies the diffusion coefficient characteristic for the observed walks.

A particular extension of the simple RW model, relevant in the context of the presented study, consists in the addition of disorder to the environment in which a walk is performed. Different approaches to implement disorder can be found in the literature: For example, one might consider percolation-like "occupied vs. free" site (or bond) disorder [24-29], in which case a fraction p of lattice sites (or bonds) is distinguished as "occupied" and where a RW is performed on a cluster of occupied nearest neighbor sites (a.k.a. "the ant in the labyrinth"). For short walking times, i.e., times within which the walker cannot fully trace the cluster, and in the vicinity of the critical point p_c , signaling the onset of percolation, anomalous diffusion was observed [25,30]. Thereby, the MSD of the RW scales according to $\langle R(t)^2 \rangle \propto t^{\alpha}$, characterized by an effective dimension $d_{\rm RW} = 2/\alpha$. At the critical point for site dilution on 2D lattices, Ref. [24] reports $d_{RW} = 2.68(5)$, i.e., $\alpha \approx 0.75$ indicating a subdiffusive behavior. As pointed out in Ref. [31], considering a percolation-like "occupied vs. free" disorder of the underlying lattice, the effective dimension of the RWs increases from $d_{RW} = 2$ in dimension d = 1 to the mean-field value $d_{RW} = 6$ for $d \ge 6$. Due to the trivial nature of 1D percolation, which inhibits anomalous diffusion, one obtains the scaling dimension that characterizes unhindered random walks for the 1D setup. Also, for the nondisordered case (i.e., in the limit p = 1), one would expect to find $d_{\rm RW} = 2$, characterizing the effective "fractal" scaling dimension of unhindered RWs in any dimension.

A different approach considers "energetic" disorder associated with sites or bonds [17,18], representing random barriers that define site-to-site transition probabilities governed by Boltzmann statistics. Therein, a model inherent temperatureparameter T controls the ability of an individual walker to overcome energy barriers. At low temperatures, a RW is likely to get trapped at local minima in the energy landscape. Albeit this trapping effect is only temporary, the RWs generally display a rather limited mobility. With increasing temperature the RWs gain mobility and behave similar to ordinary RWs in the limit $T \to \infty$. Regarding this energetic disorder one might further distinguish static (i.e., quenched) disorder, which remains unchanged in time, and dynamic disorder, which is renewed after each discrete time step; see Ref. [18]. In case of static disorder and after a temperature-dependent crossover time scale τ_c , the MSD exhibits the same asymptotic linear scaling as for unhindered RWs. Supported by analytical arguments from classical chain-reaction theory, the crossover time τ_c could be attributed to an effective activation energy that must be overcome in order to escape from an initial local minimum in the energy landscape. Intuitively, the lower the temperature, the larger the crossover time tends to be. In the early time regime, i.e., at times smaller than τ_c , the scaling of the mean-square displacement is subdiffusive. In case of dynamic disorder, no crossover time is found and the same scaling as for the unhindered RW is observed for all temperatures. The renewal of the disorder after each time step creates an averaged environment in which the temperature T directly reflects the transport efficiency in the medium, resulting only in a temperature-dependent diffusion constant D(T).

A further extension that is heavily studied in the literature consists in emphasizing a particular lattice direction along

which the RW effectively drifts. Also for this modification there are different possible implementations, differing in the particular way in which the transition probabilities are altered, see Refs. [3,32–37]. In this regard we only briefly recap one of those studies, which, for the first time, reported on a direct observation of a sharp drift-to-no-drift transition [35]. Therein the influence of a directional bias field on the diffusive behavior of random walks on 3D lattice graphs with percolation-like site disorder at p = 0.5, i.e., above the critical point $p_c \approx 0.3116$ of the 3D setup, was addressed. Initially, for a given realization of the disorder, noninteracting RWs were started on randomly chosen occupied sites and individual walkers were only allowed to advance to adjacent occupied sites. A bias parameter $0 \leq B \leq 1$ is used to control the drift of the RWs along the x direction: The probability to advance along the positive x direction is set to B, and the probability to advance to any of the six neighbors is set to 1 - B. Once such a step is proposed, it is only executed if the target site also belongs to the cluster. Hence, there is the possibility that a walker remains at its current position (speaking in terms of the popular phrases mentioned in the introduction, it would be more precise to call this the "blind biased ant in the labyrinth" since the decision on the direction of the proposed move is made irrespective of whether the target-site can actually be reached). For such noninteracting walkers and prior to the numerical study reported in Ref. [35], a drift-to-no-drift transition at a finite critical value B_c was theoretically hypothesized [38,39]. For weak bias B, the walkers exhibit a diffusive motion and their velocity v(t)assumes a constant value which tends to increase with B. For strong bias, diffusion is slowed down since the walkers tend to wind up in "dangling ends" of the clusters on which they reside. With increasing bias they have an increasingly hard time to escape from such dangling ends. Neglecting initial transients, three different regimes were observed: (i) $B < B_c$, where v(t)tends to a constant; (ii) $B = B_c$, where $v(t) \sim 1/\log(t)$; and (iii) $B > B_c$, where $v(t) \sim t^{-x(B)}$ [and x(B) > 0]. The critical bias strength at which the drift-to-no-drift transition occurs was estimated as $B_c \approx 0.53$. The behavior of the average velocity for a large range of B values, covering all three scaling regimes, could further be summarized by an extrapolated scaling form. Similar studies were, e.g., also carried out to construct a full phase diagram for the drift-to-no-drift transition in the B-pplane for the 2D square lattice [36]. Regarding the MSD, the RW component of the dynamics leads to diffusion $\langle R(t)^2 \rangle \propto t$ and the bias induces a drift $\langle R(t)^2 \rangle \propto t^2$. At this point note that under drift, the asymptotic effective dimension $d_{RW} = 1$ might be expected. At short (long) times diffusion (drift) is dominant. Under a strong bias, the dangling ends act as traps that inhibit drift by characteristic "trapping times" (depending on the precise cluster structure). Consequently, the MSD is governed by effective exponents $\alpha_{\rm eff}(p)$ that approach their asymptotic value rather slowly [3,26,36,39]. Further, biased diffusion for networks without underlying regular spatial structure (i.e., regular random graphs, Erdős-Renyi random graphs and scale-free networks) was recently investigated by means of numerical simulations and analytic calculations [40]. Therein, the bias was designed in a way that it guides the walker to a specified target node along the respective shortest path. Using a model parameter, the strength of the bias, ranging from unbiased (yielding an ordinary random-walk dynamics) to fully biased (effectively confining the walker to the shortest path), could be tuned. Among other things it was shown that for regular random graphs and Erdős-Renyi random graphs, the scaling of the mean first passage time as function of the system size changes from a power law to a logarithmic behavior at a characteristic value of the bias parameter.

In addition to simple RWs, self-avoiding walks (SAWs), conveniently used as models for randomly bent polymers that exhibit an "excluded volume" effect [41], have been studied for diluted systems exhibiting percolation disorder [42,43]. Therein, a basic observable is the mean-square radius $\langle R^2 \rangle$ of N-step SAWs which scales according to $R_N \propto N^{2\nu_{\text{SAW}}}$, where the exponent v_{SAW} can be seen as an inverse of a SAW characteristic scaling exponent $d_{\text{SAW}} = v_{\text{SAW}}^{-1}$ that is universal in each dimension d. Depending on the fraction p of occupied sites on the underlying lattice and the set of configurations over which averages are computed, different SAW scaling regimes can be found [43]. For the case of no disorder (equivalent to p = 1), results consistent with $v_{\text{SAW}} = 0.75$ ($d_{\text{SAW}} = 1.333$) are obtained in two dimensions, see, e.g., Ref. [44]. A different study, aimed at clarifying the geometric and energetic scaling behavior of minimum energy (ME) SAWs on lattices with random site energies [45], reports $v_{\text{ME-SAW}} = 0.80(2)$ $[d_{\text{ME-SAW}} = 1.25(3)]$ for the 2D case (obtained via exact enumeration methods). Thus, under the influence of quenched randomness, minimum energy SAWs tend to expand. Note that within error bars this value is consistent with the fractal dimension $d_f = 1.268(1)$ of paths in the NWP problem at the particular point $\rho = 0.1032(5)$ (also in 3D the estimates of the respective exponents stack up well: $d_{\text{ME-SAW}} = 1.41(6)$ as compared to $d_{\rm f} = 1.459(3)$ [6]).

A further variant of a simple RW, wherein a loop is erased as soon as it is formed, is referred to as loop-erased random walk (LERW) [19]. A LERW might be interpreted as a simplified version of a SAW which has some correspondence to spanning trees and, in 2D, exhibits a scaling dimension of $d_{\text{LERW}} = 1.25$ [19,46].

IV. RESULTS

As pointed out in the introduction, minimum weight paths on disordered lattices were already studied in terms of the NWP model, which, for a given realization of the disorder, can be solved numerically exact by using quite involved optimization algorithms. Instead of using such a complex algorithmic approach that takes into account global information, we here address the question how well local algorithms perform. In particular, we investigate how well the transition point $\rho_c = 0.0869(2)$ for the 2D square lattice with bimodal disorder, resulting from static simulations considering the NWP model, can be resolved using simple dynamic algorithms. We study local algorithms because the movement of particles in physical systems, or the decisions of agents in populations, typically depend on local information only. Hence, such an approach reflects real situations better. Therefore we consider the dynamic biased and greedy RW dynamics discussed previously in Sec. II A and the ant colony optimization heuristic described in Sec. II B. The results are reported in Secs. IV A and IV B, respectively.



FIG. 3. Results for type A BGRWs where a walker interacts with the lattice graph according to rule (i) described in the text. The figure shows the probability $P(\rho)$ of finding a lattice walk with a negative weight as function of the disorder parameter ρ . (a) $P(\rho)$ for different bias strengths B (listed in the key). For a given value of B the respective 4-tuple of curves correspond to the system sizes L = 100,200,400,800 (the larger the system size the steeper the curve). The common crossing point indicates a critical value $\rho_c(B)$ at which, in the thermodynamic limit, lattice walks with negative energy might be found. (b) the main plot shows a scaling plot for the particular bias strength B = 0.5, therein the abscissa was rescaled according to $(\rho - \rho_c)L^{1/\nu}$ with $\rho_c = 0.2777(6)$ and $\nu = 1.97(4)$. The inset illustrates (for fixed L = 800) the increasing steepness of data curves for a decreasing bias parameter B. I.e., the steepest curve is found for B = 0.05, while at B = 1.0 the curve increases most gradual.

A. Biased and greedy random walks

First, we consider type A BGRWs where the walker interacts with the environment via rule (i), i.e., while traversing the graph it replaces positive and negative edge weights by edge weights equal to zero. In Fig. 3 we show the probability that, after termination, a lattice walk with negative weight is found. Simulations are carried out for systems of size L = 100,200,400,800 where data points represent an average over 2×10^4 realizations of the disorder.

To facilitate intuition, for a bias strength B = 1 the greediness of the walker has no effect. It starts at the outer

node on the, say, left lattice boundary and performs a straight line (i.e., L + 1-step) walk to arrive at the outer node on the right lattice boundary (thus, the lattice walk is characterized by a scaling dimension $d_s = 1$). Since the disorder is uncorrelated, the walker basically sums up L - 1 uncorrelated random edge weights drawn from the disorder distribution P_{bi} , see Eq. (1). At the value $\rho = 0.5$ of the disorder parameter and for large values of L (on average) half of the edge weights will be negative. Hence, above $\rho_c(B=1) = 0.5$, the value of $P(\rho)$ will tend to 1 as $L \to \infty$. At $\rho = 0.5$ this is equivalent to a symmetric 1D RW, where, starting at the origin, one allows the walker to perform L-1 steps, afterwards asking for the probability that the walker is located, say, left of the origin. Also note that in this limit, the effective 1D problem statement is trivially equivalent to the 1D NWP problem. To further understand the scaling behavior of the data curves for the case B = 1 shown in Fig. 3, one might follow a real-space renormalization approach [47-49]. The basic idea is to replace a subsystem of L edges (characterized by disorder parameter ρ , which is, generally spoken, the probability to have a negative weight) by one edge characterized via an effective disorder parameter $\rho' = P(\rho)$. In this regard, note that the only condition to yield a "successful" negative-weighted L-step path is that the number of negative edge weights along the lattice walk needs to exceed the number of positive edge weights by at least 1. In this question, the probability to find a negative-weighted path for systems of odd L for a given value of ρ is simply a sum of binomial terms,

$$P(\rho) = \sum_{k=L_0}^{L} {\binom{L}{k}} \rho^k (1-\rho)^{L-k},$$
 (3)

where $L_0 = (L + 1)/2$. The fixed points ρ^* for $P(\rho)$, which charazterize phases and phase transitions in the renormalization, are determined by $\rho^* = P(\rho^*)$. Intuitively, a stability analysis of the three fixed points $\rho = 0,0.5,1$ reveals that the only unstable fixed point is the previously discussed value $\rho_c = 0.5$ (at B = 1). Linearizing $P(\rho)$ about this critical point yields the scaling power ν of the observable Eq. (3) via

$$\nu = \ln(L) / \ln(dP(\rho)/d\rho), \tag{4}$$

Albeit the expression for ν converges rather slowly, e.g., it assumes the values $\nu(L = 10) = 2.56$, $\nu(L = 100) = 2.22$, and $\nu(L = 1000) = 2.14$, it attains the limiting value $\nu = 2$ as $L \rightarrow \infty$ [obtained through a series expansion of the derivative of Eq. (3) in powers of L]. This slow convergence is presumably caused by the additive character of edge weights: The total sum of the edges must be negative, while subsystems of arbitrary size are free to have a positive weight. Consequently, if the data curves for $P(\rho)$ obtained in the simulations exhibit the property of scaling, it should thus not come as a surprise if the scaling behavior is governed by the scaling parameter $\nu = 2$. Indeed, we find that the data curves show a smooth data collapse if rescaled according to

$$P(\rho) = f[(\rho - \rho_{\rm c})L^{1/\nu}],$$
 (5)

where for B = 1 the parameter values $\rho_c = 0.5$ and $\nu = 2$ yield a good data collapse (not shown). This scaling behavior still holds for values of B < 1, where $\rho_c < 0.5$. For example, in Fig. 3 the scaled data for B = 0.5 is illustrated, where

TABLE I. Comparison of the critical values ρ_c above which the path weight turns negative and the critical exponents ν for different algorithms considered in the presented study. Here, the exemplary results for the type A BGRWs were obtained using rule (i) at the bias strength B = 0.5. The results for the ACO heuristic are derived in Sec. IV B. The results for the static NWP algorithm, which determines a global optimal path for a given realization of the disorder, are taken from Ref. [4].

Algorithm type	$ ho_{\rm c}$ (bi)	ν	
BGRW-A	0.278(1)	1.97(4)	
ACO	0.14(1)	2.0(3)	
NWP	0.0869(2)	1.47(6)	

 $\rho_c = 0.2777(6)$ and $\nu = 1.97(4)$. For comparison, these results are listed in Table I along with the corresponding values obtained for the ACO heuristic, derived later in Sec. IV B, and those for the static NWP algorithm, which determines a global optimal path for a given realization of the disorder [4].

For decreasing bias strength B < 1 the greediness of the BGRW has an increasing influence on the dynamics of the walker. Hence, at a given value of ρ the average path weight after a fixed number of steps is smaller at smaller bias strengths. Consequently, the probability to find a negative-weighted path principally increases. Albeit true for intermediate bias strengths $B = 0.2 \dots 1$, this does not hold for comparatively small values B < 0.2. As evident from Fig. 4, the critical point ρ_c above which negative paths appear first as $L \to \infty$ tends to increase for a decreasing bias strength < 0.2, that is, one can observe a minimum value of $\rho_{\rm c}=0.199(4)$ at a characteristic value $B_c \approx 0.18(1)$ [for type A BGRWs considering an interaction with the environment according to rule (i)]; see Table II. A further extremal parameter set for type A BGRWs is the case of vanishing bias strength B = 0 at $\rho > 0.5$. Since $\rho_{c,bond} = 0.5$ indicates the onset of a system spanning (percolating) cluster of sites joined by edges of weight -1, for $\rho > \rho_{c,bond}$ there is a nonzero probability that the walker crosses the lattice without accumulating a single positive edge weight.

TABLE II. Values B_c of the bias strength for which the smallest ρ_c is attained for all types of BGRWs considered. The type of interaction with the graph is specified by rule (i) [rule (ii)] in the case where positive and negative [only negative] edge weights are modified during the walk. The type of disorder distribution is indicated by "bi" [bimodal distribution; see Eq. (1)] and "sc" [semicontinuous; see Eq. (2)]. In case of type B and D walks, the values of ρ_c where estimated from a fit of the data to a polynom of order three. We here list the figures up to the third decimal, only (the estimated fit errors where notoriously small; on the order of 10^{-6}).

BGRW type	$B_{\rm c}$ (bi)	$\rho_{\rm c}$ (bi)	$B_{\rm c}~({\rm sc})$	$\rho_{\rm c}~({\rm sc})$
A [rule (ii)]	0.30(1)	0.344(4)	0.27(1)	0.464(5)
A [rule (i)]	0.18(1)	0.199(4)	0.12(1)	0.280(3)
В	0.0	0.258	0.0	0.465
C [rule (ii)]	0.23(1)	0.300(4)	0.20(1)	0.401(6)
C [rule (i)]	0.18(1)	0.174(6)	0.12(1)	0.241(4)
D	0.0	0.187	0.0	0.286



FIG. 4. The critical threshold ρ_c as a function of the bias parameter *B* for the different types of BGRWs. (a) Bimodal disorder drawn from Eq. (1). The continuous line represents the value of $\rho_c = 0.0869(2)$ obtained using the static simulations in the context of the NWP problem. (b) Semicontinuous disorder drawn from Eq. (2).

Regarding BGRW types A through D, Fig. 4 shows the critical points ρ_c as function of bias strength B for both the bimodal disorder distribution Eq. (1) [see Fig. 4(a)] and the semicontinuous disorder Eq. (2) [see Fig. 4(b)]. As evident from the figure, BGRWs of type A and C show a minimum at critical disorder parameter value ρ_c at bias strengths in the range $B = 0 \dots 0.3$. In contrast to this, BGRWs of type B and D appear to reach a minimum only when the bias strength is zero. This observation holds for both bimodal and semicontinuous disorder, although for the different disorder types the minima are located at different values ρ_c and B_c . For all the types of BGRWs considered, Table II lists the lowest value of ρ_c , attained at the corresponding bias strength B_c . This can be thought of as a measure of efficiently for the different walk dynamics. The lower the dynamical threshold ρ_c , the more successful a walker upon traversing a graph aimed at finding a path of total negative weight (under the respective dynamics).

Another quantity that allows us to characterize the different dynamics is the average number of steps taken by the different BGRW types until the walk terminates; see Fig. 5. This quantity can further be used to determine the effective scaling dimension of the walks: Given that a walk of N steps spans

an linear distance L, its associated scaling dimension is defined via $N \propto L^{d_s}$ (Note that the concept of the fractal dimensionality of a random walk in the presence of anomalous diffusion is also used in the literature, see Refs. [24,31]. As pointed out in Sec. III, one situation in which one encounters anomalous diffusion is, e.g., for RWs on percolating clusters near the percolation thresholds and for short walking times [31]).

Note that for the resulting paths, due to the possibility of loop erasure for some algorithms, the scaling dimension might differ. On a general basis, as soon as there is a bias that causes an effective drift of a walker (as in the BGRWs considered here), one can expect that on large enough lattices one trivially finds $d_s = 1$ (a RW for which the mean-squared displacement scales $\propto N^2$, leading to an equivalent of $d_s = 1$, is usually said to exhibit a ballistic movement). Note that this differs when the environment itself is fractal, as for biased RWs on percolation clusters discussed in Sec. III. Here, for completeness and to illustrate the precise dependence of the dynamics on the disorder parameter ρ and different bias strengths *B*, the rescaled number of steps N/L taken during the lattice walks under the four BGRW dynamics are shown in Fig. 5 for the bimodal disorder distribution.

Considering type A BGRWs, see Figs. 5(a) and 5(b), it can be seen that the number of steps generally decreases when the density of negative edges is increased. This can be easily explained by taking into account the greedy behavior of the walker. In case that the walker only consumes the negative edges [replacing them with a standard positive value of 1; see rule (ii) discussed abovel, this behavior appears since, after each step taken, the probability of returning to the previous node is smaller than that of proceeding forward towards other directions that might have negative edges. In the case where a type A BGRW explores a graph with a low value of the disorder parameter and alters both, positive and negative edge weights (replacing them by edges with weight zero), an effective "repulsive effect" takes place because of the comparatively large probability that it might encounter excess edges with positive weights. Hence, the walker is likely to return to its previous location by traversing the edge with weight zero created while performing the last step. This repulsive effect is emphasized especially for a bias strength B < 0.5. As can be seen in Fig. 5(b), this repulsive effect leads to an increase in the number of steps taken by more than one order of magnitude in the low bias regime, compared to the case when only the negative bonds are altered. Also note that for different system sizes L the data curves for the scaled number of steps N/L fall on top of each other. Hence, the expected scaling dimension $d_{\rm s} = 1$ can be verified from the figures.

Regarding BGRWs of type B, illustrated in Fig. 5(c), the number of steps taken in the extreme cases $\rho = 0$ (positive edge weights only) and $\rho = 1$ (negative edge weights only), where there is no disorder at all, is the same, as one would expect. Regarding the bimodal disorder distribution, a minimum in the walk length is found when the disorder parameter assumes the value $\rho \approx 0.25$. This is caused by the fact that each node is, on average, the end node of one edge with negative edge weight that acts like a local trap for the walker. Only under the influence of the bias the walker can escape from such a trap. Until then, the walker has possibly



FIG. 5. Results for the average number of steps taken by the different BGRW types until the walk terminates. The panels (a) through (f) show the average number of steps as function of the disorder parameter ρ for the bimodal disorder distribution for different bias strengths *B* (see right border of figures, respectively) and for different system sizes *L*. Note that, because the number of steps is normalized by the system size *L*, the system-size dependence is very small, the main differences are through the bias parameter *B*.

traversed the edge with negative weight multiple times. After the loop erasure process (discussed for type B BGRWs in Sec. II A), this results in an almost straight-line walk. Increasing ρ above $\rho = 0.25$ also increases the likelihood of a node having, on average, more than one incident edge with negative edge weight, thus preventing the walker from getting stuck (however, note that for the 2D setup considered here, only above $\rho = 0.5$ there is a system spanning cluster of negative edge weights). Although not shown here, in the case of a semicontinuous disorder distribution, the number of steps shows a similar behavior up to $\rho = 0.25$ but instead of increasing afterwards it has a steady decrease until it reaches $\rho = 1$ density which is caused by the lack of degeneracy in the bond weights. A close similarity between BGRWs of type A and C can be seen upon comparison of the respective subfigures in Fig. 5. The most evident difference between the two dynamics is that type C BGRWs do not reach the minimum number of steps when the graph is fully equipped with negative edge weights. This effect appears due to the additional optimization overhead that takes advantage of the larger number of negative edge weights at large values of ρ by increasing the length of the random walk.

For type D BGRWs, shown in Fig. 5(f), a behavior that strikingly differs from the other three BGRW dynamics can be observed. The respective lattice walks exhibit a minimum in the number of steps when no negative edges are present in the graph and attain a maximum when the graph is fully equipped

with negative edge weights. This is a telltale sign of an increase in the efficiency of traversing edges with a negative weight, while avoiding positive edge weights. Shorter lattice walks are performed at small values of ρ and longer paths, possibly with many detours, are found at larger values of ρ . The latter characteristic for the optimal paths in a graph predominantly equipped with negative edge weights.

It is interesting to note that the more efficient a random-walk algorithm is at finding negative-weighted paths, the higher the number of steps it will take at a low density of negative bonds if it alters both negative and positive bonds, corresponding to rule (i) discussed above. As we observe, this effect is especially pronounced in the case of BGRWs of type D. The respective algorithm could not even be used efficiently for $\rho < 0.5$ (no results for this case are shown).

As pointed out above, the expected scaling dimension for the BGRWs is $d_s = 1$. The number of steps taken during the lattice walks, normalized by the system size, verifies this in the form $\langle N \rangle / L \approx \text{const}$ for given values of ρ and B. Discrepancies in the curve overlap, corresponding to deviations from the expected scaling behavior, are accentuated at small values of B, because the size of the walk gets closer to a *B*-dependent crossover length L_c (similar to the crossover time τ_c discussed in Sec. III). For walks much shorter than this crossover length the BGRWs exhibit a scaling dimension larger than 1. Only for walks longer than the crossover length do they assume the asymptotic scaling dimension $d_s = 1$. The simple explanation for this effect is that on short length scales the diffusive (albeit greedy) behavior dominates the configuration of the lattice walk, while on longer length scales the drift induced by the bias is dominant. Note that the scaling dimension of the lattice walks obtained here differes markedly from the fractal scaling dimension of the paths in the 2D NWP problem. For the 2D setup of the NWP problem, a fractal scaling dimension of $d_f = 1.268(1)$ was observed at $\rho = 0.1032(5)$ [4], increasing towards $d_{\rm f} = 1.756(8)$ at $\rho = 1$ (measured for the loops-only setup) [50].

B. Ant colony optimization

In order to find values of α and β that reflect the underlying edge weights, leading to an exploration in the vicinity of the "best" (s,t) path found so far (referred to as *intensifi*cation) but still allowing for an efficient exploration of many different paths (referred to as *diversification*), we performed several "calibration" runs of the ACO algorithm. During our simulation we used M = L, $\alpha = 1/M$, and $\beta = 0.98$. Below we report the results of a finite-size scaling analysis for the probability that the ACO algorithm finds a negatively weighted optimal path for different values of the fraction ρ of negative edge weights on square lattice graphs with L = 8, ..., 20. At each value of ρ , we considered $n = 10^3$ different realizations of the edge weights drawn from a bimodal disorder distribution [Eq. (1)] in order to compute averages. The considered systems have periodic boundaries in the, say, vertical direction and open boundaries in the horizontal direction. Further, two extra "outer" nodes, i.e., the source node s and the target node t, are introduced. The source (target) node is connected to all nodes on the left (right) system boundary as shown in Fig. 1. Consequently, a (s,t) path spans the system along the



FIG. 6. Exemplary (s,t) paths found by the ACO algorithm for different fractions ρ of negative edge weights on a 2D square lattice of side length L = 16. (a) $\rho = 0.10$, (b) $\rho = 0.15$, and (c) $\rho = 0.90$. The black path indicates the best path found after the algorithm terminated and the linewidth of the other edges (gray) indicate how frequently the respective edge was visited by an agent.

horizontal direction. Depending on the fraction ρ of negative edge weights it might also wrap the system along the vertical direction. Figure 6 illustrates exemplary (s,t) paths found by the ACO algorithm for three different values of the disorder parameter on a 2D square lattice of side length L = 16. In the figure, the black path indicates the best path found after the algorithm terminated and the linewidth of the other edges (colored gray) indicate how frequently the respective edge was visited by an agent. As evident from Fig. 6(a), corresponding to $\rho = 0.10$, if there are only a few negative edge weights, the "best" (s,t) path found by the ACO algorithm is rather straight lined. As the fraction of negative edge weights increases, see Figs. 6(b) and 6(c), corresponding to $\rho = 0.15$ and $\rho = 0.90$, respectively, the paths roughen up and eventually also wind around the lattice in the vertical direction.

The characteristics of the weight of the best path found by the ACO algorithm as a function of the number of sweeps carried out by the algorithm is shown in Fig. 7(a). As can be seen in the curve corresponding to $\rho = 0.01$ and after a short exploration time (t = 0-500 sweeps), the ACO algorithm maintains an average best pathweight $\langle \omega_p^{\text{best}} \rangle \approx L$. Since the underlying edge-weight distribution is bimodal, i.e., it allows for $\omega_{ij} = \pm 1$ only, this reflects that at very small values of ρ almost all edges contained in (s,t) paths will have weight 1 and the paths tend to have length $\approx L$. The situation changes upon increasing the value of ρ . As can be seen from Fig. 7(a), already at $\rho = 0.15$ and after a certain initial exploration time (t = 0-1000), the ACO algorithm is able to identify (s,t) paths with a negative weight so $\langle \omega_p^{\text{best}} \rangle < 0$.

Finally, Fig. 7(b) illustrates the scaling behavior of the probability $P(\rho)$, that the ACO algorithm returns a negativeweighted (s,t) path as a function of the fraction ρ of negative edge weights. As can be seen from the inset of Fig. 7(b), the data curves for the different system sizes cross at $\rho_c \approx$ 0.13. Below (above) ρ_c and for increasing system size, the probability that the ACO algorithm returns a path weight $\omega_{\rm p}$ < 0 tends to zero (1). As can be seen from the main plot of Fig. 7(b), $P(\rho)$ can be rescaled using Eq. (5), similarly to the paths found using the BGRW dynamics discussed previously. Considering the system sizes $L \ge 10$, a best data collapse (attained in the range $(\rho - \rho_c)L^{1/\nu} \in [-0.3, 0.3]$ on the rescaled ρ axis) results in the estimates $\rho_c = 0.14(1)$ and $\nu = 2.0(3)$ with a quality S = 1.84 of the data collapse [51] (the numerical value of S measures the mean-square distance of the data points to the master curve, described by the scaling



FIG. 7. Results for the ACO heuristic on 2D square lattice graphs for the case of a bimodal disorder distribution. (a) Best path weight found so far as function of the number of sweeps carried out by the ACO algorithm, averaged over different realizations of the edge weight disorder at L = 16 for three different values of ρ . (b) Probability $P(\rho)$ that the ACO algorithm returns a (s,t) path with negative weight as a function of the fraction ρ of negative edge weights. The main plot shows the scaled data (see text for more details) and the inset illustrates the unscaled data in the vicinity of the critical point $\rho_c = 0.14(1)$.

function, in units of the standard error [52]). Performing the analysis for different intervals on the rescaled ρ axis led to various estimates in the range $\rho_c = 0.13 \dots 0.14$. Further estimates of ρ_c and ν can be obtained from the scaling behavior of the variance var(P). It assumes a peak at an effective, L-dependent critical point ρ_L^{eff} at which $P(\rho_L^{\text{eff}}) = 0.5$ (not shown). We find that ρ_L^{eff} seems to saturate at a value consistent with ρ_c . In particular, we find $\rho_{20}^{\text{eff}} = 0.137(1)$ [for the smaller system size L = 10 we find $\rho_{10}^{\text{eff}} = 0.137(2)$], obtained by fitting a Gauss-shaped curve to the data of var(P). The value of ρ_L^{eff} specifies the location of the peak of the fitting curve and the error-bar is the respective fitting error. Further, the width σ_L of the Gauss-shaped fitting function is consistent with $\sigma_L = a L^{-1/\nu}$, where we found a = 0.098(6) and $\nu = 2.1(1)$ considering $L = 8 \dots 20$. Note that the value of the exponent ν found here is also in agreement with the one found for the BGRWs in the preceeding subsection.

V. CONCLUSIONS

In the presented article we have investigated the principle characteristics of BGRWs on 2D lattices with quenched disorder on the lattice edges. Four different types of BGRWs and an algorithm based on the ACO heuristic were considered. Regarding the BGRWs, the precise configurations of the lattice walks constructed during the numerical simulations were influenced by two parameters: a disorder parameter ρ that controls the amount of negative edge weights on the lattice and a bias strength B that governs the drift of the walkers along a certain lattice direction. Focus of the presented study was the probability that, after termination, a lattice walk exhibits a negative weight as function of ρ and B. All four types of dynamics exhibit a phase transition with a characteristic disorder value $\rho_{\rm c}$ above which lattice walks with a negative weight appear first. Further, the data curves for the probability that a walk with negative energy is found as function of ρ could be rescaled well according to a simple scaling form. The respective scaling parameters in the extremal case of maximal bias B = 1, where the dynamics is effectively one dimensional, could be estimated from a simple renormalization argument. A lower bound on the value of ρ_c can be obtained from the NWP problem in 2D. Regarding those "static" simulations considering bimodal disorder on 2D square lattices, negative-weighted paths appear first above the threshold $\rho_c = 0.0869(2)$ with critical exponent v = 1.5(1). Considering the BGRW dynamics only, the type C BGRW [implementing rule (i) for the walker-edge weight interaction discussed in Sec. II A] yields the closest estimate, reading $\rho_c = 0.174(6)$ at bias strength $B_c = 0.18(1)$, followed closely by type D lattice walks with $\rho_c = 0.18(1)$ at vanishing bias; see Table II. Correspondingly, algorithm A (i) needs the largest number of steps to yield the best results. Interestingly, algorithm D is considerably faster, also compared to almost all variants. Nevertheless, as visible from Fig. 5, all algorithms, due to the bias, have an asymptotic running time O(L). Thus, the greedy random walks perform faster than the exact algorithms, which in practice perform in time $O(L^2)$ [4]. This explains why the greedy random walks are only able to achieve considerably larger values for the thresholds $\rho_{\rm c}$ compared to the exact approach.

The ACO-based algorithm outperforms these estimates, yielding $\rho_c = 0.14(1)$. Note that, in comparison to the type A and B BGRWs, which perform slightly worse, the dynamic rules governing the type C and D lattice walks and the ACO heuristic are rather involved. Also note that, albeit all the dynamic rules implement some degree of "local" optimization, they were not able to reach the threshold value obtained from the static simulations in terms of the NWP problem. Also the value $\nu = 2$ of the critical exponent, for all types of BGRW and ACO dynamics, and the fractal behavior of the walks, hence of the paths, differs from the static negative weight paths. Thus the static and dynamic behaviors of the problem differ, as for many glassy systems.

Whether one can reach this static threshold and find the same critical exponents by means of simple dynamics that implement local optimization remains an intriguing open problem. This could be possible, in principle, because the corresponding optimization problem can be solved in polynomial time. At least we could not succeed in doing so via the biased greedy random walks considered in the presented study. BIASED AND GREEDY RANDOM WALKS ON TWO- ...

Another interesting problem is whether a collective algorithm, which explores, e.g., the dynamics of an ensembles of paths connected to a heat bath at temperature T, with the path weight being the energy, exhibits phase transitions as function of the disorder parameter ρ as well. These transitions could be in the limit $T \rightarrow 0$ similar to the phase transitions found previously for the exact algorithms. For lower temperatures, maybe below a critical dynamical threshold, these transitions could be more similar to the transitions investigated in this work. Related studies are currently being performed.

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