Overrelaxation and Monte Carlo simulation

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(Received 6 March 1987)

I study a simple variation of the algorithm of Metropolis *et al.* for simulating statistical systems. The trial changes in any given variable are taken from a region of phase space far from the old value but involving only small changes in energy. This results in correlation times which are short compared to the usual applications of the algorithm of Metropolis *et al.* Tests with SU(2) and SU(3) lattice gauge theories indicate substantial possible savings in computation time relative to standard approaches.

Monte Carlo simulation has become the primary tool for the study of nonperturbative phenomena in both quantum field theory and statistical systems. Indeed, this approach has resulted in theoretical physicists becoming rather avid users of computer time. Despite severe difficulties with including anticommuting fermionic fields in such simulations, the results have been impressive. The technique has given rise to useful quantitative information in statistical mechanics about critical phenomena and in particle physics about the confining interquark potential, glueball spectra, and the transition to a quarkgluon plasma.

These simulations have generally used a version of the old algorithm of Metropolis *et al.*¹ A few years ago $Adler^2$ suggested extending overrelaxation techniques for solving linear equations to Monte Carlo simulations. His approach, which was further investigated by Whitmer,³ was limited to theories where the action or energy function appearing in the Boltzmann weight is quadratic in any single individual variable. This restriction to multiquadratic actions has prevented application to the more complicated forms appearing in lattice gauge theory.

Motivated by overrelaxation ideas, I present in this paper a simple variation of the algorithm of Metropolis *et al.* This approach is applicable to any theory where the variables are elements of a group larger than Z_2 . The basic idea is to pick a trial change for a given variable in a region of phase space which is as far as possible from the old value while not paying a severe energy penalty. This is done by approximately locating the locus of minimum energy for the variable under consideration, and then selecting the trial element on the "opposite" side of this value.

I begin with a review of the standard algorithm of Metropolis *et al.* Consider updating some variable g whose value lies in a group G. While holding other variables fixed, write the Boltzmann weight for this variable as

$$P_{\rm eq}(g) \propto \exp[-\beta H(g)] . \tag{1}$$

In this paper I will always consider variables as elements of a group. This is quite natural in the standard formulation of lattice gauge theory. This use of group elements is for notational convenience and is not a practical restriction; real variables can always be thought of as belonging to the group of numbers under addition.

The procedure of Metropolis *et al.* considers a trial new value g' to replace g. This is selected with some probability distribution $P_{T,g}(g')$. Here the subscript g on P is to remind us that the trial change can in general depend on the old element g. The trial element g' is then accepted with a conditional probability

$$P_{A} = \min \left[1, \frac{P_{T,g'}(g)}{P_{T,g}(g')} \frac{\exp[-\beta H(g')]}{\exp[-\beta H(g)]} \right].$$
(2)

If this conditional probability is not met, then the change is rejected and the old value of g is kept. This construction automatically satisfies the detailed balance condition

$$P(g \rightarrow g')e^{-\beta H(g)} = P(g' \rightarrow g)e^{-\beta H(g')}, \qquad (3)$$

where $P(g \rightarrow g')$ is the overall probability of taking element g to g'. When $g' \neq g$, $P(g \rightarrow g') = P_{T,g}(g')P_A$. Equation (3) is sufficient to ensure that the equilibrium distribution of Eq. (1) is invariant under the algorithm. That invariance plus ergodicity are necessary and sufficient conditions for a Monte Carlo algorithm to bring any ensemble towards equilibrium.

As usually implemented, g' is found by multiplying g by a group element h which is chosen with a probability distribution peaked around the identity and with equal probability for h and h^{-1} . For example, this can be done by choosing h from a table which contains the inverse of each of its elements. This particular method of choosing the trial change has the symmetry property

$$P_{T,g}(g') = P_{T,g'}(g)$$
 (4)

This is convenient because the ratio of P_T factors in Eq. (2) is then unity and can be ignored. The conventional approach contains an implicit parameter which represents the average distance the element h lies from the identity. If this distance is too large, the trial energies will likely be large and the changes will rarely be accepted, while if h always lies too close to unity, the changes will usually be accepted, but their small size will make the exploration of phase space rather slow. Lore is for a compromise with

an acceptance rate of order 50%.

The "heat bath" algorithm⁴ is a special case of this generalized approach of Metropolis *et al.*, wherein the trial element is taken randomly from the entire group manifold but with a weighting proportional to the Boltzmann factor of Eq. (1). In this case the factors appearing in the acceptance criterion of Eq. (3) cancel and the trial element is always accepted. This method is equivalent to the limit of taking a large number of repetitions, or "hits," of any ergodic algorithm to a single variable.

The algorithm proposed in this paper represents yet another way to choose the trial change g'. Suppose I have some simple way to find a group element g_0 which approximately minimizes the energy H(g). Suppose further that g_0 is obtained with no direct use of the element g; that is, g_0 only has a dependence on the remaining lattice variables. The essence of this paper is to propose selecting the trial element for a Metropolis *et al.* updating to lie on the "opposite side" of this element g_0 from the old value g. In particular I consider the trial element

$$g' = g_0 g^{-1} g_0 . (5)$$

Note that this construction also satisfies the symmetry relation of Eq. (4). Thus, just as in the usual application of the algorithm off Metropolis *et al.*, the acceptance or rejection of this element follows solely from the comparison of a random number with the exponential of the resulting energy change.

To be more specific about the selection of g_0 , consider the case where H(g) takes the form

$$H(g) = E_0 - \operatorname{Re} \operatorname{Tr}(gM) , \qquad (6)$$

where E_0 includes contributions to the energy which do not depend on g. In the case of lattice gauge theory with the Wilson action, M is a matrix which represents the sum of the three sided "staples" surrounding the link to be updated. For a nearest-neighbor spin model with the spins being matrices from the group G, M is the sum of the group elements on sites neighboring g. With this form for the energy, a natural choice for g_0 is to use the inverse of a group element obtained by projecting M onto G. As the correctness of the algorithm does not depend on the prescription for this projection onto the group, it is advantageous to keep the procedure as computationally simple as possible. In my tests with SU(N) lattice gauge theory, I use a Gram-Schmidt orthogonalization process on the matrix rows and then divide the last row by the determinant of the resulting matrix.

Note that the algorithm cannot be used for variables in the group Z_2 . Indeed, in this case Eq. (5) always gives g'=g and no changes are proposed. This will also be the case for any group when the system is totally ordered, with g and g_0 being the identity. Thus to start the algorithm, some disorder must be initially present. This does not mean that the algorithm will not be useful at low temperatures; rather, the size of proposed changes automatically decreases for cool systems.

In some cases, most notably with the groups SU(2) and U(1), the matrix M is always proportional to a group element. Projecting onto this particular element results in g' having exactly the same energy as g. In this case the

prescription of Metropolis *et al.* will always accept the change, and the algorithm is deterministic and microcanonical. This causes two minor complications. First, the total energy of the system is fixed and thus will not relax to any value other than where it is initially set. Second, the algorithm is actually independent of the temperature β^{-1} . Indeed, as with other microcanonical algorithms, the temperature should be measured during the simulation with some sort of thermometer, such as an average kinetic energy,⁵ using an auxiliary variable with simple dynamics,⁶ or from a dynamical equation involving both the temperature and measurable correlation functions.⁷

This issue can be avoided if desired by putting a small amount of randomness into g_0 . For example, g_0 could be the product of a deterministic estimate of the element minimizing H(g) with a random element h chosen near the identity. If h has a small probability of lying anywhere in the group, this would also eliminate possible worries about ergodicity. Nevertheless, I have done limited studies which suggest that correlation times tend to increase with additional noise in g_0 .

I now turn to some tests of the algorithm. Consider standard four-dimensional SU(N) lattice gauge theory with the Wilson action

$$H = \sum_{p} \frac{1}{N} \operatorname{Re} \operatorname{Tr} g_{p} \quad , \tag{7}$$

where the sum runs over all elementary squares or plaquettes p on a simple hypercubic lattice. The fundamental variables are group elements associated with the links of the lattice and the quantity g_p represents the product of such elements around the sides of the respective plaquette p.

To simplify vectorization, all simulations presented here used skew-periodic boundary conditions on a $7 \times 7 \times 7 \times 6$ lattice. The links in any given direction were updated in a checkerboard style, with all those emanating in a positive direction from odd sites being updated before those from even sites. The lattices were initially equilibrated with 100 iterations of a heat-bath algorithm for SU(2) and an optimized 10-hit algorithm of Metropolis *et al.* for SU(3). Where error bars are shown, they were obtained by repeating the respective experiments 5 or 20 times.

A simple measure of the correlation between two lattices U and V with corresponding links $(g_l)_U$ and $(g_l)_V$ is

$$C(U,V) = \frac{1}{n_l N} \sum_{l} \operatorname{Re} \operatorname{Tr}[(g_l^{-1})_U(g_l)_V], \qquad (8)$$

where the sum is over all links l and n_l is the total number of links. This quantity is unity where U and V are the same and vanishes for uncorrelated lattices. I will be considering lattice V obtained from U through a few applications of various Monte Carlo algorithms. The speed with which this correlation drops to zero is then indicative of the efficiency of the algorithm.

Figure 1 shows the falloff of this interlattice correlation as a function of the number of Monte Carlo iterations separating the lattices U and V. Here the gauge group is SU(2) and $\beta = 2.3$. The correlation is shown both for the standard heat-bath algorithm and the overrelaxation algorithm. As mentioned above, for SU(2) the latter approach



FIG. 1. The correlation between two lattices as a function of the number of simulation iterations separating them. The model is SU(2) lattice gauge theory at $\beta = 2.3$. The solid points are for the heat-bath algorithm while the open triangles are for the over-relaxation algorithm presented in the text.

is both microcanonical and deterministic. In this figure the heat-bath algorithm appears to give a simple monotonic decrease of the correlation, while the overrelaxation approach decorrelates somewhat faster in an oscillatory fashion.

Figure 2 shows the same correlation for the case of SU(3) at $\beta = 6.0$. I determine g_0^{-1} from a Gram-Schmidt



FIG. 2. The decrease with Monte Carlo iterations of the correlation between lattices with the gauge group SU(3) at β =6.0. The open triangles represent the overrelaxation algorithm. The solid points, the crosses, and the open circles represent a standard algorithm of Metropolis *et al.* with 10, 64, and 128 hits per link, respectively.

orthogonalization process on the matrix M interacting with the element being updated. Unlike for SU(2), the overrelaxation algorithm is now neither deterministic nor microcanonical. I find the acceptance rate for the trial changes is 57%.

For comparison Fig. 2 also shows the correlation falloff for a Metropolis *et al.* updating with the 10, 64, and 128 "hits" or trial changes for each element before moving to update the neighbors. Of course, as the number of such hits goes to infinity one should approach a heat-bath algorithm. These standard runs used trial elements g' selected by multiplying g with a matrix h chosen with probability $P(h) \propto \exp(k \operatorname{Re} \operatorname{Tr} h)$. I chose $k = 2\beta$ because, at least for this value of β , this empirically optimizes the correlation decrease (independent of number of hits). For this value of k and β , the acceptance per hit was 30%.

For all runs in Fig. 2 the correlation appears to be monotonically decreasing, with the overrelaxation algorithm decreasing the fastest. Indeed, noting the small change in going from 64 to 128 hits, the figure suggests that the new algorithm outperforms the heat bath. This is true even with the naive method for projecting M onto the group. Although the runs in this figure are consistent with approaching exponentials, it might be dangerous to assume that this continues. There could be hidden longtime correlations which emerge upon further running.

One potential difficulty with this measure of correlation is associated with the gauge invariance of the system. Indeed, a random gauge transformation will give a new lattice with zero expected correlation with the old one. This is in spite of the fact that all gauge-invariant quantities are identical on the two lattices. All the algorithms I will study here are based on updating single links at a time and do not take any special advantage of the gauge symmetries of the theory. Nevertheless, because of this worry, consider the correlation between gauge-invariant plaquette operators. In particular, define

$$C_p(U,V) = \frac{1}{n_p} \sum_p (W_p - \langle W \rangle)_U (W_p - \langle W \rangle)_V , \quad (9)$$

where

$$W_p = \frac{1}{N} \operatorname{Re} \operatorname{Tr} g_p \tag{10}$$

on the respective lattice, n_p is the total number of plaquettes, and the expectation value $\langle W \rangle$ is the plaquette averaged over the entire lattice. Because the plaquette has an expectation value, this is subtracted to obtain a vanishing correlation for independent lattices. One disadvantage of using this gauge-invariant measure of correlation is purely numerical; the corresponding numbers fall quite rapidly and are more difficult to measure than the quantity in Eq. (8).

In Fig. 3 I show the quantity from Eq. (9) for the SU(2) runs which gave rise to Fig. 1. Note that with this measure the behavior of the new algorithm is somewhat slower at decorrelating than the heat bath. There is also a hint of an even-odd oscillation in the convergence. This might be expected because for the new algorithm a double hit on the same group element would result in no change to the lattice. Figure 4 shows the corresponding quanti-



FIG. 3. The correlation between plaquettes vs iterations for the SU(2) lattices used in Fig. 1. The solid points represent heat-bath updating and the open triangles are for the overrelaxation algorithm.

ties for the SU(3) runs. Here the overrelaxation approach with one hit appears to be slightly worse at decorrelating plaquettes than the 10-hit standard algorithm of Metropolis *et al.*

That the new algorithm performs less well at decorrelating plaquettes than links might be expected from the fact that the SU(2) algorithm is microcanonical and the SU(3) case approximately so. Thus there is a built-in glo-



FIG. 4. The plaquette correlation for the SU(3) runs in Fig. 2. The open triangles are for the overrelaxation algorithm, while the solid points and crosses are for standard Metropolis *et al.* updating with 10 and 64 hits per link, respectively.

bal correlation between plaquettes which appears to show up in the local plaquette-plaquette correlation.

In a sense the overrelaxation following from Eq. (5) is the maximum that one would intuitively expect to be useful. It might be interesting to consider something less extreme and choose the trial element somewhere between gand g'. Indeed, the stochastic Langevin equation ⁸ can be thought of as an underrelaxed algorithm of this type wherein the trail element is chosen by multiplying the old element by a driving force towards g_0 and then introducing a noise with width selected to make the acceptance in Eq. (2) unity to lowest order in the step size.

For an intermediate case, I have also considered picking the trial element by multiplying g_0 by a random element h near the identity

$$g' = g_0 h \quad . \tag{11}$$

The elements h were selected with distribution

$$P(h) \propto \exp(k \operatorname{Re} \operatorname{Tr} h)$$
 (12)

This distribution was obtained via a separate Monte Carlo simulation. In Fig. 5 I show the acceptance rate for the



FIG. 5. (a) The acceptance rate for trial links chosen using Eq. (11) as a function of the parameter k appearing in Fq. (12). The model is SU(3) lattice gauge theory at $\beta = 6.0$. (b) The correlation between lattices separated by 10 iterations as a function of k, where the trial changes are selected as in (a).

trial changes as well as the correlation between two lattices separated by 10 sweeps as a function of the parameter k. The gauge group is SU(3) and $\beta = 6.0$. Here only one trial change was considered for each variable before moving on to the next. Note that for a narrow region of the parameter k, the decorrelation rate is quite reasonable; indeed, it is better than the standard 10-hit approach as shown in Fig. 2. This approach does not satisfy the symmetry relation of Eq. (3), so the extra ratio of trial probabilities enters nontrivially into the acceptance criterion. This is somewhat of a disadvantage because systematic errors could be introduced if the matrices h are not chosen independently with the correct distribution. At large kthe acceptance becomes low because the ratio of probabilities P_T becomes small, while at small k acceptance suffers because most changes result in large energies.

To summarize I have presented a variation of the scheme of Metropolis *et al.* for Monte Carlo simulation. The approach is motivated by overrelaxation ideas; that is, I consider trial changes which lie beyond the minimum of the energy from the old value of a variable being updated. There are two intuitive arguments in favor of this idea. The simplest is that the trial element is placed rather far from the old value without exacting a large energy penalty. Thus one might expect a rather rapid flow through phase space. A second argument is based on the overrelaxation idea as used in minimization schemes such as used in solving linear equations. The position of minimum energy for a given variable is indirectly influenced by the variable itself. When the neighbors were updated, they assumed values which tended to accommodate the position of the current variable. If, however, that variable were allowed to float, they would in general move away and one might expect that the best value for the variable being updated might lie somewhat further away than the position of lowest energy with the neighbors held fixed.

Perhaps the greatest advantage of this algorithm is its computational speed due to simplicity. For both SU(2)and SU(3) I obtained decorrelations per iteration comparable to a heat bath. Although a heat bath is rather easily implemented for SU(2), this is not the case for SU(3) and thus most lattice gauge simulations have been done with a procedure of Metropolis et al. using of order 10 trial changes on any link before proceeding. The present algorithm performs best with but a single hit; indeed, further attempts to change a given variable will just return to earlier trials. Furthermore, the construction of the trial element takes only minimally more computation than a single hit in a standard application of Metropolis et al. This advantage may be even greater for spin systems where there is substantially less overhead involved in calculating the interacting neighborhood of a variable being updated.

Note added. This algorithm and other variations of overrelaxation are discussed in a recent report of Brown and Woch. 9

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