Primitive quantum gates for an SU(2) discrete subgroup: Binary octahedral

Erik J. Gustafson^{1,2,3,4} Henry Lamm,^{1,2,*} and Felicity Lovelace^{5,†}

¹Superconducting and Quantum Materials System Center (SQMS), Batavia, Illinois 60510, USA

²Fermi National Accelerator Laboratory, Batavia, Illinois 60510, USA

³Quantum Artificial Intelligence Laboratory (QuAIL),

NASA Ames Research Center, Moffett Field, California 94035, USA

⁴USRA Research Institute for Advanced Computer Science (RIACS), Mountain View, California 94043, USA

⁵Department of Physics, University of Illinois at Chicago, Chicago, Illinois 60607, USA

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We construct a primitive gate set for the digital quantum simulation of the 48-element binary octahedral (\mathbb{BO}) group. This non-Abelian discrete group better approximates SU(2) lattice gauge theory than previous work on the binary tetrahedral group at the cost of one additional qubit—for a total of six—per gauge link. The necessary primitives are the inversion gate, the group multiplication gate, the trace gate, and the \mathbb{BO} Fourier transform.

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I. INTRODUCTION

The possibilities for quantum utility in lattice gauge theories (LGT) are legion [1–4]. Perhaps foremost, it provides an elegant solution to the sign problem which results in exponential scaling of classical computing resources [5] which precludes large-scale simulations of dynamics, at finite fermion density, and in the presence of topological terms. In order to study these fundamental physics topics, a number of quantum subroutines are required.

The first task is preparing strongly-coupled states of interest including ground states [6–12], thermal states [13–23], and colliding particles [24–34]. For applications to dynamics, the time-evolution operator $U(t) = e^{-iHt}$ must be approximated and many different choices exist; Trotterization [35,36], random compilation [37,38], Taylor series [39], qubitization [40], quantum walks [41], signal processing [42], linear combination of unitaries [38,43], and variational approaches [44–47], each with their own trade-offs. Important alongside state preparation and evolution is the need to develop efficient techniques [48–51] and formulations [52–63] for measuring physical observables. Necessary for acheiving this, one may further use algorithmic improvements such as error mitigation and correction

[64–79], improved Hamiltonians [80,81] and quantum smearing [82] to reduce errors from quantum noise and theoretical approximations. Beyond direct simulations, quantum computers could also accelerate classical lattice gauge theory simulations by reducing autocorrelation [83–86] and optimizing interpolating operators [87,88].

Across this cornucopia, there exist a set of fundamental group theoretic operations that LGT requires [52]. Via this identification of primitive subroutines, the problem of formulating quantum algorithms for LGT can be divided into deriving said group-dependent primitives [89–91] and group-independent algorithmic design [80,82,91,92].

With this notion in mind, one can turn to digitizing the infinite-dimensional Hilbert space of the gauge bosons. Many proposals exist that prioritizing theoretical and algorithmic facets of the problem differently [93–133]. For some digitized theories, there may be no nontrivial continuum limit [122,123,125,134–141]. Furthermore, the efficacy of a gauge digitization can be dimension dependent [89,90,117,142]. While all digitizations consider the relative quantum memory costs, the consequences of digitization on gate costs are more limited. Despite this, recent work has made abundantly clear that the number of expensive fixed-point arithmetic required can vary by orders of magnitude [36,143].

One promising digitization that uses comparatively few qubits and avoids fixed-point arithmetic is the discrete subgroup approximation [66,101,103–107,128,144,145]. This method was explored in the 1970s and 1980s to reduce memory and runtime of Euclidean LGT on classical computers. The replacement of U(1) by \mathbb{Z}_N was considered first [146,147] and eventually extended to the crystal-like subgroups of SU(N) [103,104,128,148–153]. Some studies

^{*}hlamm@fnal.gov [†]fl16@uic.edu

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were even performed with dynamical fermions [154,155]. Theoretical work has established that the discrete subgroup approximation corresponds to continuous groups broken by a Higgs mechanism [156–160].

LGT calculations are performed at fixed lattice spacing $a = a(\beta)$ which for asymptotically free theories approach zero as $\beta \to \infty$. Finite *a* leads to discrepancies from the continuum results, but provided one simulates in the scaling regime below $a_s(\beta_s)$, these errors are well-behaved. On the lattice, the breakdown of the discrete subgroup approximation manifests as a *freeze-out* a_f (or coupling β_f) where the gauge links become "frozen" to the identity. Despite this, the approximation error for Euclidean calculations can be tolerable provided $a_s \gtrsim a_f$ ($\beta_s \lesssim \beta_f$) [104,152]. Further, a connection between the couplings and lattice spacings of Minkowski and Euclidean lattice field theories has been shown [144,153,161], which suggests similarly controllable digitization error on large-scale quantum simulations and a way for determining viable approximations.

The freezing transitions are known in 3 + 1d when the Wilson action is used. Given the known connection between the Wilson action and the Kogut-Susskind Hamiltonian $H_{\rm KS}$ [162], this provides insight into the viable groups for quantum simuations. Approximating U(1) by $\mathbb{Z}_{n>5}$ satisfies $\beta_f > \beta_s$, with $\beta_f \propto 1 - \cos^{-1}(2\pi/n)$. In the case of non-Abelian gauge groups, there are limited number of crystallike subgroups. SU(2), with $\beta_s = 2.2$ has three; the 24-element binary tetrahedral \mathbb{BT} ($\beta_f = 2.24(8)$), the 48-element binary octahedral \mathbb{BO} ($\beta_f = 3.26(8)$), and the 120-element binary icosahedral \mathbb{BI} ($\beta_f = 5.82(8)$) [104]. Thus, while \mathbb{BT} require only 5 qubits, due to its low β_f it is unlikely $H_{\rm KS}$ can be used for quantum simulation but a modified or improved Hamiltonians H_I [80,81] could prove sufficient [104,128,151,152]. The effects on the scaling regime when including fermions is a more subtle question. The structure of discrete groups makes it difficult to leverage traditional Monte Carlo schemes to investigate the transition such as shown in Fig. 1. Qualitatively one could make the following insights; for sufficiently heavy fermions one would expect the scaling behavior to be slightly increased, while the effects for lighter fermions could drastically change the scaling regime in β due to the lower cost of pair creation.

In this work, we consider the smallest crystal-like subgroup of a SU(2) with $a_f > a_s$ — \mathbb{BO} which requires 6 qubits per register. A number of smaller non-Abelian groups have been considered previously. Quantum simulations of the 2*N*-element dihedral groups, D_N , while not crystal-like, have been extensively studied [52,89,101,163]. The 8element \mathbb{Q}_8 subgroup of SU(2) has also been investigated [145]. In [90], quantum circuits for \mathbb{BT} were constructed and resource estimates were obtained using the H_I of [80].

In the interest of studying near-term quantum simulations, we should consider 2 + 1d theories in addition to



FIG. 1. Euclidean calculations of lattice energy density $\langle E_0 \rangle$ of \mathbb{BO} as measured by the expectation value of the plaquette as a function of Wilson coupling β on 8^{*d*} lattices for (top) 2 + 1*d* (bottom) 3 + 1*d*. The shaded region indicates $\beta \leq \beta_s$.

3 + 1d ones. Using classical lattice simulations, we determined $\beta_f > \beta_s$ in both space-times for the Wilson action (See Fig. 1). Thus quantum simulations with \mathbb{BO} can be performed with the Kogut-Susskind Hamiltonian [164], although using an improved Hamiltonian can reduce either qubits or lattice spacing errors [80].

In this paper, the four necessary primitive quantum gates (inversion, multiplication, trace, and Fourier) for quantum simulation of \mathbb{BO} theories on qubit-based computers are constructed. In Sec. II, important group theory for \mathbb{BO} is summarized and the qubit encoding is presented. A brief review of entangling gates used is found in Sec. III. Section IV provides an overview of the primitive gates. This is followed by explicit quantum circuits for each gate for \mathbb{BO} ; the inversion gate in Sec. V, the multiplication gate in Sec. VI, the trace gate in Sec. VII, and the Fourier transform gate in Sec. VIII. Using these gates, Sec. IX presents a resource estimates for simulating 3 + 1d SU(2). We conclude and discuss future work in Sec. X.

II. PROPERTIES OF BO

The simulation of LGT requires defining a register where one can store the state of a bosonic link variable which we call a G – register. To construct the \mathbb{BO} -register in term of integers, it is necessary to map the 48 elements of \mathbb{BO} to the integers [0, 47]. A clean way to obtain this is to write every element of \mathbb{BO} as an ordered product of five generators with exponents written in terms of the binary variables x_i with i = [1, 6]:

$$g = (-1)^{x_1} \mathbf{j}^{x_2} \mathbf{k}^{x_3} \mathbf{u}^{2x_4 + x_5} \mathbf{t}^{x_6}, \tag{1}$$

with

$$\mathbf{u} = -\frac{1}{2}(\mathbb{1} + \mathbf{i} + \mathbf{j} + \mathbf{k})$$
 and $\mathbf{t} = \frac{1}{\sqrt{2}}(\mathbb{1} + \mathbf{i})$ (2)

and **i**, **j**, **k** are the unit quaternions which in the 2d irreducible representation (irrep) correspond to Pauli matrices. With the construction of Eq. (1), the BO-register is given by a binary qubit encoding with the ordering $|x_6x_5x_4x_3x_2x_1\rangle$. While there exist 2⁶ possible state in a 6 qubit register, we only consider the 48 states where $x_4 + x_5 \le 1$ represent the group elements. The states where $x_4 = x_5 = 1$ correspond to *forbidden states*. In this work we will use a short hand $|N\rangle$ where N is the decimal representation of the binary $x_6x_5x_4x_3x_2x_1$. For example,

$$\frac{1}{\sqrt{2}} \begin{pmatrix} -i & -1\\ 1 & i \end{pmatrix} = (-1)^1 \mathbf{j}^0 \mathbf{k}^1 \mathbf{u}^{2 \times 0 + 0} \mathbf{t}^1 \to |100101\rangle = |37\rangle.$$
(3)

The \mathbf{i} , \mathbf{j} , and \mathbf{k} generators anticommute with each other. Additional useful relations are

$$i^{2} = j^{2} = k^{2} = -1,$$
 $u^{3} = 1,$ $t^{2} = i,$
 $ij = k,$ $jk = i,$ $ki = j,$
 $iu = uk,$ $ju = ui,$ $ku = uj,$
 $it = ti,$ $-jt = tk,$ $kt = tj,$ (4)

The character table (Table I) lists important group properties; the different irreps, denoted as ρ_i , can be identified by the value of their character, which correspond to the columns in Table I, acting on each element. The final row indicates the binary state corresponding to group elements in a given character. An irrep's dimension is the value of the character of 1. There are three 1*d* irreps, three 2*d* irreps (one real and two complex), and one 3*d* irrep. To derive the Fourier transform, it is necessary to know a matrix presentation of each irrep. Based on our qubit mapping, given a presentation of -1, **i**, **j**, and **l** we can construct any element of the group from Eq. (1). With the *n*th root of unity $\omega_n = e^{2\pi i/n}$, the matrix presentations of our generators in each irrep are found in Table II.

III. QUBIT GATES

To construct our primitive gates, we chose a universal, albeit redundant, basic qubit quantum gate set. We use the

TABLE I. Character table of \mathbb{BO} from [165] and an enumeration of the elements in the given class.

Size	1	1	12	6	8	8	6	6
Order	: 1	2	4	4	6	3	8	8
$\overline{\rho_1}$	1	1	1	1	1	1	1	1
ρ_2	1	1	-1	1	1	1	-1	-1
ρ_3	2	2	0	2	-1	-1	0	0
ρ_4	2	-2	0	0	1	-1	$\sqrt{2}$	$-\sqrt{2}$
ρ_5	2	-2	0	0	1	-1	$-\sqrt{2}$	$\sqrt{2}$
ρ_6	3	3	-1	-1	0	0	1	1
ρ_7	3	3	1	-1	0	0	-1	-1
ρ_8	4	-4	0	0	-1	1	0	0
$ g\rangle$	0 angle	$ 1\rangle $	$ 34\rangle - 37\rangle$	$ 2\rangle - 7\rangle$	9 angle, 11 angle	8 angle, 10 angle	<i> </i> 32 <i>⟩</i> ,	33 ⟩ ,
							39>	38>
		ŀ	$ 44\rangle - 49\rangle$	>	$ 13\rangle$,	$ 12\rangle$,	$ 41\rangle$,	$ 40\rangle$,
					15>	14>	43>	42>
			$ 52\rangle, 53\rangle$		$ 17\rangle$,	$ 16\rangle$,	$ 50\rangle$,	$ 51\rangle$,
					$ 18\rangle$	19>	$ 54\rangle$	$ 55\rangle$
					$ 20\rangle$,	$ 21\rangle$,		
					$ 22\rangle$	$ 23\rangle$		

Pauli gates p = X, Y, Z and their arbitrary rotation generalizations $R_p(\theta) = e^{i\theta p/2}$. When decomposing onto fault-tolerant devices, how these are decomposed in terms of the $T = \text{diag}(1, e^{i\pi/4})$ gate becomes relevant to resource estimations.

We also use the SWAP gate

$$\mathrm{SWAP}|a\rangle \otimes |b\rangle = |b\rangle \otimes |a\rangle,$$

and CNOT gate

$$\mathrm{CNOT}|a\rangle \otimes |b\rangle = |a\rangle \otimes |b \oplus a\rangle.$$

We further use the multiqubit C^n NOT—of which C^2 NOT is called the Toffoli gate—and CSWAP (Fredkin) gates. The C^n NOT gate consists of one target qubit and *n* control qubits. For example, the Toffoli in terms of modular arithmetic is

$$\mathbb{C}^2 \mathrm{NOT} |a\rangle \otimes |b\rangle \otimes |c\rangle = |a\rangle \otimes |b\rangle \otimes |c \oplus ab\rangle.$$

The CSWAP gate swaps two qubit states if the control is in the $|1\rangle$ state:

$$\begin{split} \mathrm{CSWAP}|a\rangle \otimes |b\rangle \otimes |c\rangle &= |a\rangle \otimes |b(1 \oplus a) \oplus ac\rangle \\ &\otimes |c(1 \oplus a) \oplus ab\rangle. \end{split}$$

IV. OVERVIEW OF PRIMITIVE GATES

One can define any quantum circuit for gauge theories via a set of primitive gates, of which one choice is inversion \mathfrak{U}_{-1} , multiplication \mathfrak{U}_{\times} , trace \mathfrak{U}_{Tr} , and Fourier transform \mathfrak{U}_F [52]. The inversion gate, \mathfrak{U}_{-1} , takes a *G*-register to its inverse:

TABLE II. Matrix representations of the generators used for digitization of BO.

g	-1	j	k	u	t
ρ_1	1	1	1	1	1 _1
$\rho_2 \\ \rho_3$	$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$	$\left(\begin{array}{cc} \omega_3^2 & 0\\ 0 & \omega_3 \end{array}\right)$	$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$
ρ_4	$\begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}$	$\begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$	$\begin{pmatrix} i & 0 \\ 0 & -i \end{pmatrix}$	$\frac{1}{2}\begin{pmatrix} -1-i & -1+i\\ 1+i & -1+i \end{pmatrix}$	$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & -i \\ -i & 1 \end{pmatrix}$
ρ_5	$\begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}$	$\begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$	$\begin{pmatrix} i & 0 \\ 0 & -i \end{pmatrix}$	$\frac{1}{2}\begin{pmatrix} -1-i & -1+i\\ 1+i & -1+i \end{pmatrix}$	$\frac{1}{\sqrt{2}} \begin{pmatrix} -1 & i \\ i & -1 \end{pmatrix}$
$ ho_6$	$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$	$\begin{pmatrix} -1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & -1 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -1 \end{pmatrix}$	$\begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & -1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & -1 \end{pmatrix}$
$ ho_7$	$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$	$\begin{pmatrix} -1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & -1 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -1 \end{pmatrix}$	$\begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 1 & 0 \\ -1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$
	$\begin{array}{cccc} -1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{array}$	$\begin{pmatrix} 0 & -i & 0 & 0 \\ -i & 0 & 0 & 0 \\ 0 & 0 & -i & 0 \\ 0 & 0 & 0 & i \end{pmatrix}$	$\begin{pmatrix} i & 0 & 0 & 0 \\ 0 & -i & 0 & 0 \\ 0 & 0 & 0 & -i \\ 0 & 0 & -i & 0 \end{pmatrix}^{\underline{\omega}}_2$		$ \begin{pmatrix} 0 \\ 0 \\ 1+i \\ -1-i \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 & -1 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} $

$$\mathfrak{U}_{-1}|g\rangle = |g^{-1}\rangle. \tag{5}$$

 \mathfrak{U}_{\times} takes a target *G*-register and changes it to the left product controlled by a second *G*-register:

$$\mathfrak{U}_{\times}|g\rangle|h\rangle = |g\rangle|gh\rangle. \tag{6}$$

Left multiplication is sufficient for a minimal set as right multiplication can be obtained from two applications of \mathfrak{U}_{-1} and \mathfrak{U}_{\times} , albeit resource costs can be further reduced by an explicit construction [80].

Traces of group elements generally define the lattice Hamiltonian. We can implement the evolution with respect to these terms via:



FIG. 2. Example \mathfrak{U}_F from Eq. (8) using $\tilde{\rho}_{i,j} = \sqrt{d_{\rho}/|G|}\rho_{i,j}$ where $\rho_{i,j} = \rho_i(g_j)$. This example has four irreps with $d_1 = d_2 < d_3 < d_4$. \mathfrak{U}_F is square since $\sum_{\rho} d_{\rho}^2 = |G|$.

$$\mathfrak{U}_{\mathrm{Tr}}(\theta)|g\rangle = e^{i\theta\mathrm{Re}\,\mathrm{Tr}\,g}|g\rangle. \tag{7}$$

The final gate of this set is \mathfrak{U}_F . The Fourier transform, \hat{f} , of a function f over a finite G is

$$\hat{f}(\rho) = \sum_{g \in G} \sqrt{\frac{d_{\rho}}{|G|}} f(g)\rho(g), \tag{8}$$

where |G| is the size of the group, d_{ρ} is the dimensionality of the irrep ρ . The inverse transform is given by

$$f(g) = \sum_{\rho \in \hat{G}} \sqrt{\frac{d_{\rho}}{|G|}} \operatorname{Tr}(\hat{f}(\rho)\rho(g^{-1})),$$
(9)

where the dual \hat{G} is the set of irrep of G. A paradigm of this unitary matrix is shown in Fig. 2 which can then be transformed into a gate. \mathfrak{U}_F then acts on a single *G*-register with some amplitudes f(g) which rotate it into the irrep basis:

$$\mathfrak{U}_F \sum_{g \in G} f(g) |g\rangle = \sum_{\rho \in \hat{G}} \hat{f}(\rho)_{ij} |\rho, i, j\rangle.$$
(10)

V. INVERSION GATE

Consider a BO-register storing the group element given by $g = (-1)^{x_1} \mathbf{j}^{x_2} \mathbf{k}^{x_3} \mathbf{u}^{2x_4+x_5} \mathbf{t}^{x_6}$. The effect of the inversion gate on this register is to transform it to

$$|g\rangle = |x_1x_2x_3x_4x_5x_6\rangle \rightarrow |g^{-1}\rangle = |y_1y_2y_3y_4y_5y_6\rangle.$$
 (11)

Where y_i must be determined. With Eq. (1), g^{-1} is

$$g^{-1} = \mathbf{t}^{8-x_6} \mathbf{u}^{3-(2x_4+x_5)} \mathbf{k}^{x_3} \mathbf{j}^{x_2} (-1)^{x_1+x_2+x_3}.$$
(12)

A systematic way to build the \mathfrak{U}_{-1} of \mathbb{BO} is to embed the inverse gates of subgroups of \mathbb{BO} . We start with the expression of Eq. (12), and begin by reordering the \mathbb{Q}_8 subgroup so that the element is of the form:

$$\mathbf{k}^{x_3} \mathbf{j}^{x_2} (-1)^{x_1 + x_2 + x_3} = (-1)^{a_1} \mathbf{j}^{a_2} \mathbf{k}^{a_3}, \qquad (13)$$

where $a_1 = x_1 + x_2 + x_3 + x_2x_3$ and $a_2 = x_2$, $a_3 = x_3$. The next step is commuting through **u** to obtain

$$\mathbf{u}^{3-(2x_4+x_5)}(-1)^{a_1}\mathbf{j}^{a_2}\mathbf{k}^{a_3} = (-1)^{b_1}\mathbf{j}^{b_2}\mathbf{k}^{b_3}\mathbf{u}^{2b_4+b_5}, \quad (14)$$

which corresponds to the \mathbb{BT} inverse gate where

$$b_{2} = a_{2}(1 + x_{4}) + a_{3}(x_{4} + x_{5}),$$

$$b_{3} = a_{3}(1 + x_{5}) + a_{2}(x_{4} + x_{5}),$$

$$b_{4} = x_{5},$$

$$b_{5} = x_{4},$$
(15)

while $b_1 = a_1$ is unchanged. Finally, commuting through **t**,

$$\mathbf{t}^{8-x_6}(-1)^{b_1}\mathbf{j}^{b_2}\mathbf{k}^{b_3}\mathbf{u}^{2b_4+b_5} = (-1)^{y_1}\mathbf{j}^{y_2}\mathbf{k}^{y_3}\mathbf{u}^{2y_4+y_5}\mathbf{t}^{y_6}.$$
 (16)

Propagating \mathbf{t}^{x_6} through the \mathbb{Q}_8 portion yields $g^{-1} = (-1)^{c_1} \mathbf{j}^{c_2} \mathbf{k}^{c_3} \mathbf{t}^{x_6} \mathbf{u}^{2b_4+b_5}$, where



FIG. 3. Quantum circuit for \mathfrak{U}_{-1} for \mathbb{BO} . The subgates correspond to \mathfrak{U}_{-1} for \mathbb{Q}_8 (BT) in orange (blue).

$$c_1 = b_1 + x_6(1 - b_3)(1 - b_2),$$

$$c_2 = (1 - b_3)x_6 + (1 - x_6)b_2,$$

$$c_3 = (1 - b_2)x_6 + (1 - x_6)b_3.$$

Finally, propagating through \mathbf{t}^{x_6} through $\mathbf{u}^{2b_4+b_5}$ yields,

$$y_{1} = c_{1} + x_{6}(b_{4}(1 - c_{2}) + c_{3}b_{5}),$$

$$y_{2} = c_{2} + b_{4}x_{6},$$

$$y_{3} = c_{3} + (b_{4} + b_{5})x_{6},$$

$$y_{4} = x_{6}b_{5} + (1 - x_{6})b_{4},$$

$$y_{5} = x_{6}b_{4} + (1 - x_{6})b_{5},$$
 (17)

with $y_6 = x_6$. Together, this yields for $|g\rangle$ to $|g^{-1}\rangle$:

$$y_{1} = x_{1} + x_{2}x_{3}x_{6} + x_{2}x_{3} + x_{2}x_{5}x_{6} + x_{2}x_{6} + x_{2} + x_{3}x_{4}x_{6} + x_{3}x_{6} + x_{3} + x_{4}x_{6} + x_{6},$$

$$y_{2} = x_{2}x_{4} + x_{2}x_{5}x_{6} + x_{2}x_{6} + x_{2} + x_{3}x_{4}x_{6} + x_{3}x_{4} + x_{3}x_{5} + x_{3}x_{6} + x_{5}x_{6} + x_{6},$$

$$y_{3} = (x_{2} + x_{3})(x_{5} + x_{6}) + x_{6}(x_{4}(x_{3} + 1) + x_{5}(x_{2} + 1)) + x_{2}x_{4} + x_{3} + x_{6},$$

$$y_{4} = (x_{4} + x_{5})x_{6} + x_{5},$$

$$y_{5} = (x_{4} + x_{5})x_{6} + x_{4},$$

(18)

where $y_6 = x_6$ is unaffected. We provide the quantum circuit for this gate in Fig. 3.

VI. MULTIPLICATION GATE

Given two BO-registers $|g\rangle = |\prod x_i\rangle$ and $|h\rangle = |\prod y_i\rangle$, we want $|gh\rangle = |\prod z_i\rangle$. Here, we again decompose

$$\mathfrak{U}_{\times} = \mathfrak{U}_{\times,-1}\mathfrak{U}_{\times,\mathbf{j}}\mathfrak{U}_{\times,\mathbf{k}}\mathfrak{U}_{\times,\mathbf{u}_{2}}\mathfrak{U}_{\times,\mathbf{u}_{1}}\mathfrak{U}_{\times,\mathbf{t}}$$
(19)

into gates controlled by a single qubits. For each step, we use temporary variables indexed by other letters, e.g. a_i . The first gate, $\mathfrak{U}_{\times,t}$, controlled by x_6 , sets:

$$a_{1} = y_{2}y_{4}y_{6} + y_{3}y_{4}y_{6} + y_{2}y_{5}y_{6} + y_{2}y_{3} + y_{2}y_{4} + y_{3}y_{5} + y_{3}y_{6} + y_{4}y_{6} + y_{1} + y_{3} + y_{5}, a_{2} = y_{4}y_{6} + y_{3} + y_{5} + y_{6}, a_{3} = y_{5}y_{6} + y_{2} + y_{4} + y_{5} + y_{6}, a_{4} = y_{5}, a_{5} = y_{4}, z_{6} = y_{6} + 1.$$
(20)

Then one acts with $\mathfrak{U}_{\times,\mathbf{u}_1}$ which is controlled by x_5 :



FIG. 4. The decomposition of the multiplication gate, \mathfrak{U}_{\times} , into a product of multiplication by each generator.

$$b_2 = a_2 + a_3,$$

 $b_3 = a_2,$
 $b_4 = a_5,$
 $b_5 = a_4 + a_5 + 1,$ (21)

where a_1 and z_6 do not change. Next, applying $\mathfrak{U}_{\times, \mathbf{u}_2}$ controlled by x_4 ,

$$c_{2} = b_{3},$$

$$c_{3} = b_{2} + b_{3},$$

$$z_{4} = b_{4} + b_{5} + 1,$$

$$z_{5} = b_{4},$$
(22)

with a_1 and z_6 unchanged. $\mathfrak{U}_{\times,\mathbf{k}}$, controlled by x_3 , sets

$$d_1 = a_1 + c_2 + c_3,$$

$$z_3 = c_3 + 1,$$
(23)

where c_2 , z_4 , z_5 , and z_6 are unchanged. Then, $\mathfrak{U}_{\times,\mathbf{j}}$ controlled by x_2 is

$$e_1 = d_1 + c_2,$$

 $z_2 = c_2 + 1,$ (24)

with z_3 through z_6 unchanged. Finally, $\mathfrak{U}_{\times,-1}$ needs $z_1 = e_1 + 1$ controlled by x_1 i.e. a CNOT. The full \mathfrak{U}_{\times} is in Fig. 4.

VII. TRACE GATE

For simulating LGT, the Hamiltonian requires gaugeinvariant operators. Without matter, all terms can be constructed from traces of g—and thus \mathfrak{U}_{Tr} —in the fundamental irrep, ρ_4 . Table I provides us with ReTr (g_i) . In previous works [52,89,90], \mathfrak{U}_{Tr} was derived from a Hamiltonian, H_{Tr} with $\operatorname{ReTr}(g_i)$ as eigenvalues. For our \mathbb{BO} -register, H_{Tr} would require 20 Pauli strings which results in $\mathfrak{U}_{\operatorname{Tr}} = e^{i\theta H_{\operatorname{Tr}}}$ decomposing into at least 20 $R_Z(\theta)$ gates. In a fault-tolerant calculation, R_Z gates require synthesis from T gates, and thus can be unduly expensive.

Here, we explore a different implementation of \mathfrak{U}_{Tr} that is advantageous for discrete groups where $\operatorname{ReTr}(g)$ are limited by the number of conjugacy classes. In this case, \mathfrak{U}_{Tr} could be decomposed into two gates; a gate U_{conj} to map the 48 $|g\rangle$ to the 8 conjugacy classes $|c\rangle$, and a gate U_{Tr} which computes the traces for each conjugacy class. Together, we obtain a circuit for $\mathfrak{U}_{Tr} = U_{conj}U_{Tr}(\theta)U_{conj}^{\dagger}$ which should have fewer $R_Z(\theta)$ at the cost of additional C^N NOT gates which have ϵ -independent T gate costs. Taking the assignment of Table III for the traces we derive

$$(-1)^{x_1} \mathbf{j}^{x_2} \mathbf{k}^{x_3} \mathbf{u}^{2x_4+x_5} \mathbf{t}^{x_6} \mapsto (v_1, v_2, v_3),$$

where

$$v_{1} = (x_{6} + 1)(x_{1} + (x_{4} + x_{5})(1 + x_{2})(1 + x_{3})) + x_{6}(x_{5}x_{1} + x_{4}(1 + x_{1}) + (1 + x_{4} + x_{5})(x_{2}x_{3} + x_{1})), v_{2} = (x_{6} + 1)((1 + x_{2})(1 + x_{3})(1 + x_{4} + x_{5})) + x_{6}(x_{5}x_{2} + x_{4}(1 + x_{3}) + (1 + x_{4} + x_{5})(1 + x_{2} + x_{3})), v_{3} = (x_{6} + 1)(x_{4} + x_{5}) + x_{6}(x_{5}x_{2} + x_{4}(1 + x_{3}) + (1 + x_{4} + x_{5})(1 + x_{2} + x_{3})).$$
(25)

TABLE III. Trace elements and associated bit strings for the conjugacy classes stored in the ancilla. Note, v_1 is a sign bit.

$\operatorname{Re}\operatorname{Tr}\left(g ight)$	0	0	2	-2	1	-1	$\sqrt{2}$	$-\sqrt{2}$
<i>v</i> 1	0	1	0	1	0	1	0	1
v2	0	0	1	1	0	0	1	1
<i>v</i> 3	0	0	0	0	1	1	1	1



FIG. 5. A qubit implementation of \mathfrak{U}_{Tr} .

A qubit-based circuit for \mathfrak{U}_{Tr} is shown in Fig. 5. With this, we will estimate the reduction in fault-tolerant resource gate costs in Sec. IX compared to the H_{Tr} method.

VIII. FOURIER TRANSFORM

The standard *n*-qubit quantum Fourier transform (QFT) [166] corresponds to the quantum version of the fast Fourier transform of \mathbb{Z}_{2^n} . Quantum Fourier transforms over some non-Abelian groups exist [89,167–170]. Alas, for the crystal-like subgroups, efficient QFT circuits are currently unknown [171]. In general, no clear algorithmic way to construct the QFT exists. Thus, we instead construct a suboptimal \mathfrak{U}_F from Eq. (8) from the irreps.

Since \mathbb{BO} has 48 elements, on a qubit device \mathfrak{U}_F must be embedded into a larger $2^d \times 2^d$ unitary. The columns index qubit bitstrings from $|0\rangle$ to $|63\rangle$ where the physical states are given by the subset of $|g\rangle$ in Table I. We then index the irreducible representation ρ_i sequentially from i = 1 to i = 8, and construct the matrix, including appropriate padding for the forbidden states. This matrix was then passed to the Qiskit v0.43.1 transpiler, and an optimized version of \mathfrak{U}_F needed 1839 CNOTs, 166 R_X , 1996 R_Y , and 3401 R_Z gates¹; the Fourier gate is the most expensive qubit primitive and dominates simulation costs.

IX. RESOURCE ESTIMATES

Clearly quantum practicality with \mathbb{BO} will require error corrections. Since the Eastin-Knill theorem restricts QEC codes from having a universal transversal sets of gates [172], compromises must be made. Typically, the Clifford gates are transversal [173–177] while the T gate is not. Thus T gate counts are often used in fault-tolerant resource analysis [173,178]. Recently, novel universal sets have been

proposed with transversal \mathbb{BT} , \mathbb{BO} , \mathbb{BI} gates [179–182] which deserve investigation for use with LGT.

The Toffoli gate requires 7 T gates [173] and one method for constructing the CⁿNOT gates is with $2\lceil \log_2 n\rceil - 1$ Toffoli gates and n - 2 dirty ancilla qubits² which can be reused later [173,183,184]. We arrive at the cost for the R_Z gates via [185] where these gates can be approximated to precision ϵ with on average $1.15\log_2(1/\epsilon)$ T gates (and at worst $-9 + 4\log_2(1/\epsilon)$ [186]). Further, R_Y and R_X can be replaced by at most $3 R_Z$. With these, we can construct gate estimates for BO (See Table IV). We note that for $\epsilon \leq 10^{-6}$, the $20 R_Z(\theta)$ required for $e^{i\theta H_{\rm Tr}}$ is more expensive than the $\mathfrak{U}_{\rm Tr}$ constructed using $U_{\rm conj}$.

Primitive gate costs for implementing $H_{\rm KS}$ [164] and H_I [80], per link per Trotter step δt are shown in Table V. Using these, we can determine the total T gate count $N_T^H = C_T^H \times dL^d N_t$ for a *d* spatial lattice simulated for a time $t = N_t \delta t$. We find that for $H_{\rm KS}$

$$C_T^{\text{KS}} = 2863(d-1) + (22737.8 + 2.3d)\log_2\frac{1}{\epsilon}.$$
 (26)

With this, the total synthesis error ϵ_T can be estimated as the sum of ϵ from each R_Z . In the case of H_{KS} this is

$$\epsilon_T = 2(9886 + d)dL^dN_t \times \epsilon. \tag{27}$$

If one looks to reduce lattice spacing errors for a fixed number of qubits, one can use H_I which would require

$$C_T^I = 11949d - 10157 + (45473.3 + 6.9d) \log_2 \frac{1}{\epsilon}, \quad (28)$$

where the total synthesis error is

¹In [90], the Qiskit v0.37.2 was used, yielding for the BT \mathfrak{U}_F of 1025 CNOTs, 1109 R_Y , and 2139 R_Z gates. With the v0.43.1 transpiler, we obtain 442 CNOTs, 40 R_X , 494 R_Y , and 835 R_Z .

²A *clean* ancilla is in state $|0\rangle$. *Dirty* ancillae have unknown states.

TABLE IV. Number of physical T gates and clean ancilla required to implement logical gates for (top) basic gates taken from [173] (bottom) primitive gates for \mathbb{BO} .

Gate	T gates	Clean ancilla
C ² NOT	7	0
C ³ NOT	21	1
C ⁴ NOT	35	2
CSWAP	7	0
R_Z	1.15 $\log_2(1/\epsilon)$	0
\mathcal{U}_{-1}	112	1
$\mathcal{U}_{ imes}$	392	4
${\cal U}_{ m Tr}$	$340^{a} + 4.6 \log_{2}(1/\epsilon)$	2
\mathcal{U}_{FT}	11370.1 $\log_2(1/\epsilon)$	0

^aWith 6 additional ancilla, this can be reduced to 188.

TABLE V. Number of primitive gates per link per δt neglecting boundary effects as a function of d for H_{KS} and H_I .

Gate	$N[H_{\rm KS}]$	$N[H_I]$
$\overline{\mathfrak{U}_F}$	2	4
$\mathfrak{U}_{\mathrm{Tr}}$	$\frac{1}{2}(d-1)$	$\frac{3}{2}(d-1)$
\mathfrak{U}_{-1}	$\bar{3}(d-1)$	2 + 11(d - 1)
\mathfrak{U}_{\times}	6(d-1)	4 + 26(d - 1)

$$\epsilon_T = 2(19771 + 3d)dL^dN_t \times \epsilon. \tag{29}$$

Following [57,90,143], we will make resource estimates based on our primitive gates for the calculation of the shear viscosity η on a $L^3 = 10^3$ lattice evolved for $N_t = 50$, and total synthesis error of $\epsilon_T = 10^{-8}$. Considering only the time evolution and neglecting state preparation (which can be substantial [12,14,17–21,28,30–33,51,82,84,87,187]), Kan and Nam estimated 3×10^{19} T gates would be required for an pure-gauge SU(2) simulation of $H_{\rm KS}$. This estimate used a truncated electric-field digitization which requires substantial fixed-point arithmetic—greatly inflating the T gate cost. Using BT, we presently estimate that 1.1×10^{11} T gates would be sufficient³ but only H_I could be used. Here, the BO group would require 4.1×10^{11} T gates if H_I were used—albeit with smaller lattice spacing errors than Kan and Nam since they used H_{KS} . Thus \mathbb{BO} reduces the gate costs of [143] by 10⁸ for fixed *L* by avoiding quantum fixed-point arithmetic while allowing for smaller lattice spacings that \mathbb{BT} . If H_{KS} were used the T gate cost is reduced to 2.0×10^{11} . Alternatively, if we take the heuristic of [80] that for $a = \mathcal{O}(0.1 \text{ fm})$ using H_I could reduce *L* by half for fixed lattice spacing errors with only 4.9×10^{10} T gates. Similar to \mathbb{BT} , \mathfrak{U}_F dominates the simulations—99% and 98% of the total cost of simulation H_{KS} and H_I , respectively.

X. CONCLUSIONS

In this paper, we constructed the necessary primitive gates for the simulation of \mathbb{BO} —a crystal-like subgroup of SU(2)—gauge theories and quantum resource estimates were made for the simulation of pure SU(2) shear viscosity. Compared to previous fault-tolerant qubit estimates of electric basis truncations, we require 10^8 fewer T gates by avoiding quantum fixed point arithmetic via the discrete group approximation. Further, reducing digitization error compared to \mathbb{BT} increase the total cost by a factor of ~4 perhaps suggesting a $|G|^2$ scaling.

Looking forward, primitive gates should be constructed for BI and to the subgroups of SU(3). At the cost of more qubits, BI would allow smaller digitization error and lattice spacings. To further reduce the qubit-based simulation gate costs for all discrete subgroup approximations, the formalism for deriving the quantum Fourier transform for crystallike groups remains of paramount interest.

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³This supersedes [90] due to the improved transpiler, proper consideration of ϵ_T , and better cost estimates of R_X , R_Y .

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