

Coprocessors for quantum devices

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Quantum devices, from simple fixed-function tools to the ultimate goal of a universal quantum computer, will require high-quality, frequent repetition of a small set of core operations, such as the preparation of entangled states. These tasks are perfectly suited to realization by a coprocessor or supplementary instruction set, as is common practice in modern CPUs. In this paper, we present two quintessentially quantum coprocessor functions: production of a Greenberger-Horne-Zeilinger state and implementation of optimal universal (asymmetric) quantum cloning. Both are based on the evolution of a fixed Hamiltonian. We introduce a technique for deriving the parameters of these Hamiltonians based on the numerical integration of Toda-like flows.

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I. INTRODUCTION

The task of quantum state synthesis [1,2] lies at the heart of any potential quantum technology: before a quantum protocol can be run, be it a Bell test [3], quantum key distribution [4], quantum cloning [5–7], random-number generation [8], or quantum computation [9], a nontrivial quantum resource must be prepared. The required resource might be a fixed quantum state such as a Bell state, W state, cluster [9], or Greenberger-Horne-Zeilinger (GHZ) state, or it might depend on a small input, such as the unknown state of a qubit. The availability of these resource states is the source of the power of quantum technologies. Repeated demands for the same resource state make it vital to concentrate on their accurate functioning. This suggests developing a device that accomplishes that single task, replacing a complex sequence of quantum gates. These might provide the first step in a quantum protocol (i.e., the core functionality of a particular quantum technological device) or operate as a fixed-function subroutine within a quantum computer, much as today's classical processors provide enhanced instruction sets (e.g., Streaming SIMD Extensions (SSE) and Advanced Vector Extensions (AVX)) or coprocessors. Our aim is to produce the desired states and transformations by the free evolution of a Hamiltonian whose parameters have been specifically tuned for the task. By doing this directly with the system's Hamiltonian for any relevant experimental scenario, whether this is in the solid state [10,11], trapped ions [12], or even photonic systems [13–15], we ensure that the state is produced as accurately and as quickly as possible, reducing the opportunities for external influences to degrade the resource.

One special case of this has been extensively studied: perfect quantum state transfer [16–20], wherein an unknown quantum state can be transported between the two extremes of a one-dimensional chain of spins. This example demonstrates the power of the approach: it is twice as fast as the equivalent quantum gate sequence [21], and many of the error modes are

relegated to the manufacturing process; they can be identified prior to use and corrected or simply rejected until a higher-quality version is produced [15]. In addition to perfect state transfer, the same device can create graph states [22], of which the cluster states and GHZ states are special cases (albeit in an unusual basis). Minor modifications [19,23,24] have also permitted the creation of Bell states between the extremal sites of the chain. More recently [1,2], new systems have been created to facilitate the synthesis of arbitrary one-excitation states on the chain, such as W states between subsets of sites, while more exotic interactions have been shown to cause signal amplification, ideal for enhancing measurement signals [25].

In this paper, we develop a coprocessor that creates GHZ states in a particularly straightforward manner (see Sec. II). Moreover, the one-dimensional transverse Ising model that we introduce is highly appropriate for many experimental scenarios from superconducting qubits [10,11] to trapped ions [12] and demonstrates a reasonable robustness to experimental imperfections (see Sec. II A). In Sec. III, we also show how the state synthesis solutions of [1,2] can be combined with the GHZ coprocessor to implement optimal universal cloning of one unknown qubit to N clones [7,26,27]. While the circuits in [27] were probabilistic in nature, the present scheme constitutes a reliable implementation of optimal universal asymmetric cloning, demonstrating that fixed-function devices can perform transformations based on a small input space and realize highly nontrivial quantum properties. In Sec. III A 1, we also introduce the state synthesis problem for uniformly coupled networks (as compared to chains with engineered couplings) and demonstrate that some hypercubes are useful for generating the uniform superpositions (W states) that are desirable for symmetric cloning.

Crucial to the specification of both coprocessors is the numerical discovery of appropriate Hamiltonian parameters. We introduce a technique based on the numerical integration of a differential equation, the Toda flow. This is the main focus of Sec. III A, along with discussions of convergence issues in the Appendix. Variants of this [28] and good techniques for its numerical integration [29–31] have been extensively studied in the numerical-analysis and numerical-methods literature.

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anticommutation of the fermions. Hence, the product involving pairs of fermions is unchanged, and the final state must become

$$\frac{1}{2^N} \left[\prod_{n=1}^{N-1} (\mathbb{1} + i c_{2n} c_{2n+1}) \right] (\mathbb{1} - c_{2N}).$$

This is the same (up to a global phase) as the pure state

$$|\text{GHZ}\rangle = \frac{1}{\sqrt{2}} (|0\rangle^{\otimes N} - i|1\rangle^{\otimes N}). \quad (3)$$

We have successfully engineered an Ising chain that creates GHZ states by its natural dynamics. The transfer time scales linearly with N if a maximum coupling strength is imposed, which is the best possible scaling for a one-dimensional system [38].

The evolution after subsequent periods of $t_0/2$ is readily determined since

$$(\mathbb{1} + c_1) \rightarrow (\mathbb{1} - c_{2N}) \rightarrow (\mathbb{1} - c_1) \rightarrow (\mathbb{1} + c_{2N}) \rightarrow (\mathbb{1} + c_1),$$

so the state evolves as

$$|0\rangle^{\otimes N} \rightarrow |\text{GHZ}\rangle \rightarrow -i|1\rangle^{\otimes N} \rightarrow -iZ_1|\text{GHZ}\rangle \rightarrow -i|0\rangle^{\otimes N} \quad (4)$$

(indeed, the phase Z_1 might equally well be applied on any qubit). The only part that is not justified is the global phase. However, we know that

$$|0\rangle^{\otimes N} \rightarrow |\text{GHZ}\rangle, \quad |1\rangle^{\otimes N} \rightarrow Z_1|\text{GHZ}\rangle,$$

so by linearity, we can evaluate $|\text{GHZ}\rangle \rightarrow -i|1\rangle^{\otimes N}$.

It is also now straightforward to determine the evolution of other basis states, which we will make use of in Sec. III. Consider an input state $|x\rangle$ for $x \in \{0, 1\}^N$. We can alternatively write this as $X_x|0\rangle^{\otimes N}$, where X_x has Pauli X operators applied on the sites where the bit value of the string x is 1 and identity on the other sites. Each $X_n = i c_{2n-1} c_{2n}$ and therefore evolves to X_{N+1-n} . So, if x_R is the reversal of bit string x , we have

$$|x\rangle = X_x|0\rangle^{\otimes N} \mapsto X_{x_R}|\text{GHZ}\rangle. \quad (5)$$

These models are well suited to near-term experimental realization. For example, [36] uses Rydberg atoms in a chain to produce a Hamiltonian which can, in principle, be tuned to give an Ising model of up to 51 qubits. The main challenge is to make the fields B_n and J_n a similar strength. Currently, $|B_n/J_n| \leq \delta \lesssim 0.1$. While incompatible with the standard perfect-state-transfer couplings [17], other techniques such as that in [33] can return the couplings for suitable perfect-transfer schemes. However, this means that there are two eigenvalues of h_1 which are separated by $O(\delta^N)$ (corresponding to the unpaired fermions of the Majorana chain described by Kitaev [39]). Consequently, the time for generating the GHZ state scales as $\Omega(\delta^{-N})$, which is currently prohibitive.

A. Robustness of GHZ synthesis

While we have identified the main experimental challenge, the necessary accuracy of engineering for the system parameters could be a concern. This is particularly acute in the case of GHZ-state synthesis because the state in Eq. (2) is described in terms of a large number (up to $2N - 1$) of Majorana fermions. If each is transported with some subunitary fidelity,

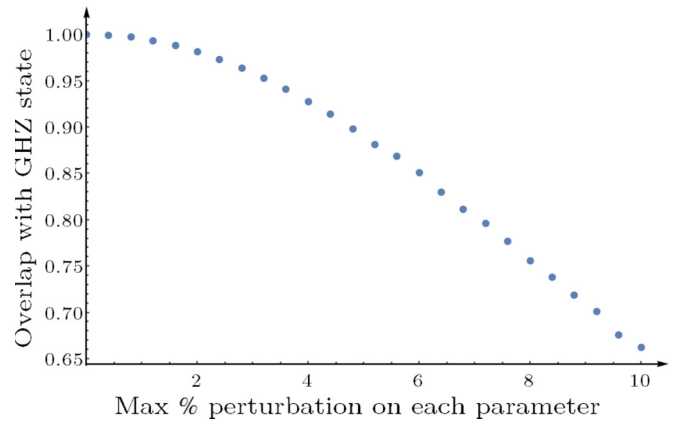


FIG. 1. Average overlap with $|\text{GHZ}\rangle$ for $N = 21$ when the parameters B_n and J_n are all chosen uniformly at random within a range of $\pm x\%$ of the value they should be (the “standard” perfect transfer couplings [17]), using 1000 samples.

the overall success of the synthesis could be quite minimal. In fact, the situation is not nearly so bleak, and the system has a good tolerance of these imperfections. A good estimate of the overlap of the output state $\langle \text{GHZ} | e^{-iH_1 t_0/2} | 0 \rangle^{\otimes N}$ is

$$\approx \frac{1 + |F|}{2^N} \sqrt{\det(e^{-ih_1 t_0} h_0 e^{ih_1 t_0} h_0 - \mathbb{1})},$$

where $F = \langle 2N | e^{-ih_1 t_0} | 1 \rangle$ is the single-excitation transfer fidelity and

$$h_0 = \sum_{n=1}^{N-1} |2n\rangle \langle 2n+1| - |2n+1\rangle \langle 2n|.$$

This exactly evaluates the evolution of $\prod_{n=1}^{N-1} (\mathbb{1} + Z_n Z_{n+1})$, comparing it to itself, by writing it as a fermionic Gaussian state [40,41]. The additional effect of the single excitation c_1 is then approximated, ignoring possible interactions with the Gaussian component. This approximation facilitates numeric simulations, and Fig. 1 demonstrates the effect on a chain of size 21.

It must be emphasized that these results are very basic, merely measuring the overlap with the target state. When this device is made, we will characterize the state that is produced and adapt for its imperfections, such as applying an optimized choice of local unitaries. This can serve only to increase the figure of merit. Or can we witness the presence of different types of entanglement, in particular, k -body entanglement for $k \sim N$? Existing entanglement witnesses are not yet sophisticated enough to be able to discriminate this.

On the other hand, one thing that we cannot easily do is replace the perfect-transfer couplings with a coupling scheme that achieves nearly perfect transfer (but with a shorter transfer time, making the system less susceptible to noise and some perturbations), such as those suggested in [42]. Those schemes are tuned specifically for end-to-end transfer. They generate high transfer fidelity between $|1\rangle$ and $|2N\rangle$ at a higher speed, at the cost of the transfer fidelity between intermediate sites. However, GHZ synthesis requires high-quality transfer for all pairs $|n\rangle \rightarrow |2N+1-n\rangle$. For example, the optimal $N = 21$ solution from [42] has a $1 \rightarrow 42$ excitation transfer fidelity of

0.993 (which is roughly reproduced by the 3% perturbed chains in Fig. 1) but generates the GHZ state with only an overlap of 0.762 due to the vastly lower transfer fidelities in the middle of the chain, such as for $5 \rightarrow 38$, which is less than 0.4.

B. GHZ creation in the XY model

A generalization of H_I can be written as

$$H_{ZY} = \sum_{n=1}^N B_n X_n + \sum_{n=1}^{N-1} J_n (1 + \gamma_n) Z_n Z_{n+1} + J_n (1 - \gamma_n) Y_n Y_{n+1}.$$

This model is also a free-fermion model and has the same Majorana fermions as H_I . Starting from the same initial state as described by Eq. (2), the only possible difference is what those Majorana fermions can evolve into, which is again governed by a $2N \times 2N$ matrix h_γ , similar to h_1 . We are interested in whether this broader class of Hamiltonians can also produce the GHZ state, again in the hope of improving experimental viability. For pedagogical simplicity, we will fix $\gamma_n = \gamma$ for all n , although there is no such restriction arising in the mathematics. We could equally well consider the Hadamard-transformed version of this Hamiltonian, which is the more familiar XY model:

$$H_{XY} = \sum_{n=1}^N B_n Z_n + \sum_{n=1}^{N-1} J_n (1 + \gamma_n) X_n X_{n+1} + J_n (1 - \gamma_n) Y_n Y_{n+1}.$$

In this case, the initial state would be $(H|0\rangle)^{\otimes N}$, and the final state would be $H^{\otimes N}|\text{GHZ}\rangle$.

We already know two solutions for this matrix. At $\gamma = 1$ (vanishing YY terms), we have already fixed $B_n = J_{2n-1}^{(2N)}$ and $J_n = J_{2n}^{(2N)}/2$, while for $\gamma = 0$, $B_n = N$ and $J_n = J_n^{(N)}/2$ come from perfect state transfer. Indeed, this last solution is the usual perfect-state-transfer chain (using H_{XY}), and it was already observed in [22] that this system is capable of creating the GHZ state (in a nonobvious basis).

Solutions for both values of γ have the same eigenvalues, $\pm 1, \pm 3, \pm 5, \dots, \pm(2N - 1)$. We are hence interested in intermediate values of γ , with the same spectrum. We shall do this by providing a numerical routine to interpolate between the two known solutions. We believe the form of the isospectral transformation is new to the spin-chain community, although it is well studied in the numerical-methods and -analysis literature [28–31].

We permute the elements of h_γ , grouping odd-numbered and even-numbered basis elements together. The matrix then decomposes as

$$h_\gamma = i|0\rangle\langle 1| \otimes X(\gamma) - i|1\rangle\langle 0| \otimes X^T(\gamma),$$

where $X(\gamma)$ is a nonsymmetric matrix,

$$X = \sum_{n=1}^N B_n |n\rangle\langle n| + \sum_{n=1}^{N-1} J_n (1 + \gamma) |n\rangle\langle n+1| + \sum_{n=1}^{N-1} J_n (1 - \gamma) |n+1\rangle\langle n|.$$

The spectrum of h_γ is $\pm\lambda_i$, where λ_i are the singular values of X . Hence, it is sufficient to perform an isospectral transformation on X . To achieve this, we observe that if A and B are anti-Hermitian, then

$$X(t) = e^{-B} X e^A$$

describes an isospectral flow. Taking the derivative,

$$\frac{dX}{dt} = XA - BX. \quad (6)$$

So any small, anti-Hermitian A, B will achieve an isospectral transformation; we just need to select them so that X retains the properties that we want it to:

(1) X is tridiagonal; that is, $\langle m|XA|n\rangle = \langle m|BX|n\rangle$ for any n, m such that $|n - m| > 1$.

(2) C is centrosymmetric in the sense $\langle n|X|m\rangle = \langle N + 1 - m|X|N + 1 - n\rangle$. We anticipate this being a necessary condition in the same way that it is for state transfer [19], although this is unproven.

(3) We require $\langle n + 1|X|n\rangle / \langle n|X|n + 1\rangle = \frac{1-\gamma}{1+\gamma}$ to be the same for all n .

As conditions on the matrix X , if it satisfies them at any given value, we can ensure they are upheld on subsequent values by imposing them on the derivatives. For a given X , each condition is linear in the coefficients of A and B , and the number of constraints coincides perfectly with the number of coefficients, permitting solution. By performing a numerical integration starting from a known solution for $X(0)$ (which, as already observed, we know for $\gamma = 0, 1$) any desired value of γ can be arrived at.

Equation (6) can be integrated following two different philosophies. First, one can integrate it directly, i.e., setting the next X to be

$$X \mapsto X + \delta(XA - BX).$$

The structural aspects of X are preserved exactly, but the isospectral transformation is accurate only to $O(\delta^2)$ for each step of size δ , giving an overall accuracy of $O(\delta)$. Alternatively, one can perform the update

$$X \rightarrow e^{-B} X e^A.$$

This unitary transformation is isospectral, but the structural properties such as tridiagonality are accurate only to $O(\delta^2)$. Nevertheless, there is the facility to compensate for any error in the next step, preventing it from accumulating during the integration. Moreover, if δ is shrunk as a solution converges, then the accuracy is arbitrarily good. Throughout this paper, we use first-order (Euler) integrations. While they appear to serve very well, isospectral flows of this form are often challenging to integrate [43], and novel techniques such as Runge-Kutta-Munthe-Kaas, have been developed [29,30]. These may be used to improve performance in the future.

We conclude that any model H_{ZY} can be tuned to achieve GHZ-state generation. An explicit demonstration is given in [44] for 21 qubits and $\gamma = 0.7$ (integrating from $\gamma = 0$, using the direct method). The modest choice of size derives only from memory limits of simulating a full Hamiltonian for verification.

III. OPTIMAL CLONING

Production of the GHZ state has shown that although studies of state transfer are often constrained to the single-excitation subspace, the same ideas can be applied to generate interesting evolution in multiple excitation sectors. For the GHZ state, this was a fixed input providing a fixed output. Are there other protocols that we might consider? An arbitrary unitary seems to be out of the question, even within the single-excitation subspace; if we use a state synthesis routine [1,2], then we can choose the evolution of a particular excitation, say, $|1\rangle \rightarrow |\psi\rangle$. But then the possible evolution of other input states is tightly constrained. For example, since $|2\rangle = H|1\rangle/J_1$,

$$e^{-iHt}|2\rangle = \frac{1}{J_1} H e^{-iHt}|1\rangle = \frac{1}{J_1} H |\psi\rangle.$$

So, if $|\psi\rangle$ is confined to a small set of sites, $|2\rangle$ evolves only onto those and neighboring sites. Perfect state transfer demonstrates this: any system that transfers $|1\rangle \rightarrow |N\rangle$ must also transfer $|n\rangle \rightarrow |N + 1 - n\rangle$; there is no freedom to choose these transformations. Still, there may be interesting protocols that depend upon a small input subspace. Perfect state transfer is one such example, wherein the possible inputs are spanned by a basis of two states. Another example is $1 \rightarrow N$ cloning. The optimal $1 \rightarrow N$ universal asymmetric cloning machine [7,26,27] implements the (not necessarily unique) transformation

$$\begin{aligned} |0\rangle &\mapsto A|0\rangle^{\otimes N} + \sum_{n=1}^N \beta_n |1\rangle^{\otimes(n-1)} |0\rangle |1\rangle^{\otimes(N-n)}, \\ |1\rangle &\mapsto A|1\rangle^{\otimes N} + \sum_{n=1}^N \beta_n |0\rangle^{\otimes(n-1)} |1\rangle |0\rangle^{\otimes(N-n)}, \end{aligned} \quad (7)$$

where

$$A = \sum_{n=1}^N \beta_n, \quad B^2 = \sum_{n=1}^N \beta_n^2,$$

and β_n determine the asymmetry of the cloning quality via the single-copy average fidelities,

$$F_n = \frac{2 + (\beta_n + A)^2}{6},$$

and satisfy the normalization $A^2 + B^2 = 1$. We will now show how this can be implemented using spin chains. There is not a single spin chain that achieves this entire transformation, but we can use spin chains as tools that massively simplify the sequence of quantum gates that need to be applied. To that end, consider a set of $M = 2N - 1$ qubits. One qubit, $k + 2$, is the unknown state to be cloned $|\psi\rangle$, but rotated by a phase gate, and we still aim to produce N clones on the odd-numbered qubits. The rest are prepared in the separable state:

$$|0\rangle^{\otimes k} (A|0\rangle + iB|1\rangle) \otimes \sqrt{Z}|\psi\rangle \otimes |0\rangle^{\otimes(M-2-k)}.$$

This can be decomposed into the four basis states

$$|0\rangle^{\otimes k} |x\rangle |0\rangle^{\otimes(M-2-k)}$$

for $x \in \{0, 1\}^2$. We evolve with any M -qubit, GHZ-generating H_{ZY} for the generation time. According to Eq. (5), this

produces the states

$$X_{M-k}^{x_1} X_{M-k-1}^{x_2} |\text{GHZ}\rangle.$$

Now, we apply a controlled-phase gate between the qubits $M - k$ and $M - k - 1$. The effect is that if the two bits of x are the same value, this works like a phase gate on the $|\text{GHZ}\rangle$ state [skipping two steps in Eq. (4)], while it does nothing if the two bit values are different. Then, we repeat the GHZ evolution. Referring to Eq. (4), this returns $X_{k+1}^{x_1} X_{k+2}^{x_2} |0\rangle^{\otimes M}$ if the two bits of x are equal and $X_{k+1}^{x_1} X_{k+2}^{x_2} |1\rangle^{\otimes M}$ otherwise. Overall,

$$\begin{aligned} |0\rangle &\mapsto A|0\rangle^{\otimes M} + B X_{k+1} |1\rangle^{\otimes M}, \\ |1\rangle &\mapsto A X_{k+2} |1\rangle^{\otimes M} + B X_{k+1} X_{k+2} |0\rangle^{\otimes M} \end{aligned}$$

have been implemented (the inputs being the basis of the qubit to be cloned). Finally, we apply a controlled-NOT controlled from qubit $k + 1$ and targeting qubit $k + 2$. This gives the overall transformation

$$\begin{aligned} |0\rangle &\mapsto A|0\rangle^{\otimes M} + B|1\rangle^{\otimes k} |0\rangle |1\rangle^{\otimes(M-1-k)}, \\ |1\rangle &\mapsto A|1\rangle^{\otimes M} + B|0\rangle^{\otimes k} |1\rangle |0\rangle^{\otimes(M-k-1)}. \end{aligned}$$

In fact, the entire transformation up to this point can be implemented by a single Hamiltonian evolution, using a less physically motivated Hamiltonian, based on a tuned three-body cluster-state Hamiltonian [25]. Alternatively, as already observed, one can replace the Ising-generating Hamiltonian with a perfect-state-transfer Hamiltonian by applying a Hadamard transform before and after. This has the advantage of making the form of the Hamiltonian for the GHZ generation and state synthesis parts the same, up to modification of the coupling strengths, at the cost of adding some local Hadamard gates.

From here, we can get our overall desired cloning transformation, creating the N clones on the odd-numbered qubits of the chain, if we can implement

$$\begin{aligned} |0\rangle^{\otimes M} &\mapsto |0\rangle^{\otimes M}, \\ |1\rangle^{\otimes M} &\mapsto |1\rangle^{\otimes M}, \\ |k + 1\rangle &\mapsto \frac{\sum_{n=1}^N \beta_n |2n - 1\rangle}{B}, \\ |\overline{k + 1}\rangle &\mapsto \frac{\sum_{n=1}^N \beta_n |\overline{2n - 1}\rangle}{B}, \end{aligned}$$

where $|0\rangle^{\otimes(k-1)} |1\rangle |0\rangle^{\otimes(M-k)} = |k\rangle$ and $|\overline{k}\rangle = X^{\otimes M} |k\rangle$. The first two transformations are automatic for an exchange-coupled spin chain,

$$H_{XX} = \sum_{n=1}^{M-1} \frac{J_n}{2} (X_n X_{n+1} + Y_n Y_{n+1}), \quad (8)$$

because $|0\rangle^{\otimes M}$ and $|1\rangle^{\otimes M}$ are null vectors of H_{XX} . Assume that couplings can be found to implement the third transformation in a time t_0 [1,2]. This will be discussed in Sec. III A. Indeed, a suitable solution was found in [2] for an initial state in the middle, $k = N - 1$, and is reproduced in Fig. 2. For the last condition, observe that $[H_{XX}, X^{\otimes M}] = 0$. Hence,

$$e^{-iH_{XX}t_0} |\overline{k}\rangle = X^{\otimes M} e^{-iH_{XX}t_0} |k\rangle,$$

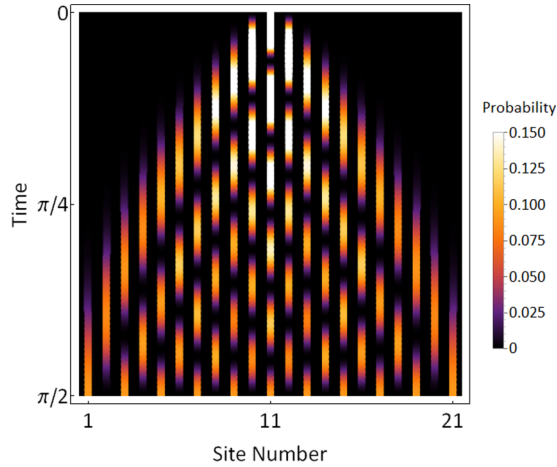


FIG. 2. A single excitation input to the central spin of a 21-qubit spin chain (top) evolves into a uniform superposition over the odd-numbered sites (bottom), as required for optimal symmetric universal $1 \rightarrow 11$ cloning.

which simply yields that the transformations in the 1 and $M - 1$ excitation subspaces are essentially identical, so the final condition will also be satisfied.

All of the complexity of producing these clones is conveniently wrapped up in just two helper functions. The corresponding circuit diagrams are contrasted in Figs. 3 and 4. While a quantum circuit for cloning has previously been explicitly stated for small sizes [45], we are not aware of a version, other than probabilistic versions [27], that works deterministically for general $1 \rightarrow N$ universal symmetric cloning, let alone the asymmetric case. This is probably because the cloning map in Eq. (7), specialized to symmetric cloning, is not the map usually stated [46,47]: [7] reveals that the cloning map is associated with the ground state of a particular matrix, and the symmetric case is highly degenerate. The version that we have chosen, Eq. (7), extends consistently from the asymmetric case and lends itself well to implementation with a quantum circuit, as depicted in Fig. 4, providing such a definition for symmetric cloning. This circuit can be modified for asymmetric cloning. Assuming that the architecture exhibits only nearest-neighbor

couplings (the only instance where it makes sense to consider implementation via a nearest-neighbor Hamiltonian), the depth of the circuit is N and comprises $O(N^2)$ gates.

Why do we create clones only on every second site of the chain? Imagine we have a Hamiltonian like Eq. (8) but include magnetic fields as well:

$$H = \sum_{n=1}^{M-1} \frac{J_n}{2} (X_n X_{n+1} + Y_n Y_{n+1}) + \sum_{n=1}^N \frac{B_n}{2} (\mathbb{1} - Z_n). \quad (9)$$

Let

$$U = \prod_{n=1}^{(N+1)/2} X_{2n-1} \prod_{n=1}^{(N-1)/2} Y_{2n}.$$

We have that $UHU = -H + (\sum_{n=1}^N B_n)\mathbb{1}$. Moreover, at t_0 , $e^{-iHt_0} = e^{iHt_0}$ because $e^{-i\lambda_n t_0} = \pm 1$ for every eigenvalue (neglecting, for simplicity, a possible global phase). Thus, time evolution in the higher-excitation subspace is given by

$$e^{-iHt_0} |\bar{n}\rangle = -(-i)^{(N-1)/2+2n} U e^{-i[H+(\sum_{n=1}^N B_n)\mathbb{1}]t_0} |n\rangle.$$

If the state produced in the single-excitation subspace is

$$|\psi\rangle = \sum_n \alpha_n |n\rangle,$$

in the higher-excitation subspace we get

$$\begin{aligned} & (-i)^{(N+3)/2+2n} e^{-it_0 \sum_m B_m} U \sum_m \alpha_m |m\rangle \\ &= e^{-it_0 \sum_m B_m} \sum_m (-1)^{n+m} \alpha_m |\bar{m}\rangle. \end{aligned}$$

This is consistent with the desired transform only if $\alpha_m = 0$ on every second site to eliminate the effect of the $(-1)^{n+m}$ term. In doing so, it transpires that for a symmetric target spectrum, it suffices to set $B_n = 0$. One can readily verify that although the even-numbered qubits effectively act as ancillas and although the transformation we implement does not leave them separable, it does not adversely affect the cloning fidelity.

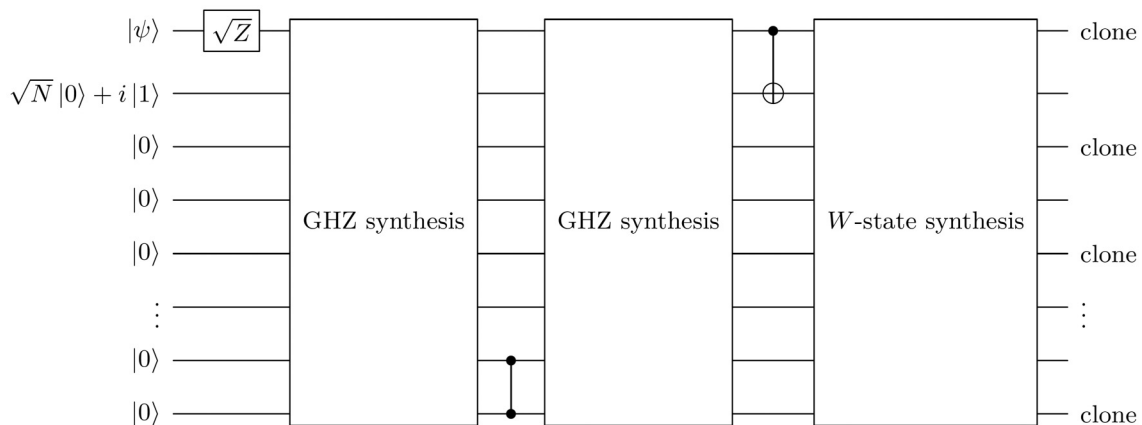


FIG. 3. Quantum circuit diagram for quantum cloning when supplemented by two helper Hamiltonians, acting on $M = 2N - 1$ input qubits. Coefficients are specifically chosen for optimal universal symmetric cloning. The input qubit can be arbitrary.

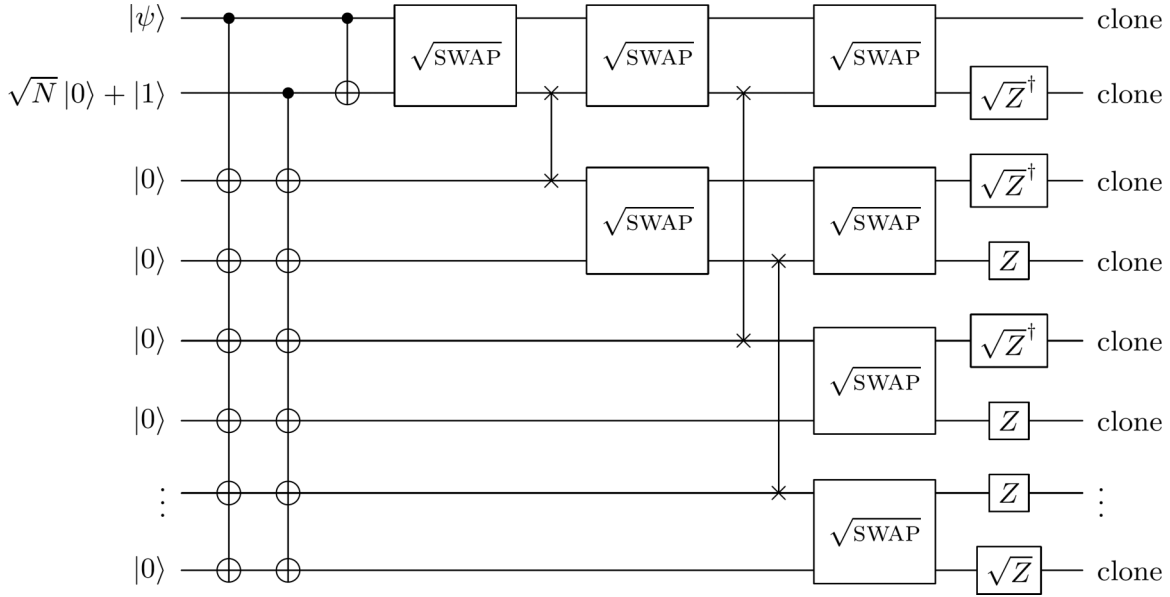


FIG. 4. Quantum circuit diagram for quantum cloning without the two helper Hamiltonians. Moving beyond the symmetric cloner requires replacement of the $\sqrt{\text{SWAP}}$ with partial SWAP operations of different amounts.

A. Designing state synthesis systems

While we can use the algorithms of [1,2] to generate figures such as Fig. 2, they are limited by very small radii of convergence. Instead, we would now like to examine if the isospectral flow ideas outlined above can be applied. The two papers [1,2] provide different philosophies for how to produce a useful chain, but for our purposes, much of the calculation in the same.

Let us start with a first guess at a Hamiltonian H_0 (in the single-excitation subspace). It has a desired spectrum, but its evolution produces

$$e^{-iH_0t_0}|\phi\rangle = |\psi_0\rangle,$$

starting from the separable state $|\phi\rangle$ (i.e., a single excitation on a particular site) and evolving for a time t_0 , where $|\psi_0\rangle$ is not our target state $|\psi_t\rangle$. In practice, we will set $|\phi\rangle = |\frac{N+1}{2}\rangle$ to minimize the evolution time. However, if $|\psi_t\rangle$ is symmetric, we can reduce the task to finding a chain of half the length, starting with an excitation at site 1 [2]. For that reason, we will typically assume $|\phi\rangle = |1\rangle$.

How can we make a better guess, H_1 , which should have the same spectrum as H_0 ? Again, we use the isospectral transformation

$$H_1 = e^{-B} H_0 e^A$$

for some anti-Hermitian matrices A, B . We have several properties that we want to impose. As in the Ising case, we can rearrange the matrix H_0 into a block structure of {all odd elements, all even elements}, so that H_0 takes the form

$$H_0 = \begin{pmatrix} 0 & X_0 \\ X_0^T & 0 \end{pmatrix}.$$

A block diagonal $A = \text{diag}(A_o, A_e)$ preserves the structure of H_0 , with X_0 evolving as $X_0 \rightarrow e^{-A_e} X_0 e^{A_o}$.

Next, we want to impose that the tridiagonal structure of H_0 is preserved. This just requires $\langle i|\dot{H}_0|j\rangle = 0$ for all $|i - j| > 1$.

We have already partially achieved this with our block diagonal of A ensuring that it is true for all $|i - j|$ being even. The remainder is simply a set of simultaneous linear equations in the parameters A_o and A_e :

$$\langle i|X_0A_o - A_eX_0|j\rangle = 0 \quad \forall j \neq i, i - 1.$$

We are hence building up a set of linear conditions which, so far, just ensure that subsequent matrices maintain the important properties of the initial matrices. Now we must impose that each subsequent iteration moves us towards a better evolution. Since we are updating

$$H_0 \mapsto e^{-A} H_0 e^A$$

and assuming A is small, the evolution updates as

$$e^{-A} e^{-iH_0t_0} e^A |\phi\rangle \approx |\psi_0\rangle - A|\psi_0\rangle + e^{-iH_0t_0} A|\phi\rangle.$$

To fix the new evolution to be $|\psi_t\rangle$, we might solve

$$-A|\psi_0\rangle + e^{-iH_0t_0} A|1\rangle = |\psi_t\rangle - |\psi_0\rangle,$$

subject to the structure constraints that we have already described. However, a solution of this form is unlikely to keep A small; it is perhaps better to describe it as a constrained optimization problem (linear programming):

$$\max_{A: \|A\| \leq \delta} (\langle \psi_t | - \langle \psi_0 | [e^{-iH_0t_0}, A] | \phi \rangle).$$

Having found A , we update H_0 according to the unitary transformation update (rather than direct integration).

For the particular design philosophy of [1], we can go further. There, due to the chosen spectrum,

$$e^{-iH_0t_0} = \mathbb{1} - 2|\lambda_0\rangle\langle\lambda_0|,$$

the aim is to fix the null vector to $|\lambda_t\rangle$, where

$$|\psi_t\rangle = (\mathbb{1} - 2|\lambda_t\rangle\langle\lambda_t|)|1\rangle,$$

knowing that the null vector of the next iteration is $e^{-A}|\lambda_0\rangle$. In successive iterations, we aim to maximize the overlap with

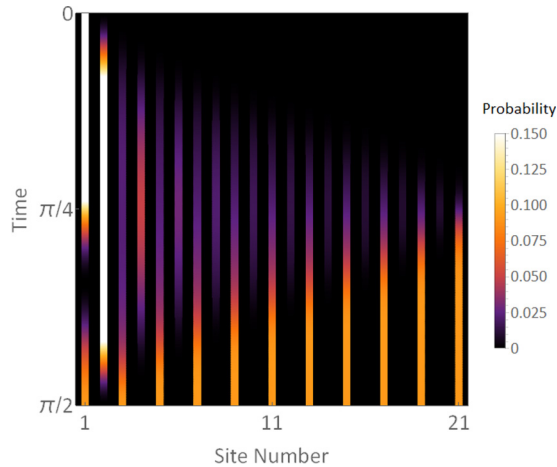


FIG. 5. A single excitation input to the first spin of a 21-qubit spin chain (top) can be caused to evolve into a uniform superposition over the odd-numbered sites (bottom), as required for symmetric $1 \rightarrow 11$ cloning.

$|\lambda_t\rangle$. Hence, we have to solve the linear programming problem

$$\min_{\|A\|_\infty \leq \delta} \langle \lambda_t | A | \lambda_0 \rangle.$$

Convergence is well motivated: unless there is a reason that either $|\lambda_0\rangle$ or $|\lambda_t\rangle$ must be a null vector of A , we can always find a nonzero value of the overlap, and choosing the sign of A ensures that the outcome is always negative and hence iterates towards an improved solution. The solution must converge, and further justification that it converges globally on the correct solution is given in the Appendix. A typical output is shown in Fig. 5 and is used in [44] to demonstrate the proper functioning of the entire evolution sequence, giving the optimal symmetric cloning fidelity of $F = \frac{23}{33}$.

1. Hypercubes

So far, we have discussed engineering a one-dimensional chain, choosing the coupling strengths to achieve the evolution that we desire. This perfectly parallels studies of perfect state transfer. However, another avenue for study in perfect state transfer is the set of uniformly coupled graphs that can accomplish the same task. For example, hypercubes of arbitrary dimension k and side length 2 or 3 achieve perfect state transfer at distance k or $2k$, respectively [18]. Various other graphs have since been shown to provide perfect transfer, including a variety of modifications of the hypercubes [18,48–50]. Can graphs also be used for the state synthesis tasks that could be useful for quantum cloning? We specifically focus on generating uniform superpositions across some subset of sites, with a preference for those where the phase on each of the superposed sites is the same.

Let G be a graph with edges E and N vertices V . The graph Hamiltonian is defined as

$$H_G = \frac{1}{2} \sum_{(i,j) \in E} (X_i X_j + Y_i Y_j).$$

Like before, there is a subspace structure based on the number of excitations, $|0\rangle^{\otimes N}$ and $|1\rangle^{\otimes N}$ are null vectors, and

$[H, X^{\otimes N}] = 0$; we have to get only the evolution in the single-excitation subspace correct. If the graph is bipartite, the phase choice can be consistent only if the superposed sites are all part of the same bipartition by a generalization of the argument around Eq. (9).

In the single-excitation subspace, the Hamiltonian is represented by the adjacency matrix A of the graph. The conditions on state synthesis [2] are remarkably similar to those of perfect transfer [19,51]: to start from a site n , producing a state $|\psi\rangle$ in time t_0 , if $|\psi\rangle \in \mathbb{R}^N$, then

$$\langle \lambda_m | n \rangle = \pm \langle \lambda_m | \psi \rangle$$

for every eigenvector $|\lambda_m\rangle$, and for those eigenvectors for which $\langle \lambda_m | n \rangle \neq 0$, the eigenvalues λ_m must satisfy

$$e^{-i\lambda_m t_0} = \pm e^{i\phi}$$

for some phase ϕ . This has some further consequences for the spectrum of the adjacency matrix A [52]. For example, with one extra assumption about the nature of the state synthesis task (that all vertices have a perfect revival at the same time), we know that the spectrum for a nonbipartite graph must be integral, while for a bipartite graph, the spectrum is either integral or rational multiples of $\sqrt{\Delta}$ for a square-free integer Δ . We will not develop this theory more generally here but will focus on some special cases that we have found.

Several instances of the path P_n (i.e., a uniformly coupled chain of n vertices) generate uniform superpositions:

$$\begin{aligned} e^{-iA(P_2)\pi/4} |1\rangle &= \frac{|1\rangle - i|2\rangle}{\sqrt{2}}, \\ e^{-iA(P_3)\pi/\sqrt{8}} |2\rangle &= \frac{|1\rangle + |3\rangle}{\sqrt{2}}, \\ e^{-iA(P_3)\cos^{-1}(\frac{1}{\sqrt{3}})/(\sqrt{2}\pi)} |2\rangle &= \frac{|1\rangle + i|2\rangle + |3\rangle}{\sqrt{3}}, \\ e^{-i2A(P_5)/\sqrt{27}} |3\rangle &= \frac{|1\rangle + i|2\rangle + i|4\rangle + |5\rangle}{2}. \end{aligned}$$

The second of these is ideally suited to $1 \rightarrow 2$ symmetric cloning (and is closely related to a previous construction [53,54]). We can extend these cases by using the hypercube construction [18]. For a graph G , the adjacency matrix of the corresponding k -dimensional hypercube is

$$A(G^k) = \sum_{n=1}^k \mathbb{1}_N^{\otimes(n-1)} \otimes A \otimes \mathbb{1}_N^{\otimes(k-n)}.$$

This describes independent evolution along each of the k dimensions. Taking a basis $x \in [N]^k$ (a k -dimensional vector where each element takes an integer value from 1 to N), each $|x\rangle$ corresponds to a single excitation being on a particular vertex of the graph. Starting from $|y\rangle$, the final amplitude on a vertex $|x\rangle$ is

$$\prod_{i=1}^k \langle x_i | e^{-iA t} | y_i \rangle.$$

If A gives a uniform superposition, so does the hypercube. The hypercube with a side length of 2 (i.e., derived from P_3) is particularly compelling. For example, a 3×3 square lattice of uniformly coupled qubits generates a uniform superposition

of all nine sites. Or, more applicable to symmetric cloning, P_3^k produces (at a different time) the uniform superposition across all 2^k corners.

IV. CONCLUSIONS

We have shown how a fixed transverse Ising system can produce a GHZ state, which is a key quantum resource for use in future technologies. This could act as a stand-alone device or as a special unit, a coprocessor, within a larger quantum device. The fixed-function coprocessor replaces what would otherwise be a complex sequence of unitary gates, with the inaccuracies inherent in the multiple separate steps that have to be taken in its implementation. By extending the results from H_I to H_{ZY} , we have potentially opened GHZ synthesis to a much broader range of experiments. Realizing that the transformation required to achieve GHZ-state synthesis in a transverse Ising model essentially reduces to a state-transfer-like condition on h_1 would also significantly simplify optimal control studies such as [55], moving away from the assumption of perfect engineering.

We have also specified a second transformation. This two-step procedure, where the first step uses the GHZ synthesizer, implements optimal asymmetric universal cloning of qubits. A nonprobabilistic strategy has been given for these cloning machines. Our transformation is implemented by operations that are local in a one-dimensional chain of qubits. Our use of the GHZ synthesizer to implement a single-control, multiple-target controlled-NOT gate may be of further interest in relation to the Fourier transform.

Both transformations, when restricted to a nearest-neighbor architecture, exhibit an optimal $O(N)$ scaling in run time and have running times essentially identical to their quantum circuit equivalents.

Central to these results was an isospectral transformation algorithm with fine-grained control over directing consecutive iterations. Global convergence of the algorithm is well motivated. *Mathematica* scripts that implement the reported results for chains of 21 qubits are available from [44]. The algorithm demonstrates considerable potential for further development and should be broadly applicable.

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APPENDIX: CONVERGENCE OF THE ISOSPECTRAL ALGORITHM

We aim to motivate that the algorithm described in Sec. III A has a single point of convergence provided the target null vector $|\lambda_t\rangle$ can be a null vector of a system with the fixed spectrum.

At each step, we iterate with a matrix A and impose that the elements of A are bounded by some step size δ , $\|A\|_\infty \leq \delta$. It is important that we pick δ such that the second-order term in the expansion of $\langle \lambda_t | e^{-A} | \lambda_0 \rangle$ is negligible. Since $\langle \psi | A | \psi \rangle = 0$ for all real-valued states $|\psi\rangle$, the maximum value of $\langle \lambda_t | A | \lambda_0 \rangle$

is achieved with

$$A|\lambda_0\rangle \sim |\lambda_t\rangle - \langle \lambda_0 | \lambda_t \rangle |\lambda_0\rangle.$$

This is generically possible to fix: $A|\lambda_0\rangle$ has $(N + 1)/2$ components, and the $(N - 1)/2$ free parameters can control the output in the space orthogonal to $|\lambda_0\rangle$. By aligning these vectors with the left- and right-maximum singular vectors of A (singular value σ), we have $\langle \lambda_t | A | \lambda_0 \rangle = -\sigma \sqrt{1 - \langle \lambda_0 | \lambda_t \rangle^2}$. Meanwhile, $\langle \lambda_t | A^2 | \lambda_0 \rangle < \sigma^2$, i.e., $\sigma \ll \sqrt{1 - \langle \lambda_0 | \lambda_t \rangle^2}$. Through the following bounds, we relate δ to σ :

$$\sqrt{\frac{N+1}{2}}\delta \geq \sqrt{\frac{N+1}{2}}\|A\|_\infty \geq \|A\|_2 \geq \sigma,$$

where $\|A\|_2$ is the Frobenius norm. Thus, by ensuring that $\delta = \epsilon \sqrt{1 - \langle \lambda_0 | \lambda_t \rangle^2}$ for small $\epsilon > 0$, the second-order term is always negligible. Indeed, we can directly bound the value $\chi = \langle \lambda_t | \lambda_0 \rangle$:

$$\frac{d\chi}{dt} \leq -\sigma \sqrt{1 - \chi^2} \leq -\epsilon(1 - \chi^2).$$

This solution tends exponentially towards $\chi = 1$. Assuming that $A|\lambda_0\rangle$ can be picked as specified, we have convergence, and accuracy ϵ is achieved with an average complexity of $O(N^6 \log(\sqrt{N}/\epsilon))$, the leading term arising from solving $O(N^2)$ linear constraints.

Generically, $A|\lambda_0\rangle$ can be any state in the odd space that is orthogonal to $|\lambda_0\rangle$ since there are $(N - 1)/2$ free parameters with which to achieve this. So, when does this fail? Is this compatible with the observation that some states cannot be the null vector for a tridiagonal system of a particular spectrum [1]?

We start by noting that although [1] indicated that none of the vector elements on the odd space can be zero (as this would give consecutive zero elements on the complete vector), this was an artificial imposition resulting from requiring nonzero coupling strengths. However, this consideration is not built into the algorithm, so we are not prevented from reaching these forms of $|\lambda_t\rangle$.

Since the state $A|\lambda_0\rangle$ is linear in the free parameters, the space described when the corresponding vectors do not span the space (and are hence linearly dependent on each other) must be a convex space. There is a single inaccessible region, which must therefore include any inaccessible null vectors. The only question is whether this region is tight with that of the inaccessible null vectors.

Case study

In the absence of universal answers, we investigate the special case of $N = 5$ and spectrum $0, \pm 3, \pm 5$ since this is a case where there are forbidden null vectors [1]. The odd space is dimension 3, and there are two free parameters, a and b . For a particular h with coupling strengths J_1, \dots, J_4 ,

$$A_o|\lambda_0\rangle = a \begin{pmatrix} J_1 J_2 (J_3^2 + J_4^2) \\ J_1^2 J_3^2 - J_2^2 J_4^2 \\ J_3 J_4 (J_1^2 + J_2^2) \end{pmatrix} + b \begin{pmatrix} J_1 (J_3^2 + J_4^2 - J_1^2) \\ -J_2 (J_1^2 - J_4^2) \\ -J_2 J_3 J_4 \end{pmatrix}.$$

The ratios $\gamma_1 = \frac{J_1}{J_2}$ and $\gamma_2 = \frac{J_4}{J_3}$ parametrize the possible null vectors

$$|\lambda_0\rangle = \begin{pmatrix} J_2 J_4 \\ -J_1 J_4 \\ J_1 J_3 \end{pmatrix},$$

from which we can derive that the only time that we do not have access to the whole space is when $J_1^2 + J_2^2 = J_3^2 + J_4^2$, and thus,

$$(1 + \gamma_1^2)(1 + \gamma_2^2) = \frac{17^2}{8^2}, \quad (\text{A1})$$

where we used the eigenvalue relations (such as $J_1^2 + J_2^2 + J_3^2 + J_4^2 = 34$) to eliminate the remaining terms. This defines a barrier in the possible space of $|\lambda_t\rangle$ that the algorithm cannot cross.

Now let us consider the region of $|\lambda_t\rangle$ for which there is no h with the correct spectrum and that null vector. Using the

explicit eigenvalue relations for the coupling strengths, we can write that

$$[J_2^2(1 + \gamma_1^2) - 17]^2 = \left\{ 64 - \frac{J_2^2}{1 + \gamma_2^2} [34 - J_2^2(1 + \gamma_1^2)] \right\}.$$

This has a non-negative solution for J_2^2 when

$$(1 + \gamma_1^2)(1 + \gamma_2^2) \geq \frac{17^2}{8^2}.$$

We conclude that our algorithm is capable of converging on any valid null vector for $N = 5$. If we demand an invalid null vector, the algorithm will converge somewhere on the surface of closest approach, defined by Eq. (A1). It is reasonable to expect similar behavior in larger spaces, but this remains unproven.

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