Tracking an open quantum system using a finite state machine: Stability analysis

R. I. Karasik and H. M. Wiseman*

Centre for Quantum Computation and Communication Technology (Australian Research Council), Centre for Quantum Dynamics, Griffith University, Brisbane, Queensland 4111, Australia

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A finite-dimensional Markovian open quantum system will undergo quantum jumps between pure states, if we can monitor the bath to which it is coupled with sufficient precision. In general these jumps, plus the between-jump evolution, create a trajectory which passes through infinitely many different pure states, even for ergodic systems. However, as shown recently by us [Phys. Rev. Lett. **106**, 020406 (2011)], it is possible to construct *adaptive* monitorings which restrict the system to jumping between a finite number of states. That is, it is possible to track the system using a *finite state machine* as the apparatus. In this paper we consider the question of the stability of these monitoring schemes. Restricting to cyclic jumps for a qubit, we give a strong analytical argument that these schemes are always stable and supporting analytical and numerical evidence for the example of resonance fluorescence. This example also enables us to explore a range of behaviors in the evolution of individual trajectories, for several different monitoring schemes.

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I. INTRODUCTION

An open system is one which continuously exchanges information with its environment [1-3]. For a Markovian (memoryless) system, this amounts to a loss of information from the system into the environment. By *monitoring* the environment, it is possible to regain this lost information and hence to *track* the system. If the monitoring process is perfect, then one can expect to perfectly track the system, that is, to know as much as is possible to know about the system, as would be the case were it a closed system prepared at will.

These very general considerations apply equally to classical and quantum systems. However, there are some very significant differences between the two cases. In the classical case, there is only one (best) way to monitor the environment. Also, if the classical system has only finitely many possible states this is known as a finite state machine [4]—then obviously it is possible to keep track of its state using only a finite classical memory—another finite state machine of the same dimension. In the quantum case, by contrast, there are infinitely many inequivalent ways of monitoring the environment that enable the experimenter to deduce what pure state the system is in [1,3,5]. This is because of the entanglement between the system and its environment. But in almost all cases, an infinite classical memory is required to store that pure state, even for a finite-dimensional quantum system.

This last point can perhaps only be understood by introducing a little formalism. We consider finite-dimensional systems that undergo Markovian open quantum system dynamics, described by a Lindblad-form master equation [2,3]:

$$\dot{\rho} = \mathcal{L}\rho \equiv -i[\hat{H}_{\text{eff}}\rho - \rho\hat{H}_{\text{eff}}^{\dagger}] + \sum_{l=1}^{L}\hat{c}_{l}\rho\hat{c}_{l}^{\dagger}, \qquad (1)$$

where $\hat{H}_{\text{eff}} \equiv \hat{H} - i \sum_{l} \hat{c}_{l}^{\dagger} \hat{c}_{l}/2$. Here \hat{H} is Hermitian (it is the Hamiltonian) but the jump operators $\{\hat{c}_{l} : l\}$ are completely arbitrary. This decomposition defines the evolution of the

system conditioned on a perfect monitoring of the bath quanta. If a quantum in channel *l* is observed at time t_n , the system state jumps from the prejump state $|\psi(t_n^-)\rangle$ to the postjump state $|\psi(t_n)\rangle \propto \hat{c}_l |\psi(t_n^-)\rangle$. Then until the next jump its subsequent evolution would be generated by the effective (non-Hermitian) Hamiltonian \hat{H}_{eff} appearing in Eq. (1).

In general, the postjump state will depend on the prejump state $|\psi(t_n^-)\rangle$, and it will not remain stationary until the next jump, unless it happens to be an eigenstate of \hat{H}_{eff} . It is thus not at all obvious whether for a general open quantum system it would be possible to keep track of its pure state, even in principle, with a finite classical memory. On the face of it, it would seem necessary to store the nature and exact times of each jump—requiring a sequence of real numbers $\{t_n : n\}$ each of which would require, in principle, an infinite memory to store. Alternately one could store the conditioned quantum state $|\psi(t)\rangle$ itself, but this (a finite vector of complex numbers) would also require an infinite memory.

In Ref. [6] we showed that, for an arbitrary *D*-dimensional quantum system obeying a Markovian master equation with a unique stationary (mixed) state, one can expect there to exist a monitoring such that it is possible to track the system's conditional (pure) state with a *K*-state machine as apparatus, with some $K \ge (D-1)^2 + 1$. For a qubit we proved that this is indeed always the case: a two-state apparatus can be found that ensures the qubit jumps between just two states. This apparatus must implement an *adaptive* monitoring, choosing how to measure the environment depending on its own internal state. This result shows that the infinite amount of information required to track a quantum system under a generic monitoring scheme arises from the poor choice of scheme and is not intrinsic to the coupling between the quantum system and its environment.

The goal of this paper is to investigate the stability of such finite-state monitoring and provide details on how such schemes can be identified and constructed. The first goal is necessary to establish that the schemes introduced in Ref. [6] are not just mathematical constructions, but are physically realizable. For this purpose we restrict to qubit evolution and cyclic-jump schemes. All the schemes we study are stable in a

^{*}h.wiseman@griffith.edu.au

mean-square sense, and we give a strong analytical argument plus supporting numerical evidence (for two- and three-state machines) that such schemes are always stable.

The paper is organized as follows. In Sec. II we review the background for this paper and summarize relevant results from Ref. [6]. In Sec. III we develop our main results for stability. We discuss these results in the context of specific examples of resonance fluorescence in Sec. IV. We conclude with a summary and a statement of open problems in this area.

II. THE PREFERRED ENSEMBLE

The idea of tracking the system with finite classical memory is closely related to the concept of a preferred ensemble. To explain the latter notion, we consider an evolution generated by a Markovian master equation with a unique steady state defined by $\mathcal{L}\rho_{ss} = 0$ and we assume that ρ_{ss} is a mixed state. This mixed state can be decomposed in terms of an ensemble of pure states $|v_k^e\rangle$ with some positive weights \wp_k such that

$$\rho_{\rm ss} = \sum_{k=1}^{K} \wp_k \left| v_k^e \right\rangle \! \left\langle v_k^e \right|. \tag{2}$$

Note that there are infinitely many such decompositions as the states $|v_k^e\rangle$ need not be orthogonal. A valid interpretation for any such decomposition of a mixed state is that in the long time limit a measurement performed on the environment will collapse the system into one of the possible states $|v_k^e\rangle$ and this will happen with probability \wp_k . Such a measurement will, in general, require simultaneously measuring parts of the environment that interacted with the system at *different* times in the past.

A different question arises if we require the measurements on the environment to be continuous in time, so that the system state is continually being collapsed, and so (in the long time limit) will be in a stochastically evolving pure state. Because the system evolution is Markovian, this can be done while leaving unchanged the average evolution of the system (1). Now the natural interpretation for the ensemble $\{\wp_k, |v_k^e\rangle\}$ is that the system will only ever be in one of the states $|v_k^e\rangle$ and will spend a proportion of time in that state equal to \wp_k . In this case, this interpretation is *not* valid for most ensembles. Decompositions $\{\wp_k, |v_k^e\rangle\}$ that can be realized through an experiment via continuous measurement are called *physically* realizable (PR). The fact that some ensembles are not PR is known as the preferred ensemble fact [7]. We note here that if we know of a PR ensemble with K finite, then if at some time t the system is a pure state from the PR ensemble, $|v_k^e\rangle$, and is subject to the continuous monitoring that realizes this PR ensemble, then in its subsequent conditional evolution the system will only occupy states from the PR ensemble and we can therefore track such evolution with a classical register with only K states. Such a classical device is known as a finite state machine.

From Ref. [7] we know that the ensemble $\{\wp_k, |v_k^e\rangle\}$ is PR if and only if (iff) there exists rates $\kappa_{jk} > 0$ such that

$$\forall k, \mathcal{L} | v_k^e \rangle \langle v_k^e | = \sum_{k=1}^K \kappa_{jk} (| v_j^e \rangle \langle v_j^e | - | v_k^e \rangle \langle v_k^e |), \qquad (3)$$

where the system jumps between *K* different states. Typically, most ensembles $\{\wp_k, |v_k^e\rangle\}$ representing ρ_{ss} are not PR, including the K = D ensemble composed from the diagonal basis for ρ_{ss} [3].

For a qubit, the conditions in Eq. (3) can be simplified further. Using the Bloch representation, Eq. (1) becomes

$$\vec{r} = A\vec{r} + \vec{b},\tag{4}$$

where A is a 3×3 matrix that dictates the evolution of the state and \vec{b} , a three-vector, determines the steady state of the system, $\vec{r}_{ss} = -A^{-1}\vec{b}$. This equation has a unique steady state iff (if and only if) the real parts of all eigenvalues of A are negative. We can track this system with a K-state memory iff there exists an ensemble $\{\wp_k, \vec{r}_k\}$ and rates $\kappa_{ik} \ge 0$ such that

$$\vec{r}_k \cdot \vec{r}_k = 1 \ \forall k, \tag{5}$$

$$A\vec{r}_j + \vec{b} = \sum_{k=1}^K \kappa_{jk} (\vec{r}_k - \vec{r}_j) \;\forall j.$$
(6)

These equations are better than conditions from Eq. (3) for numerical search for PR ensembles because they generically reduce the number of different equations, the number of different variables, and the maximal degree of this system of equations. Thus the search for PR ensembles reduces to finding a real solution to a system of quadratic polynomials with real coefficients. Unfortunately, this is still a hard problem, which in general is known to be NP complete [8]. Here NP stands for nondeterministic polynomial time algorithm. Equations (5) and (6) are used in Sec. IV C to analyze a PR ensemble for K = 3 for a qubit.

For the moment, we concentrate on a K = 2 PR ensemble. In Ref. [6], we showed that such an ensemble always exists for a qubit. This can be easily deduced from Eqs. (5) and (6), which for K = 2 imply that $A(\vec{r}_1 - \vec{r}_2) = (\kappa_{12} + \kappa_{21})(\vec{r}_2 - \vec{r}_1)$. This is an eigenvalue equation and we can conclude that

$$\vec{r}_1 = \vec{r}_{\rm ss} + \eta_1 \hat{u},\tag{7}$$

$$\vec{r}_2 = \vec{r}_{\rm ss} - \eta_2 \hat{u},\tag{8}$$

where \hat{u} is the normalized eigenvector of matrix A and parameters η_1 and η_2 relate to probabilities \wp_1 and \wp_2 for occupying states \vec{r}_1 and \vec{r}_2 , respectively, and $\wp_j = \eta_j/(\eta_1 + \eta_2)$. Detailed expressions for η_1 and η_2 can be found in Ref. [6], but are not relevant to this paper. Here we just need the general structure of the solutions. Because the Bloch vectors, \vec{r}_1 and \vec{r}_2 , must be real, only real eigenvectors \hat{u} of A can contribute to the solution. As a 3 × 3 matrix, A has three eigenvalues and, by a fundamental theorem of algebra, at least one eigenvalue (and consequently one eigenvector) is real. Therefore, a qubit always has a preferred ensemble comprising just two states and there can be up to three different solutions.

Different PR ensembles (including three different PR ensembles to track a qubit with one bit [6]) arise from the freedom that experimentalists have (in principle) to monitor the system's environment in different ways. This freedom exists because of the invariance properties of the master equation,

Eq. (1), which is invariant with respect to transformations

$$\{\hat{c}_l\} \to \left\{ \hat{c}'_m = \sum_{l=1}^L S_{ml} \hat{c}_l + \mu_m \right\},\tag{9}$$

$$\hat{H} \to \hat{H}' = \hat{H} - \frac{i}{2} \sum_{m=1}^{M} (\mu_m^* \hat{c}_m' - \mu_m \hat{c}_m'^{\dagger}),$$
 (10)

where $\vec{\mu} = (\mu_1, \dots, \mu_M)$ is an arbitrary complex vector and **S** is an arbitrary semiunitary matrix— $\sum_{m=1}^{M} S_{l'm}^{*} S_{ml} = \delta_{l',l}$. Realizations of this master equation with $\{\hat{c}'_m\}$ as the jump operators and $\hat{H}'_{\text{eff}} = \hat{H}' - i \sum_{m=1}^{M} \hat{c}'_m^{\dagger} \hat{c}'_m/2$ as the effective non-Hermitian Hamiltonian generate the same average evolution as the original master equation, but clearly give rise to different stochastic evolution. To obtain the most general pure-state unraveling of the master equation, we require $\vec{\mu}$ and **S** to depend upon the previous record of jumps. That is, we require an adaptive monitoring. Of course when we use this to achieve jumping between a finite number of states, the classical *K*-state memory that stores which state the system is currently in carries all the information necessary for determining $\vec{\mu}$ and **S**. That is, the adaptive unraveling is specified by *k* different values for $\vec{\mu}$ and **S**.

The physical meaning of these parameters is most easily explained in a quantum optics context. The matrix **S** describes a linear interferometer taking the field outputs from the system as inputs and interfering them prior to detection. This is to be understood in the most general sense, including frequency shifters if the system has outputs in different frequency bands. The vector $\vec{\mu}$ describes adding (weak) local oscillators to the output fields from the interferometer prior to detection by photon counting. There have been for many years theoretical proposals for adaptively controlling the local oscillator amplitude [9] or phase [5,10], and more recently a number of experiments have been performed [11–13], one of which (Ref. [12]) used a weak local oscillator, with amplitude comparable to that of the system.

One characteristic that sets apart different solutions is the Shannon entropy. Under continuous monitoring, the system will occupy states $|v_k^e\rangle$ with probabilities \wp_k . The Shannon entropy for an ensemble $\{\wp_k, |v_k^e\rangle\langle v_k^e|\}$ that represents ρ_{ss} is $h(\{\wp_k\}) = -\sum_k \wp_k \log_2 \wp_k$. This is lower bounded by the von Neumann entropy of ρ_{ss} :

$$h(\{\wp_k\}) \ge S(\rho_{\rm ss}) = -\operatorname{Tr}[\rho_{\rm ss}\log_2 \rho_{\rm ss}],\tag{11}$$

where equality holds iff $\{\wp_k, |v_k^e\rangle \langle v_k^e|\}$ is the diagonal ensemble. We showed in Ref. [6] that some of the K = 2 and K = 3ensembles for a qubit can have entropy h much smaller than 1. In this case one can track the state of the qubit with less than one bit on average, meaning that the state of N identical qubits subject to the same independent monitoring can be tracked with $Nh \ll N$ bits. This is one reason why studying different PR ensembles for the same system is of great interest. In the next section we explore the stability of such adaptive monitoring.

III. JUMP DYNAMICS AND STABILITY

As explained above, the existence of a PR ensemble $\{\wp_k, |v_k^e\rangle\langle v_k^e|\}$ ensures that if we start the system in state $|v_k^e\rangle$

and subject it to the adaptive monitoring determined by the parameters κ_{jk} from Eq. (3), then the system will always jump between states $|v_k^e\rangle$ with k = 1, ..., K and throughout its evolution will never leave the PR ensemble. But what happens if the system is not initialized perfectly at the start of the monitoring procedure?

We address this question in this section. We consider a general qubit system, whose evolution is governed by the master equation (1), with one jump operator, \hat{c} . We assume that this system is subject to adaptive monitoring that allows the system to jump between K states in a cyclic manner. Since the resulting evolution is stochastic, we cannot say what will happen to each specific trajectory, but we can determine what happens to many different realizations on average. We show that on average any initial state subject to the adaptive monitoring will eventually converge to the jumping between states $|v_k^e\rangle$ with $k = 1, \ldots, K$ from the PR ensemble. We first develop our results assuming that adaptive monitoring leaves the qubit jumping between K = 2 states and then generalize to K-state cyclic jumps.

A. Two-state jumping

For the two-state jumping scenario, the system jumps between pure states $|v_1^e\rangle$ and $|v_2^e\rangle$. In principle, we can compute these states using Eqs. (7) and (8). This approach tends to yield useful numerical answers, but extracting a simple analytical expression is not easy. An alternative approach for identifying jumping states (for the specific example of resonance fluorescence) was undertaken in Refs. [14,15]. This approach uses the properties of transformations, Eqs. (9) and (10), that leave the master equation (1), invariant to explicitly construct the adaptive monitoring schemes that generates a PR ensemble with the desired number of elements. This is the approach we use in this section.

The system we study in this paper is simple. It has only one jump operator and, therefore, the only degree of freedom for generating Eqs. (9) and (10) is the strength of the local oscillator, a complex scalar μ . For a given μ , the effective Hamiltonian for the system is [3]

$$\hat{H}(\mu) \equiv \hat{H} - \frac{i}{2}\hat{c}^{\dagger}\hat{c} - i\mu^{*}\hat{c} - i\frac{|\mu|^{2}}{2}$$
(12)

and the jump operator is

$$\hat{c}'(\mu) = \hat{c} + \mu.$$
 (13)

Note that although for convenience we are not using the "eff" subscript anymore, $\hat{H}(\mu)$ is the *non-Hermitian* Hamiltonian \hat{H}_{eff} introduced earlier.

Under two-step adaptive monitoring, the signal from the system (a qubit) is mixed with the local oscillator with strength μ_1 prior to the photon detection. Before the photon is detected the system undergoes smooth evolution governed by $\hat{H}_1 = \hat{H}(\mu_1)$. Then the system experiences a jump governed by $\hat{s}_1 = \hat{c}'(\mu_1)$ when the photon is detected. At this point an experimentalist switches the strength of the local oscillator from μ_1 to μ_2 and the consequent smooth evolution is generated by the effective Hamiltonian $\hat{H}_2 = \hat{H}(\mu_2)$ and the next jump is caused by $\hat{s}_2 = \hat{c}'(\mu_2)$. As soon as this jump is

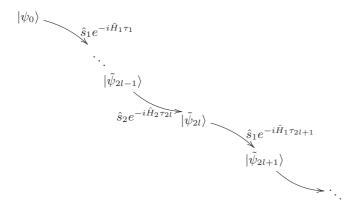


FIG. 1. A schematic of the conditioned state evolution under twostep adaptive monitoring.

detected, the experimentalist switches the strength of the local oscillator back to μ_1 and the evolution cycle repeats itself.

We can solve for the evolution of the system as follows. Let the initial state of the system be $|\psi_0\rangle$. We assume that the *n*th jump happens at time t_n after which the system will be in state $|\psi_n\rangle$ and the waiting times between jumps are given by $\tau_n = t_n - t_{n-1}$. The evolution is stochastic and for each realization of such evolution jumps will happen at different times. The probability density function, $p(\vec{\tau})$, for waiting times $\vec{\tau} = (\tau_1, \ldots, \tau_n)$ is given by

$$p(\vec{\tau}) = |||\tilde{\psi}_n\rangle||^2. \tag{14}$$

Here the unnormalized state $|\tilde{\psi}_n\rangle$ is defined recursively by

$$|\tilde{\psi}_n\rangle = \hat{s}_j e^{-i\hat{H}_j\tau_n} |\tilde{\psi}_{n-1}\rangle, \quad \text{with } j = \begin{cases} 1 & \text{for } n \text{ odd,} \\ 2 & \text{for } n \text{ even.} \end{cases}$$
(15)

This construction implies that $|\tilde{\psi}_{n-1}\rangle$ implicitly depends on $\tau_1, \ldots, \tau_{n-1}$ and, consequently, state $|\tilde{\psi}_n\rangle$ depends on all prior waiting times $\vec{\tau}$. The diagram for this evolution is depicted in Fig. 1. At an arbitrary time $t = t_{n-1} + s$, with $0 \le s < \tau_n$, for some *n*, the state of the system is given by

$$|\psi(t)\rangle = \frac{e^{-i\hat{H}_1s}|\tilde{\psi}_{n-1}\rangle}{||e^{-i\hat{H}_1s}|\tilde{\psi}_{n-1}\rangle||}, \quad \text{with } j = \begin{cases} 1 & \text{for } n \text{ odd,} \\ 2 & \text{for } n \text{ even.} \end{cases}$$
(16)

We can force the system evolution to be restricted to just two states, $|v_1^e\rangle$ and $|v_2^e\rangle$, if we require that the state $|v_1^e\rangle$ is an eigenstate of \hat{H}_1 , the state $|v_2^e\rangle$ is an eigenstate of \hat{H}_2 , the jump operator \hat{s}_1 maps $|v_1^e\rangle$ to $|v_2^e\rangle$, and \hat{s}_2 performs the reverse action. A consequence of these assumptions is that after two jumps the state $|v_1^e\rangle$ should return back to itself, i.e.,

$$\hat{s}_1 \hat{s}_2 \left| v_1^e \right\rangle \equiv \left[\mu_1 \mu_2 + (\mu_1 + \mu_2) \hat{c} + \hat{c}^2 \right] \left| v_1^e \right\rangle \propto \left| v_1^e \right\rangle.$$
(17)

Without loss of generality, we can assume that the jump operator \hat{c} is traceless and this means that the square of this operator is proportional to the identity, $\hat{c}^2 \propto \hat{I}$. Thus, Eq. (17) can happen only in two ways: either (i) $|v_1^e\rangle$ is an eigenstate of \hat{c} or (ii) $\mu_1 + \mu_2 = 0$. Case (i) holds iff $|v_1^e\rangle$ is an eigenstate of \hat{s}_1 , which would prevent any jumping from happening. Since case (i) is ruled out, case (ii) must hold. Thus, $\mu_1 = -\mu_2$ and the full cycle jump operator $\hat{s}_1 \hat{s}_2$ is proportional to the identity, which is a very useful fact for the discussion later on. We also comment here how Eqs. (14)–(16) can be used for simulation of a piecewise deterministic process and, in particular, for two-state jumping. We start the system in the normalized state $|\psi_0\rangle$. We know that after the *n*th jump the normalized state of the system is given by Eq. (16) with s =0. We determine a random waiting time τ_n according to the cumulative waiting time distribution [2]:

$$F(\tau_n) = 1 - ||e^{-i\dot{H}_j\tau_n}|\psi(\tau_{n-1})\rangle||^2,$$
(18)

with j = 1 if *n* is odd and j = 2 if *n* is even. This is done as follows. We draw a random number η from the uniform distribution over the interval [0, 1] and solve $\eta = 1 - F(\tau_n)$ for τ_n . Between the jumps the state of the system is given by Eq. (16). The evolution depicted in Figs. 3 and 4 is generated using this procedure.

The cumulative waiting time distribution, $F(\tau_n)$, can be, of course, related back to the probability density $p(\vec{\tau})$. Using the induction step

$$p(\tau_n) = p(\tau_n | \tau_{n-1}) p(\tau_{n-1}) = \frac{d}{d\tau_n} F(\tau_n) p(\tau_{n-1})$$
(19)

together with the fact that

$$\frac{d}{d\tau_n}F(\tau_n) = -\frac{d}{d\tau_n} ||e^{-i\hat{H}_j\tau_n}|\psi(\tau_{n-1})\rangle||^2
= \langle \psi(\tau_{n-1})|e^{i\hat{H}_j^{\dagger}\tau_n}\hat{s}_j^{\dagger}\hat{s}_j e^{-i\hat{H}_j\tau_n}|\psi(\tau_{n-1})\rangle
= \frac{|||\tilde{\psi}_n\rangle||^2}{|||\tilde{\psi}_{n-1}\rangle||^2},$$
(20)

we can deduce Eq. (14).

B. Mean-square stability

We now know how to describe the state of the system subject to two-state monitoring and we know the probability density for jumps occurring at times t_1, \ldots, t_n with waiting times $\vec{\tau} = (\tau_1, \ldots, \tau_n)$. The resulting evolution is a stochastic process, as we don't know when the jumps occur. Thus it is sensible to look at what happens to the system on average, over many different realizations, by considering the mean-square stability.

The mean-square fidelity between the desired system state and the actual system state immediately following the *n*th jump, starting from the initial state $|\psi_0\rangle$ is

$$\langle |\langle v_j^e | \psi_n \rangle|^2 \rangle = \int d^n \vec{\tau} \ p(\vec{\tau}) |\langle v_j^e | \psi_n \rangle|^2$$

$$= \int d^n \vec{\tau} \ |||\tilde{\psi}_n\rangle||^2 \frac{|\langle v_j^e | \tilde{\psi}_n \rangle|^2}{|||\tilde{\psi}_n\rangle||^2}$$

$$= \int d^n \vec{\tau} \ |\langle v_j^e | \tilde{\psi}_n \rangle|^2,$$
(21)

with j = 1 if n is odd and j = 2 if n is even. Thus this quantity is easy to compute. At the same time, it is easy to see that if

$$\lim_{n \to \infty} \left\langle \left| \left\langle v_j^e \middle| \psi_n \right\rangle \right|^2 \right\rangle = 1, \quad \text{with } j = \begin{cases} 1 & \text{if } n \text{ is even,} \\ 2 & \text{if } n \text{ is odd,} \end{cases}$$
(22)

the system will converge on average to a perfect jumping scenario. We call this *mean-square stability*.

We now investigate under what conditions the system has mean-square stability. To do this, we need a way to evaluate the dependence of $|\tilde{\psi}_n\rangle$ on the jump times, $\vec{\tau}$. We let $|v_j^e\rangle$ and $|v_j^o\rangle$ be the eigenstates of $i\hat{H}_j$ with respective eigenvalues λ_j^e and λ_j^o , where states $|v_j^e\rangle$ are part of the PR ensemble and states $|v_j^o\rangle$ are other states, not part of the ensemble. It is easiest to compute the time dependence of $|\tilde{\psi}_n\rangle$ if we decompose this state in terms of eigenstates of \hat{H}_1 for even *n* (i.e., $|v_1^e\rangle$ and $|v_1^o\rangle$) and in terms of eigenstates of \hat{H}_2 for odd *n* (i.e., $|v_2^e\rangle$ and $|v_2^o\rangle$). There is no loss of generality in doing this, since these pairs of eigenstates each form a basis (albeit nonorthogonal) for the qubit Hilbert space.¹

Thus we let the state of the system after 2l and 2l + 1 jumps be, respectively,

$$|\tilde{\psi}_{2l}\rangle = \alpha_{2l} |v_1^e\rangle + \beta_{2l} |v_1^o\rangle, \qquad (23)$$

$$|\tilde{\psi}_{2l+1}\rangle = \alpha_{2l+1} |v_2^e\rangle + \beta_{2l+1} |v_2^o\rangle.$$
 (24)

The dependence on the jump times is hidden in the coefficients α_{2l} , α_{2l+1} , β_{2l} , and β_{2l+1} , which we determine by establishing the relationship between them and their dependence on α_0 and β_0 . To do this, we first define that the action of jump operators \hat{s}_1 and \hat{s}_2 is on the basis states $\{|v_1^e\rangle, |v_1^o\rangle\}$ and $\{|v_2^e\rangle, |v_2^o\rangle\}$:

$$\begin{aligned} \hat{s}_{1} | v_{1}^{e} \rangle &= Q_{11}^{1} | v_{2}^{e} \rangle, \\ \hat{s}_{1} | v_{1}^{o} \rangle &= Q_{21}^{1} | v_{2}^{e} \rangle + Q_{22}^{1} | v_{2}^{o} \rangle, \\ \hat{s}_{2} | v_{2}^{e} \rangle &= Q_{21}^{2} | v_{1}^{e} \rangle, \\ \hat{s}_{2} | v_{2}^{o} \rangle &= Q_{21}^{2} | v_{1}^{e} \rangle + Q_{22}^{2} | v_{1}^{o} \rangle. \end{aligned}$$

$$(25)$$

The scalar elements Q_{jk}^1 and Q_{jk}^2 can be determined once all the parameters of the system, \hat{H} , \hat{c} , and μ_1 , are specified. Equations (25) simply state that jumping operators (\hat{s}_1 and \hat{s}_2) map jumping states (\hat{v}_1^e and \hat{v}_2^e) to each other and the mapping for the other basis states is quite general.

Using Eqs. (23)–(25) and the propagation between states given in Eq. (15), we can derive

$$\beta_{2l+1} = \beta_{2l} e^{-\lambda_1^o \tau_{2l+1}} Q_{22}^1, \tag{26}$$

$$\beta_{2l} = \beta_{2l-1} e^{-\lambda_2^o \tau_{2l}} Q_{22}^2.$$
(27)

Now we can show by induction that

$$\beta_{2l} = \beta_0 \left(Q_{22}^1 Q_{22}^2 \right)^l \prod_{j=1}^l e^{-\lambda_1^o \tau_{2j-1}} e^{-\lambda_2^o \tau_{2j}}, \qquad (28)$$

$$\beta_{2l+1} = \beta_0 Q_{22}^1 (Q_{22}^1 Q_{22}^2)^l \prod_{j=0}^l e^{-\lambda_2^o \tau_{2j}} e^{-\lambda_1^o \tau_{2j+1}}.$$
 (29)

This information is sufficient to derive the mean-square fidelity $\langle |\langle v_i^e | \psi_n \rangle|^2 \rangle$, using the following useful identity:

$$\left| \left\langle v_{j}^{e} \middle| \tilde{\psi}_{n} \right\rangle \right|^{2} = |||\tilde{\psi}_{n}\rangle||^{2} - |\beta_{n}|^{2}(1 - |O_{j}|^{2}), \qquad (30)$$

where j = 1 if *n* is even and j = 2 if *n* is odd and $O_j = \langle v_j^e | v_j^o \rangle$ is the overlap function. Now calculating the meansquare fidelity becomes easy. According to Eq. (14), $|||\tilde{\psi}_n\rangle||^2$ is the probability density for jumps occurring with waiting times $\vec{\tau} = (\tau_1, \dots, \tau_n)$ and thus when integrated with respect to $d^n \vec{\tau}$ yields 1. Thus

$$\left\langle \left| \left\langle v_{j}^{e} \left| \psi_{n} \right\rangle \right|^{2} \right\rangle = 1 - (1 - |O_{j}|^{2}) \int d^{n} \vec{\tau} \left| \beta_{n} \right|^{2}, \qquad (31)$$

with j = 1 if n is even and j = 2 if n is odd. Using Eq. (28), we learn

$$\int d^{2l} \vec{\tau} |\beta_{2l}|^2 = |\beta_0|^2 \left(\frac{|Q_{22}^1 Q_{22}^2|^2}{4\text{Re}\lambda_1^o \text{Re}\lambda_2^o} \right)^l \equiv |\beta_0|^2 C^l. \quad (32)$$

Thus the mean-square fidelity is

$$\left| \left\langle v_1^e | \psi_{2l} \right\rangle \right|^2 \right\rangle = 1 - |\beta_0|^2 (1 - |O_1|^2) C^l$$
(33)

and

$$\left\langle \left| \left\langle v_{2}^{e} \right| \psi_{2l+1} \right\rangle \right|^{2} \right\rangle = 1 - \left| Q_{22}^{1} \beta_{0} \right|^{2} (1 - |O_{2}|^{2}) C^{l}.$$
 (34)

Thus coefficient C determines the convergence rate to the perfect two-state jumping.

The absolute value of fidelity is a number between 0 and 1 and its average with respect to all possible waiting times $\vec{\tau}$ is still between 0 and 1. Thus Eq. (33) implies that $C \leq 1$. Thus unless C = 1, Eq. (34) implies that the mean-square fidelity always converges to 1 and the system has mean-square stability.

We now determine under what conditions parameter *C* is strictly less than one. To answer this question, we need one more calculation, which is performed in Appendix A and establishes that $|Q_{11}^1Q_{11}^2|^2/4\text{Re}\lambda_1^e\text{Re}\lambda_2^e = 1$ and that $Q_{11}^1Q_{11}^2 = Q_{12}^2Q_{22}^2$. Thus

$$C = \frac{\text{Re}\lambda_1^e \text{Re}\lambda_2^e}{\text{Re}\lambda_1^e \text{Re}\lambda_2^o}.$$
(35)

In other words, the system will have mean-square stability iff the geometric mean of the real part of the eigenvalues of the states *in* the ensemble is smaller than the geometric mean of the real part of the eigenvalues of the states *not in* the ensemble. This makes sense, as the real parts of these eigenvalues give the rate at which the amplitudes of their respective states decay during the between-jump stages of the evolution, as we explore in Sec. III D. Thus, the smaller they are, the more stable the respective states. As we see in Sec. IV B, the condition for mean-square stability can be satisfied even though every second stage is *unstable*.

¹This discussion leaves out one technical detail. Because \hat{H}_1 and \hat{H}_2 are non-Hermitian operators, there could be special circumstances when one (or both) of these operators cannot be reduced to a diagonal matrix via similarity transformation. In these cases (of measure 0), such an operator \hat{H}_j can still be reduced to a Jordan normal form. For this there is only one eigenvector $|v_j^e\rangle$ of $i\hat{H}_j$, with eigenvalue λ_j^e , which is a member of the PR ensemble. But there is a generalized eigenvector $|w_j\rangle$ satisfying $(i\hat{H}_j - \lambda_j^e I)|w_j\rangle = |v_j^e\rangle$. The set $\{|v_j^e\rangle, |w_j\rangle\}$ spans the qubit Hilbert space, and it is easy to verify that the calculation goes through exactly the same using $|w_j\rangle$ instead of $|v_j^o\rangle$. We emphasize that, in all the examples we consider, all of the \hat{H}_n do have a complete set of eigenvalues so this issue does not arise.

C. K-state jumping

We now explain how our results can be generalized to cyclic *K*-state jumping, as first studied in Ref. [6] for K > 2. We simply let the index *j* for expressions in Secs. III A and III B range over $\{1, \ldots, K\}$ instead of just taking values 1 and 2. All considerations still hold and we can derive Eq. (31) as before, but the index *j* in this expression is now determined according to $j = (n \mod K) + 1$. We also need to reevaluate $\int d^n \vec{\tau} |\beta_n|^2$. The unnormalized state for the *j*th step in the cycle after *n* repeats is

$$\left|\tilde{\psi}_{N}\right\rangle = \alpha_{N}\left|v_{j+1}^{e}\right\rangle + \beta_{N}\left|v_{j+1}^{o}\right\rangle$$
(36)

with N = nK + j and j + 1 stands for $(j + 1) \mod K$. Equation (25) becomes

$$\hat{s}_{j}|v_{j}^{e}\rangle = Q_{11}^{j}|v_{j+1}^{e}\rangle,
\hat{s}_{j}|v_{j}^{o}\rangle = Q_{21}^{j}|v_{j+1}^{e}\rangle + Q_{22}^{j}|v_{j+1}^{o}\rangle,$$
(37)

and the expression for β_{nK} becomes

$$\beta_{nK} = \beta_0 (Q_{22}^1 \cdots Q_{22}^K)^n \prod_{j=0}^{n-1} e^{-\lambda_1^o \tau_{jK+1}} \cdots e^{-\lambda_K^o \tau_{(j+1)K}}$$
(38)

and thus

$$\int d^{nK} \vec{\tau} |\beta_{nK}|^2 = |\beta_0|^2 \left(\frac{\mathcal{Q}_{22}^1 \cdots \mathcal{Q}_{22}^K}{2^n \operatorname{Re} \lambda_1^o \cdots \operatorname{Re} \lambda_K^o} \right)^n.$$
(39)

Thus parameter C which determines stability is redefined to be

$$C = \frac{Q_{22}^1 \cdots Q_{22}^K}{2^n \operatorname{Re} \lambda_1^o \cdots \operatorname{Re} \lambda_K^o}.$$
 (40)

Just as before we can show that C can be reexpressed as

$$C = \frac{\operatorname{Re}\lambda_1^e \cdots \operatorname{Re}\lambda_K^e}{\operatorname{Re}\lambda_1^o \cdots \operatorname{Re}\lambda_K^o},\tag{41}$$

so that C < 1 whenever $\operatorname{Re}\lambda_1^o \cdots \operatorname{Re}\lambda_K^o > \operatorname{Re}\lambda_1^e \cdots \operatorname{Re}\lambda_K^e$. The argument proceeds in the same way as for two-state jumping, and the details appear in Appendix A.

We conjecture that the fact that $|v_j^e\rangle$ form a PR ensemble for ρ_{ss} , the unique steady state, ensures that $\prod_{k=1}^{K} \operatorname{Re} \lambda_k^o > \prod_{k=1}^{K} \operatorname{Re} \lambda_k^e$. Even more generally, we conjecture that every finite PR ensemble is stable.

D. Stability of an individual trajectory

It is important to note that mean-square stability does not imply that in an individual trajectory there will be monotonic convergence of the system toward the desired states. The fidelity could decrease between the jumps and/or could decrease due to a jump. We begin with the first issue.

We say that the system is *piecewise stable* if $|\langle v_j^e | \psi(t) \rangle|^2$ is a monotonically increasing function of t during all stages $t_{n-1} \leq t < t_n$, where j = 1 if n is odd and j = 2 if n is even. Otherwise, if the fidelity with the desired state decreases between jumps, this constitutes an *unstable stage* in the evolution. As we will show, this happens only if n is even or n is odd, not both. That is, stable and unstable stages will alternate.

The stability of the evolution between jumps is determined by the relationship between eigenvalues λ_j^e and λ_j^o . Let us assume without loss of generality that the state of the system right after the jump is

$$|\psi_0\rangle = \alpha_0 |v_i^e\rangle + \beta_0 |v_i^o\rangle. \tag{42}$$

Then the unnormalized state of the system at time τ after the jump, but before the next one, is

$$\left|\tilde{\psi}(\tau)\right\rangle = \alpha_0 e^{-\lambda_j^e \tau} \left|v_j^e\right\rangle + \beta_0 e^{-\lambda_j^o \tau} \left|v_j^o\right\rangle. \tag{43}$$

We now want to know what happens to the fidelity between $|\psi(\tau)\rangle$ (normalized state of the system between jumps) and the ideal jumping state $|v_i^e\rangle$:

$$\left| \left\langle v_j^e | \psi(\tau) \right\rangle \right|^2 = \frac{\left| \left\langle v_j^e | \tilde{\psi}(\tau) \right\rangle \right|^2}{|||\tilde{\psi}(\tau)\rangle||^2}.$$
(44)

Using Eq. (30), we can compute this fidelity to be

$$\left| \left\langle v_{j}^{e} \right| \psi(\tau) \right\rangle \right|^{2} = 1 - \frac{|\beta_{0}|^{2} (1 - |O_{j}|^{2})}{|\alpha_{0} e^{-(\lambda_{j}^{e} - \lambda_{j}^{o})\tau} |v_{j}^{e}\rangle + \beta_{0} |v_{j}^{o}\rangle|^{2}}.$$
 (45)

Thus we can see that if $\operatorname{Re}\lambda_j^e < \operatorname{Re}\lambda_j^o$ then in the long time limit fidelity approaches 1. That is, between jumps the system will converge to the desired ensemble state with exponential rate and the fidelity will monotonically increase to the maximum value of 1. On the other hand, if $\operatorname{Re}\lambda_j^e > \operatorname{Re}\lambda_j^o$, then the system state converges toward the non-ensemble eigenstate, and the fidelity converges toward $|O_j|^2$, which is a quantity less than 1. Thus the fidelity may decrease during such a stage, and we call this an unstable stage in the evolution. This is illustrated for the case of resonance fluorescence in Sec. IV B.

We now address the issue of whether the jump itself increases or decreases fidelity. We again introduce some notation. Let $|\tilde{\phi}_1\rangle = \alpha_1 |v_j^e\rangle + \beta_1 |v_j^o\rangle$ be the state of the system right before the jump and $|\tilde{\phi}_2\rangle = \alpha_2 |v_k^e\rangle + \beta_2 |v_k^o\rangle$ be the unnormalized state of the system right after the jump, where $k = (j + 1) \mod K$. We also know that $|\tilde{\phi}_2\rangle = \hat{s}_j |\tilde{\phi}_1\rangle$. We want to compare the fidelity right before and immediately after the jump and we let

$$F_1 = \frac{\left| \left\langle v_j^e | \tilde{\phi}_1 \right\rangle \right|^2}{|||\tilde{\phi}_1\rangle||^2} \quad \text{and} \quad F_2 = \frac{\left| \left\langle v_k^e | \tilde{\phi}_2 \right\rangle \right|^2}{|||\tilde{\phi}_2\rangle||^2}.$$
(46)

We evaluate $|\langle v_j^e | \tilde{\phi}_1 \rangle|^2$ and $|\langle v_k^e | \tilde{\phi}_2 \rangle|^2$ using Eq. (30). Then fidelity will decrease after the jump, $F_1 < F_2$, iff

$$\frac{|||\tilde{\phi}_2\rangle||^2}{|||\tilde{\phi}_1\rangle||^2} < \frac{|\beta_2|^2}{|\beta_1|^2} \frac{1 - |O_k|^2}{1 - |O_j|^2} \equiv B.$$
(47)

We can easily determine tight bounds for $|||\tilde{\phi}_2\rangle||^2/|||\tilde{\phi}_1\rangle||^2$ by observing that

$$\lambda_{\min}|||\tilde{\phi}_1\rangle||^2 \leqslant |||\tilde{\phi}_2\rangle||^2 = \langle \tilde{\phi}_1|\hat{s}_j^{\dagger}\hat{s}_j|\tilde{\phi}_1\rangle \leqslant \lambda_{\max}|||\tilde{\phi}_1\rangle||^2,$$
(48)

where λ_{\min} and λ_{\max} are the smallest and the largest eigenvalues of the operator $\hat{s}_j^{\dagger} \hat{s}_j$. Thus the lower bound of λ_{\min} , or upper bound of λ_{\max} , is attained if $|\tilde{\phi}_1\rangle$ is the corresponding eigenstate of $\hat{s}_j^{\dagger} \hat{s}_j$. If $\lambda_{\min} < B$, it is possible to observe jumps that decrease the fidelity. Again, we illustrate this for the resonance fluorescence example in Sec. IV B. Note however that for mean-square stable evolution, in the long time limit, $|\phi_1\rangle$ converges to $|v_j^e\rangle$, and so F_1 and F_2 also converge to one. Thus we expect that as the system converges toward the ideal ensemble states, observing jumps that decrease fidelity becomes less and less likely.

IV. RESONANCE FLUORESCENCE

In this section we apply our results for stability in the mean-square sense, and for individual trajectories, to a specific physical example. In particular, we show that there are PR ensembles which are mean-square stable, but whose cycles are composed of stable and unstable stages. Others are composed of only stable stages, and so the fidelity is piecewise monotonically increasing. It is only piecewise because this example also illustrates that jumps can decrease the fidelity.

The example we consider is resonance fluorescence. We consider a qubit (for example, a two-level atom) with basis states $|0\rangle$ and $|1\rangle$, with a transition frequency ω_0 . We assume that qubit is coupled to the continuum of electromagnetic radiation and, therefore, decays to $|0\rangle$ at rate γ . At the same time, it is driven by a classical field oscillating at frequency ω_0 . The strength of the driving is quantified by the Rabi frequency Ω . In the interaction frame [3] with respect to the atomic transition frequency ω_0 , the evolution of the qubit is given by the master equation of the form of Eq. (1) with $\hat{H} = (\Omega/2)\hat{\sigma}_x$ and one jump operator $\hat{c} = \sqrt{\gamma}\hat{\sigma}$, where $\hat{\sigma} = |0\rangle\langle 1|$ and $\hat{\sigma}_x = \hat{\sigma} + \hat{\sigma}^{\dagger}$. Then matrix A and \vec{b} in the Bloch vector equation, Eq. (4), are

$$A = \begin{pmatrix} -\gamma/2 & 0 & 0\\ 0 & -\gamma/2 & -\Omega\\ 0 & \Omega & -\gamma \end{pmatrix} \quad \text{and} \quad \vec{b} = \begin{pmatrix} 0\\ 0\\ \gamma \end{pmatrix}.$$
(49)

The steady state $\vec{r}_{ss} = (0, 2\gamma\Omega, -\gamma^2)^T/(\gamma^2 + 2\Omega^2)$ is a mixed state for $\Omega \neq 0$. The unnormalized eigenvectors of A are $\vec{u}_1 = (1, 0\,0)^T$ and $\vec{u}_{\pm} = (0, \gamma \pm \sqrt{\gamma^2 - 16\Omega^2}, 4\Omega)^T$. For $\varepsilon \equiv \Omega/\gamma$ and $|\varepsilon| < 1/4$ all three eigenvectors of A are real, while for $|\varepsilon| > 1/4$ only \vec{u}_1 is real. Note that ε here is different from $\epsilon = (\Omega/\gamma)^2$ in Ref. [6].

A. Two-state jumping

We now analyze stability properties for two-state jumping in this system and we do this with the method described in Sec. III B and in Refs. [14,15]; i.e., we know that the PR ensemble is constructed from the eigenvectors of $\hat{H}(\mu)$, or equivalently, of $i\hat{H}(\mu)$, where $\hat{H}(\mu)$ is the effective (non-Hermitian) Hamiltonian given by Eq. (12).

For resonance fluorescence, the operator $-i\hat{H}(\mu)$ has eigenvectors given by

$$|v_{\pm}(\mu)\rangle = \varepsilon|1\rangle + \left(\pm\sqrt{\varepsilon^2 - \frac{1}{4} - 2i\varepsilon\mu^*} + \frac{i}{2}\right)|0\rangle.$$
 (50)

The corresponding eigenvalues are

$$\lambda_{\pm}(\mu) = \frac{1+2|\mu|^2}{4} \pm \frac{i}{2}\sqrt{\varepsilon^2 - \frac{1}{4} - 2i\varepsilon\mu^*},$$
 (51)

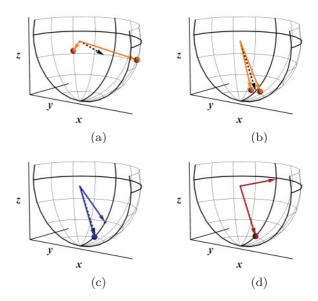


FIG. 2. (Color online) Solid arrows show Bloch vectors for twostate jumping. The volume of the sphere at the tip of each arrow represents the probability that the qubit occupies the corresponding pure state. The dashed arrow is \vec{r}_{ss} . Solutions in panels (a) and (b) arise from \vec{u}_1 with $\varepsilon = 1$ and $\varepsilon = 0.23$, respectively. This is the solution that exists for all ε . Solutions depicted in panels (c) and (d) exist only for $\varepsilon \leq 0.25$ and are shown here for $\varepsilon = 0.23$. The solution in panel (c) is generated by \vec{u}_+ and that in panel (d) by \vec{u}_- .

where here, and in the remainder of the paper, we have set $\gamma = 1$ for simplicity. Whether $|v_+(\mu)\rangle$ or $|v_-(\mu)\rangle$ is part of the PR ensemble depends on a particular value of μ_1 . We know already from Sec. III A that $\mu_2 = -\mu_1$, so without loss of generality we can choose μ_1 to have a non-negative real part. Once we impose additional conditions from Sec. IV C associated with jump operators \hat{s}_1 and \hat{s}_2 , we learn that μ_1 only assumes values from the set $\{1/2, \nu_+, \nu_-\}$, where

$$\nu_{\pm} = i \frac{\sqrt{1 \pm \sqrt{1 - 16\varepsilon^2}}}{2\sqrt{2}}.$$
 (52)

The last two values in the set contribute to solutions only for $\varepsilon \leq 1/4$. That is, one only obtains a PR ensemble for values of μ_1 that are either real or purely imaginary. Detailed derivation of these conditions can be found in Refs. [14,15], and the Bloch vectors for all three ensembles are shown in Fig. 2.

For $\mu_1 = 1/2$, the jumping states in Eqs. (7) and (8) are spanned from eigenvector \vec{u}_1 and, as explained in Ref. [6], the system spends an equal amount of time in each state. Equivalently, the probability of occupying each state is 1/2and the Shannon entropy is 1. Thus one bit is sufficient to track the state of the system. As shown in Figs. 2(a) and 2(b), the solution is symmetric with respect to the x = 0 plane and as ε increases jumping states move away from each other and away from the bottom of the Bloch sphere. For $\mu_1 = v_+$ and $\mu_1 = v_-$ the jumping states are generated by \vec{u}_+ and \vec{u}_- , respectively. Figures 2(c) and 2(d) show that such solutions lie in the x = 0 plane. For $\mu_1 = v_-$, jumping states are spread as far apart as possible; the system spends most of the time in one state, which is nearly aligned with the direction of the steady state, \vec{r}_{ss} . For $\mu_1 = v_+$, jumping states cluster together around the direction of the steady state, \vec{r}_{ss} ; the probability of occupying each state is more similar and becomes equal as ε approaches 0. Since the probability of occupying each state is no longer equal, the Shannon entropy is less than 1. This means that one could store the state of the qubit in *less than* one bit *on average*. That is, one could keep track of the state of a collection of N identically monitored qubits using only Nh bits. Actually, the entropy of the solution due to \vec{v}_{-} is very low. That is, it is very close to the von Neumann entropy of the steady state, which is a lower bound for the Shannon entropy of any ensemble representing for the steady state. For this solution, the Shannon entropy in a Taylor expansion in terms of ε is

$$h(\vec{u}_{-}) = (\log_2 10 - 4\log_2 \varepsilon)\varepsilon^4 - 48(\log_2 \varepsilon)\varepsilon^6 + O(\varepsilon^8)$$
(53)

compared to

$$S(\rho_{\rm ss}) = (\log_2 10 - 4\log_2 \varepsilon)\varepsilon^4 + 6(\log_2 \varepsilon)\varepsilon^6 + O(\varepsilon^8).$$
(54)

B. Stability of two-state jumping

For $\mu_1 = 1/2$, the PR states are $|v_+(\mu_1)\rangle$ and $|v_+(-\mu_1)\rangle$ and the full set of eigenvalues from Eq. (51) can be simplified to

$$\lambda_{+}\left(\pm\frac{1}{2}\right) = \frac{1}{8} \pm \frac{i\varepsilon}{2}$$
 and $\lambda_{-}\left(\pm\frac{1}{2}\right) = \frac{5}{8} \pm \frac{i\varepsilon}{2}$. (55)

Thus $\text{Re}\lambda_1^e = \text{Re}\lambda_2^e = 1/8$, while $\text{Re}\lambda_1^o = \text{Re}\lambda_2^o = 5/8$. Therefore C = 1/25, and the system is mean-square stable. Moreover, it is piecewise deterministically stable (that is, the fidelity increases monotonically except perhaps at jumps.)

For $\mu_1 = \nu_-$, the situation is quite different. The PR ensemble is composed of $|v_{+}(\mu_{1})\rangle$ and $|v_{-}(-\mu_{1})\rangle$. It is still mean-square stable, but the short-lived state $|v_2^e\rangle = |v_-(-\mu_1)\rangle$ has an eigenvalue such that $\operatorname{Re}\lambda_2^e > \operatorname{Re}\lambda_2^o$, where $|v_2^o\rangle =$ $|v_{+}(-\mu_{1})\rangle$. That is, every second stage is unstable. We illustrate this situation in Figs. 3 and 4. Here we assume that evolution starts with the unstable step; that is, the initial strength of the local oscillator is $-\mu_1$. We also let $\varepsilon = 0.1$ and the initial state is $|\tilde{\psi}_0\rangle = -1.09|v_-(-\mu_1)\rangle + 0.5|v_+(-\mu_1)\rangle$. Figure 3 shows highly untypical evolution, where the shortlived state $|v_{+}(-\mu_{1})\rangle$ exists for an improbably long time and the long-lived state $|v_{-}(-\mu_{1})\rangle$ lasts an improbably short time. One can see here that for stages governed by $-\mu_1$ the fidelity decays as the system moves away from the desired PR ensemble state. Figure 4 shows typical evolution under the same conditions. Here the instability can be seen only during the first step, which is short (at time $\tau_1 \approx 1.1 \gamma^{-1}$). The next jump does not happen for a long time ($\tau_2 \approx 792.5\gamma^{-1}$) and it is very short (it lasts for $\tau_3 \approx 0.3 \gamma^{-1}$). Thus the unstable stages contribute very little to the overall evolution of the fidelity. As a result it does not disturb mean-square stability for the system.

For $\mu_1 = \nu_+$ the situation is even more curious. For $\varepsilon < \varepsilon_0 \approx 0.243$, the PR ensemble is generated from $|v_+(\mu_1)\rangle$ and $|v_+(-\mu_1)\rangle$, and for $\varepsilon > \varepsilon_0$, the PR ensemble is generated from $|v_+(\mu_1)\rangle$ and $|v_-(-\mu_1)\rangle$. The system is always mean-square stable and is piecewise deterministically stable for

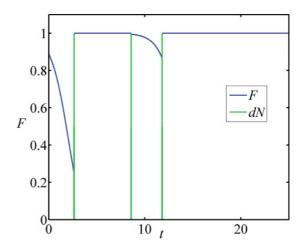


FIG. 3. (Color online) Alternating stages of stable and unstable two-jump evolution, exaggerated by an atypical trajectory. The blue (darker) line represents the fidelity of the state with the ideal PR ensemble and the green (lighter) line shows when jumps happen (photon count increment dN = 1). This is a highly improbable evolution, as the systems spends an uncharacteristically long time in an unstable step and an uncharacteristically short time in a stable step. Fidelity is decreasing during the unstable step. At the second jump there is also a decrease in fidelity due to the jump itself. Time *t* is measured in units of γ^{-1} .

 $\varepsilon < \varepsilon_0$. However for $\varepsilon > \varepsilon_0$ the short-lived stage associated with $|v_-(-\mu_1)\rangle$ becomes unstable. At $\varepsilon = \varepsilon_0$, $\lambda_+(-\mu_1) = \lambda_-(-\mu_1)$. In Fig. 5, we show the entropy for three different solutions with two-state jumping and the type of stability they exhibit.

Recall from Eq. (32) that the deviation of the fidelity from unity decays as C^l , after 2l jumps. The critical constant for all three values of μ_1 can be written as

$$C = \frac{|\mu_1|^4}{4\mathrm{Re}\lambda_1^{o}\mathrm{Re}\lambda_2^{o}}.$$
(56)

We show the convergence coefficient *C* as a function of ε for all three possible values of μ_1 in Fig. 6. There is a clear (although not perfect) correlation between the low entropy solutions and the most quickly converging (in the mean-square sense)

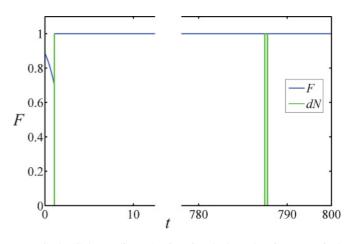


FIG. 4. (Color online) As in Fig. 3, but showing a typical trajectory. The fidelity decreases visibly only during the first stage, as stable stages last much longer than unstable stages.

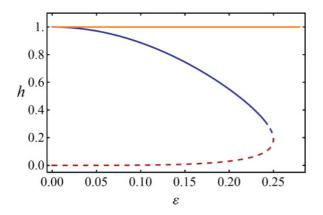


FIG. 5. (Color online) Ensemble Shannon entropy *h* for the three different two-state jumping solutions. A solid line indicates mean-square stability as well as piecewise deterministic stability. A dashed line indicates that the solution has only mean-square stability, with one of its stages being unstable. The uppermost (yellow) line shows entropy for the PR ensemble with $\mu_1 = 1/2$. The middle line (blue) describes entropy for solutions with $\mu_1 = \nu_+$ and the lowest line (red) is for $\mu_1 = \nu_-$.

solutions. Note however that the most rapidly converging solution is typically *not* piecewise deterministically stable.

We note here that from Eq. (33), one can see that convergence coefficient *C* characterizes only certain aspects of mean-square stability. In particular, it determines how many stages are needed for convergence to the PR ensemble. Thus the smaller *C* in Fig. 6 corresponds to fewer stages needed to achieve a certain level of convergence in the mean-square sense. This plot, however, does not tell us about the time needed to achieve convergence. For example, although the solution with $\mu_1 = \nu_-$ has a much smaller value for *C* than the $\mu_1 = 1/2$ solution, and so requires more jumps to converge, it could have shorter stages between every jump so that it approaches fidelity in a shorter time than the $\mu_1 = \nu_-$ solution. To show that this is a plausible scenario we consider an asymptotic rate of convergence to the PR ensemble, which we define as $R = -\ln(C)/\langle t_{as} \rangle$, where $\langle t_{as} \rangle$ is the expected

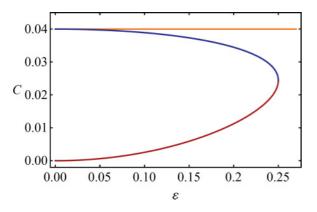


FIG. 6. (Color online) The convergence coefficient *C* as a function of ε for the three different values for μ_1 . The yellow (straight) line corresponds to $\mu_1 = 1/2$, the red (lowest) line to $\mu_1 = \nu_-$, and the blue (intermediate) line to $\mu_1 = \nu_+$. Note that the lowest entropy solution, which is not piecewise deterministically stable, is the most stable (smallest *C*) in the mean-square sense.

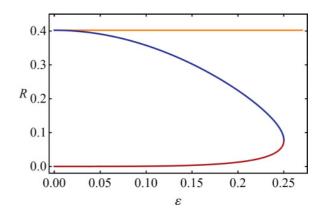


FIG. 7. (Color online) The asymptotic convergence rate *R* as a function of ε for the three different values for μ_1 . The yellow (straight) line corresponds to $\mu_1 = 1/2$, the red (lowest) line to $\mu_1 = \nu_-$, and the blue (intermediate) line to $\mu_1 = \nu_+$. Note that a high value of *R* indicates faster convergence (per unit time), so this measure of convergence reverses the order of the three schemes relative to that in Fig. 6.

value for the duration of the complete cycle of evolution for the corresponding PR ensemble. Because asymptotically the actual ensemble converges to the PR ensemble, and using the law of large numbers, the number of cycles undergone in the limit $t \rightarrow \infty$ converges to $t/\langle t_{as} \rangle$. Thus from Eq. (33), in the long time limit, the fidelity F(t) between the actual conditioned state and the record-determined PR state at time *t* behaves as

$$\langle F(t) \rangle \sim 1 - |\beta_0|^2 (1 - |O_1|^2) \exp(-Rt).$$
 (57)

For two-state jumping this asymptotic convergence rate is

$$R = -\ln(C) / \left[\left(2 \text{Re} \lambda_1^{\text{e}} \right)^{-1} + \left(2 \text{Re} \lambda_2^{\text{e}} \right)^{-1} \right], \quad (58)$$

and we plot it in Fig. 7 as a function of ε for all three values of μ_1 . Here the solution with $\mu_1 = 1/2$ has the highest asymptotic convergence rate, with $R = \ln(5)/4$, while the solution $\mu_1 = \nu_-$ has the lowest, with $R \to 0$ as $\varepsilon \to 0$. Specifically, for $\mu_1 = \nu_-$, although $\ln(C) = O[\log(\varepsilon^2)] \gg 1$, $\operatorname{Re}\lambda_1^e = O(\varepsilon^4)$ so that $R = O[\varepsilon^4 \log(\varepsilon^2)] \to 0$. This is a complete reversal from the results reported in Fig. 6, based on *C*.

The vast difference in rates in the limit $\varepsilon \to 0$ can be understood from the nature of the ensembles as illustrated in Fig. 2. In all cases the the average ρ_{ss} differs from $|0\rangle\langle 0|$ only at $O(\varepsilon)$ and differs from a pure state only at $O(\varepsilon^4)$, as reflected in Eq. (54). In the case $\mu = 1/2$, the ensemble comprises a pair of Bloch vectors located near \vec{r}_{ss} in a symmetric fashion in the plane perpendicular to the x = 0 plane. For $\mu_1 = \nu_+$ the Bloch vectors lie in the x = 0 plane around \vec{r}_{ss} in a nearly symmetric fashion. In both cases, the system spends the same time in each state and thus the expected time for the duration of each cycle is independent of ε , so the rate of convergence is nonzero. For $\mu_1 = \nu_-$, one of the jumping states is nearly aligned with ρ_{ss} [i.e., it has the form $|0\rangle + O(\varepsilon)|1\rangle$] and the other state approaches $|1\rangle$ as $\varepsilon \to 0$. For small ε , the system spends almost all of the time in the first state. It jumps to the excited state with a rate of $O(\varepsilon^4)$ and jumps back with a finite rate. The former process is the rate-limiting step, so the system

has a cycle whose duration tends to infinity as $\varepsilon \to 0$, giving a convergence rate near 0.

However, we need to be careful in interpreting the results from Fig. 7 as it only shows the asymptotic rate. Most of the convergence happens in the initial evolution—during the first few stages, while the state of the system is drastically different from the PR ensemble. The asymptotic rate becomes relevant only when the state of the system closely resembles the PR ensemble and, in a way, the convergence has already occurred. We, therefore, also consider another measure for the rate of convergence: $R_1 = -\ln(C)/\langle T_1 \rangle$. Here $\langle T_1 \rangle$ is the expected time for the duration of the first cycle. For two-state jumping, this quantity was investigated numerically.

Our numerics reveal that the initial convergence rate R_1 shows strong dependence on the initial condition. For some initial conditions, R_1 for the $\mu_1 = \nu_+$ solution is above that for the $\mu_1 = 1/2$ solution for all $\varepsilon \leq 1/4$. Because the $\nu_$ and ν_+ solutions coincide for $\varepsilon = 1/4$, this implies that, for these initial conditions, and for ε sufficiently close to 1/4, the $\mu_1 = 1/2$ solution has the slowest initial rate of convergence for some range of ε near 1/4. On the other hand, for other initial conditions, R_1 for the $\mu_1 = \nu_+$ solution is below that for the $\mu_1 = 1/2$ solution for all $\epsilon \leq 1/4$, and the same ordering as seen in the asymptotic limit (Fig. 7) occurs for some range of ε near 1/4. Because $-\ln(C)$ diverges (slowly) to $+\infty$ as $\varepsilon \to 0$ for $\mu_1 = \nu_-$, for a fixed initial condition this last solution always has the largest initial rate of convergence R_1 as $\varepsilon \to 0$.

The last question we address in this section is the effect jumps have on fidelity. For the model we are considering we can easily evaluate the upper and lower bounds in Eq. (48) to be

$$\lambda_{\min} = |\mu_1|^2 + 1/2 - \sqrt{|\mu_1|^2 + 1/4},$$
(59)

$$\lambda_{\max} = |\mu_1|^2 + 1/2 + \sqrt{|\mu_1|^2 + 1/4}.$$
 (60)

For $\mu_1 = 1/2$, these give the bounds

$$0.043 \approx \frac{3}{4} - \frac{1}{\sqrt{2}} \leqslant \frac{|||\tilde{\phi}_2\rangle||^2}{|||\tilde{\phi}_1\rangle||^2} < \frac{3}{4} + \frac{1}{\sqrt{2}} \approx 1.46.$$
(61)

To calculate *B* from Eq. (47), we observe that for $\mu_1 = 1/2$ we have observed that $|O_1| = |O_2|$ and $|\beta_2| = |\mu_1| \times |\beta_1|$. This gives $B = |\beta_2|^2/|\beta_1|^2 = 1/4$. Because this is greater than 0.043 it follows that sometimes jumps can decrease the fidelity. For $\mu_1 = \nu_-$, we determine *B* and the bounds numerically. We find similar results, as plotted in Fig. 8. This time however *B* depends on which jump operator one is considering, and a fidelity-decreasing jump is much less likely following a stage of *unstable* evolution. Our results confirm what was seen in the second jump in the untypical evolution depicted in Fig. 3, where a jump following a stable stage of evolution visibly decreased the fidelity.

C. Three-state jumping

We now consider three-state PR ensembles with cyclic jumps as in Ref. [6], still for a qubit subject to resonance fluorescence. In this case, in Eq. (6) only κ_{12} , κ_{23} , and κ_{31} are nonzero and Eqs. (5) and (6) yield a total of 12 equations (9 equations for jumping conditions and 3 equations for normalization condition) and 12 unknowns. Since this system

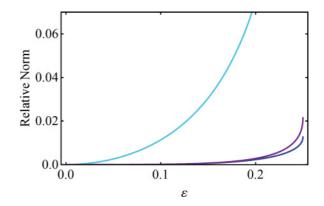


FIG. 8. (Color online) Quantities relevant to a fidelity decrease upon jumping in the two-state jumping solution with $\mu_1 = \nu_-$. The blue (lowest) line shows the achievable lower bound for the relative norm $|||\tilde{\phi}_2\rangle||^2/|||\tilde{\phi}_1\rangle||^2$. The cyan (uppermost) line shows the bound *B* the jump generated by \hat{s}_1 and the purple (middle) line is the same bound for jumps generated by \hat{s}_2 . Since the curves for *B* are both above the lower bound, it is possible for the system to experience a decrease in fidelity under either type of jump.

involves quadratic equations, simple analytic solutions no longer exist. We find *all* three-state cycles by numerical search for all real solutions to Eqs. (5) and (6) using symbolicnumerical algorithms based on computing a Groebner basis [16,17], as described in detail in Appendix B.

One of the reasons to study K-state PR ensembles for K > 2 is the search for low entropy solutions that allow for efficient tracking. As we have shown, for $\varepsilon > 1/4$ there are no low entropy solutions for K = 2. However by moving to K = 3 we open the possibility for ensembles with different properties. The intuitive reason for the greater flexibility is as follows. Recall that for two-state jumping, we could use only one eigenvector of A to generate the PR states and we could not use complex eigenvectors because Bloch vectors must have real components. Observe that complex eigenvectors of A come in conjugate pairs. Thus, for three-state jumping, we can construct (real) Bloch vectors for PR states if we use conjugate pairs of eigenvectors of A. This method imposes additional constraints however and cannot be used for all possible ε . We now explore when complex eigenvectors can yield a PR ensemble.

We first show that any cyclic jumps between three states is generated from two eigenvectors of *A*. We observe that Eq. (6) implies that for $\vec{s}_j = \vec{r}_j - \vec{r}_{ss}$, with $j \in \{1,2,3\}$,

$$(A - \kappa_{12})(A - \kappa_{23})(A - \kappa_{31})\vec{s}_j = -\kappa_{12}\kappa_{23}\kappa_{31}\vec{s}_j.$$
 (62)

Since A is invertible, we can reformulate this condition as $g(A)\vec{s}_j = 0$, where we have defined a quadratic function $g(A) = A^2 - (\kappa_{12} + \kappa_{23} + \kappa_{31})A + (\kappa_{12}\kappa_{23} + \kappa_{12}\kappa_{13} + \kappa_{23}\kappa_{31})$. This is an eigenvalue equation for \vec{s}_j . Note that all eigenvectors of A are also eigenvectors of g(A) and vice versa. Thus \vec{s}_j is a linear combination of eigenvectors of A, whose eigenvalues λ satisfy equation $g(\lambda) = 0$. If all eigenvalues of A are distinct, then only two eigenvalues of A can satisfy equation $g(\lambda) = 0$ and only two eigenvectors of A are used to construct \vec{s}_j .

However, not every conjugate pair of complex eigenvectors of A can be used to construct \vec{s}_j . Suppose A has two

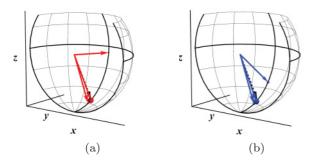


FIG. 9. (Color online) Solid arrows show Bloch vectors for three-state jumping. The volume of the sphere at the tip of each arrow represents the probability that the qubit occupies the corresponding pure state. The dashed arrow is \vec{r}_{ss} . Solutions in panels (a) and (b) are shown for $\varepsilon = 0.27$. These are low entropy solutions for three-state jumping that are generated from complex conjugate pairs of eigenvectors of *A*. Note that a low entropy solution fortwo-state jumping does not exist for such ε . The solution in panel (a) has entropy lower than that in panel (b).

complex eigenvalues λ and λ^* . Equation $g(\lambda) = 0$ implies that $\kappa_{12} + \kappa_{23} + \kappa_{31} = 2\text{Re}\lambda$ and $\kappa_{12}\kappa_{23} + \kappa_{12}\kappa_{13} + \kappa_{23}\kappa_{31} = |\lambda|^2$. By construction, κ_{12} , κ_{23} , and κ_{31} are real, which is only possible if $(\text{Re}\lambda)^2 > 3(\text{Im}\lambda)^2$. In particular, for the resonance fluorescence example, complex eigenvectors cannot be used to generate three-state PR states if $|\varepsilon| > 1/2$. This bound is actually not tight as was shown by numerical search for solutions to Eqs. (5) and (6) for three-state jumping [6]. These determined that such solutions exist iff $|\varepsilon| \leq 0.282$.

Solutions for cyclic three-state jumping come in pairs. As $|\varepsilon|$ approaches 0, the Shannon entropy h for half of the solutions approaches 1.2, whereas h for the other half approaches 0. For $0.247 < |\varepsilon| \leq 0.282$, there are two preferred ensembles generated from complex eigenvectors \vec{u}_+ of A and shown in Fig. 9 with $\varepsilon = 0.27$. In the region $0.183 < |\varepsilon| \leq$ 0.247, there are six solutions, which are shown in Fig. 10 with $\varepsilon = 0.23$. Solutions with the same entropy are shown in one subplot. These entropy-degenerate solutions are mirror images with respect to the x = 0 plane and are constructed from \vec{u}_1 and \vec{u}_- . The other two solutions have the lowest and the highest entropy and are generated by \vec{u}_+ and \vec{u}_- . In the last region, $|\varepsilon| < 0.183$, there are eight solutions, which are shown in Fig. 11 with $\varepsilon = 0.18$. Solutions with the same entropy still appear in the same subplot and are mirror images with respect to the x = 0 plane. All other solutions are constructed from \vec{u}_{+} and \vec{u}_{-} . Solutions that have unique entropy always lie in the x = 0 plane.

In ensemble $\{\wp_k, \vec{r}_k\}_{k=1,2,3}$, we can order the states by probability \wp_k in decreasing order. The system always jumps from one state to another in this order. For solutions with lower entropy, the systems spends most of the time in state one, which is nearly aligned with the steady state. For the lowest entropy ensemble, the states are spread out on the Bloch sphere as far away from each other as possible. As the entropy increases, states with small probability tend to move closer to the steady state. And for high entropy ensembles, the probability for occupying each state tends to equalize, and none of the states align with the steady state, but

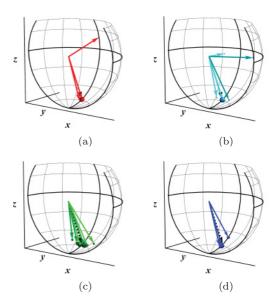


FIG. 10. (Color online) As in Fig. 9, but showing all three-state jumping solutions for $\varepsilon = 0.23$. The solution in panel (a) has the smallest entropy. Every consequent solution has a larger entropy. Solutions (a) and (d) correspond to solutions (a) and (b) in Fig. 9. Solutions (b) and (c) are actually pairs of solutions, symmetric about the x = 0 plane.

instead cluster around it. In Table I, for $\varepsilon = 0.15$ we report the geometric characteristics of the six distinct solutions. The total angle $\angle S_a = \angle(\vec{r}_1, \vec{r}_2) + \angle(\vec{r}_2, \vec{r}_3) + \angle(\vec{r}_3, \vec{r}_1)$ between the Bloch vectors shows inverse correlation with the entropy.

D. Stability of three-state jumping

We now analyze the stability of the above three-state jumping schemes. We do this using results from Secs. III C and III D and numeric solutions for three-state jumping computed in Sec. IV C. As shown before, the stability of an individual trajectory is determined by the properties of the effective Hamiltonian $\hat{H}(\mu_i)$ from Eq. (12), where index *i* ranges from 1 to 3 and μ_i are the settings for adaptive monitoring that generate three-state jumping for a resonance fluorescence system. We extract values for μ_i using the expression derived in Ref. [6] for cyclic jumps that states that

$$\hat{c}\left|v_{k}^{e}\right\rangle = -\mu_{k}\left|v_{k}^{e}\right\rangle + b_{k}\left|v_{k+1}^{e}\right\rangle,\tag{63}$$

TABLE I. This table reports geometric characteristics of six distinct solutions for $\varepsilon = 0.15$. It shows that the total angle (column 1) for the ensemble shows the inverse correlation with the entropy *h*.

$\angle S_a$	$\angle(\vec{r}_1,\vec{r}_{\rm ss})$	$\angle(\vec{r}_2,\vec{r}_{\rm ss})$	$\angle(\vec{r}_3,\vec{r}_{\rm ss})$	h
235.489	115.323	2.420 85	0.031 354 6	0.020
221.528	109.238	3.580 5	0.039 062 6	0.023
189.578	94.7316	5.874 93	0.057 596 5	0.026
26.934 5	0.654 795	12.812 4	5.303 14	0.466
14.243 5	1.663 5	5.457 59	2.779 01	1.171
13.061 4	3.065 05	1.486 3	3.465 69	1.299

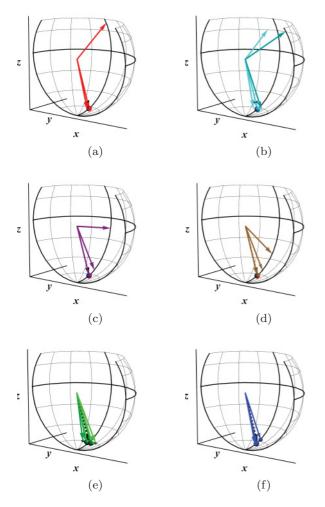


FIG. 11. (Color online) As in Fig. 10, but showing all three-state jumping solutions for $\varepsilon = 0.18$. Again, the solution in panel (a) has the smallest entropy and the entropy increases with every consequent solution. Solutions (a) and (b) correspond to solutions (a) and (b) in Fig. 10. Solutions (e) and (f) correspond to solutions (c) and (d) in Fig. 10.

where k + 1 stands for $k + 1 \mod 3$ and b_k is some constant. Once $|v_k^e\rangle$ and $|v_{k+1}^e\rangle$ are known, coefficients μ_k and b_k are uniquely determined. Thus our approach for determining the stability of three-state jumping proceeds as follows. We convert the three-state jumping solutions in terms of Bloch vectors \vec{r}_1 , \vec{r}_2 , and \vec{r}_3 into state vectors $|v_1^e\rangle$, $|v_2^e\rangle$, and $|v_3^e\rangle$. Using Eq. (63), we compute μ_k and $\hat{H}(\mu_k)$ with $k = 1, \ldots, 3$. At this point we can conclude that the evolution stage from state $|v_k^e\rangle$ to state $|v_{k+1}^e\rangle$ is stable if state $|v_k^e\rangle$ has the smallest eigenvalue with respect to operator $\hat{H}(\mu_k)$. Otherwise, such a stage is unstable. From the eigenvalues λ_j^e and λ_j^o one can determine mean-square stability using Eq. (41).

Results for this analysis are reported in Fig. 12 and resemble the results for two-state jumping quite closely. Figure 12 shows ensemble Shannon entropy h for all possible three-state jumping solutions. Instability of individual trajectories for three-state jumping always arises in the same way as for two-state jumping: one of the three stages (corresponding to jumping from the least probable state) becomes unstable. Solu-

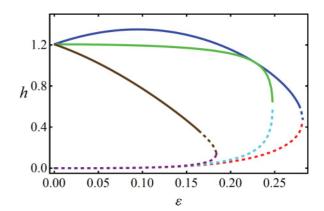


FIG. 12. (Color online) Ensemble Shannon entropy h for the eight different three-state jumping solutions (two of which come in degenerate pairs). The solid line indicates mean-square stability as well as piecewise deterministic stability. The dashed line indicates that the solution has only mean-square stability, with one of its stages being unstable.

tions with small entropy (i.e., solutions with $h \rightarrow 0$ as $\varepsilon \rightarrow 0$) have mean-square stability, but not piecewise deterministic stability. Solutions with large entropy have both mean-square and piecewise deterministic stability, except in small regions, where the entropy curves approach lower entropy solutions and develop one unstable stage in the evolution.

V. CONCLUSION

In this paper we considered a qubit undergoing evolution given by a Markovian master equation, subject to continuous monitoring that resolves every jump and allows the system to stay in a pure state. We studied special adaptive monitoring schemes that, in principle, allow an experimenter to track the evolution of such a system with a finite state machine as an apparatus. That is, the system jumps between only finitely many different states, the states in the associated *PR ensemble*. The main contribution of this paper beyond Ref. [6] is analysis of the stability of such monitoring schemes. This is necessary to establish that the finite PR ensembles introduced in Ref. [6] are not just mathematical constructions, but really are PR.

Because the evolution of the system is stochastic-it undergoes jumps at random times-we concentrated on the average properties of the system over many different realizations. Specifically, we derived conditions for the mean-square stability (that is, the long time convergence of the fidelity of the system with the correct ensemble state). We showed that there is a positive parameter C, which is never more than unity, that guarantees mean-square stability as long it is not equal to unity. For the specific example of resonance fluorescence we considered 11 different finite state monitorings (with two or three different states). In all cases C was strictly less than unity (by a long way, in fact). Based on this, we conjecture that all finite PR monitorings are mean-square stable. However we also considered the long time rate R of convergence (which is exponential in time) and found that some monitoring schemes converge much faster than others, in a way contrary to what is suggested by considering C alone.

Although all monitorings we considered were mean-square stable, there is a variety of behavior in the convergence of individual trajectories. Some monitorings give trajectories that are guaranteed to be piecewise deterministically stable (i.e., the between-jump evolution always increases the fidelity). However others do not—in some deterministic (between jump) stages of the evolution the system moves *away* from the ensemble state that it "should" be in at that stage (and will be in, in the long time limit). Moreover, the stochastic jumps can contribute to both stability and instability of individual trajectories, even under piecewise deterministically stable monitorings. However in the long time limit, as the system approaches the ideal ensemble states in the meansquare sense, a fidelity decrease upon a jump becomes very unlikely.

We illustrated these effects for the specific example of resonance fluorescence with an adaptively controlled weak local oscillator. Interestingly we found that those monitorings that are *not* piecewise deterministically stable tend to be those that are *most* stable in the mean-square sense (by the measure of *C* being very small). Moreover, these are the monitorings that produce low entropy solutions, where the system spends most of its time in one state in the ensemble. Such monitorings would make it possible, in principle, to track a large number *N* of qubits, using much less than *N* bits of memory (i.e., a finite state machine with far fewer than 2^N states).

It is important to note that there are many open questions in this field of quantum state tracking with finite state machines. First, given a PR finite ensemble, does there exist an explicit construction for an adaptive monitoring scheme that realizes it? Second, is any such monitoring mean-square stable, as conjectured here? Third, is it true that any D-dimensional ergodic Markovian quantum system can be tracked by a K-state machine with some finite $K \ge (D-1)^2 + 1$, as conjectured in Ref. [6]. Fourth, is $K = (D-1)^2 + 1$ always sufficient? Fifth, can an example system be found that proves that K = D is not sufficient in general (D = 3 would be)the minimum system size for such a search). If the last can be proven, then there may be a relation to the recent result that there are classical stochastic processes that can be generated using quantum systems of entropy lower than that required using only classical systems [18]. Finally, it seems likely that any given master equation would have additional structure that would enable one to use a K smaller than that conjectured above, and this idea also remains to be explored.

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APPENDIX A: USEFUL IDENTITIES

We now explain how to derive expressions for the convergence coefficient C that determines how quickly the state subject to adaptive monitoring approaches the states in the

corresponding PR ensemble. We begin with the two-state jumping scenario. In this case, the expression for *C* is given in Eq. (33). We evaluate it through relating $Q_{11}^1 Q_{11}^2$ and $Q_{22}^1 Q_{22}^2$. We do this by comparing the general expression $\int d^n \vec{\tau} |\beta_n|^2$ to $\int d^n \vec{\tau} |\alpha_n|^2$ for the special case when the initial state $|\psi_0\rangle$ is the jumping stat $|v_1^e\rangle$, i.e., $\alpha_0 = 1$ and $\beta_0 = 0$. Then with Eqs. (23)–(25) and (15), we get

$$\alpha_{2l+1} = \alpha_{2l} e^{-\lambda_1^e \tau_{2l+1}} Q_{11}^1, \tag{A1}$$

$$\alpha_{2l} = \alpha_{2l-1} e^{-\lambda_2^e \tau_{2l}} Q_{11}^2.$$
 (A2)

Again by induction we deduce that

$$\alpha_{2l} = \left(Q_{11}^1 Q_{11}^2\right)^l \prod_{j=1}^l e^{-\lambda_1^e \tau_{2j-1}} e^{-\lambda_2^e \tau_{2j}}, \qquad (A3)$$

$$\alpha_{2l+1} = Q_{11}^1 (Q_{11}^1 Q_{11}^2)^l \prod_{j=0}^l e^{-\lambda_2^e \tau_{2j}} e^{-\lambda_1^e \tau_{2j+1}}.$$
 (A4)

Then

$$\int d^{n}\vec{\tau}|||\tilde{\psi}_{2l}\rangle||^{2} = \int d^{n}\vec{\tau}|\alpha_{2l}|^{2} = \left(\frac{\left|Q_{11}^{1}Q_{11}^{2}\right|^{2}}{4\mathrm{Re}\lambda_{1}^{e}\mathrm{Re}\lambda_{2}^{e}}\right)^{l}.$$
 (A5)

By construction $\int d^{2l} \vec{\tau} |||\tilde{\psi}_{2l}\rangle||^2 = 1$. Therefore, Eq. (A5) proves that

$$\frac{|Q_{11}^1 Q_{11}^2|^2}{4\text{Re}\lambda_1^e \text{Re}\lambda_2^e} = 1.$$
 (A6)

To calculate the value of *C*, we relate $Q_{11}^1 Q_{11}^2$ to $Q_{22}^1 Q_{22}^2$. We recall the discussion after Eq. (17), which shows that operator $\hat{s}_2 \hat{s}_1$ is proportional to the identity operator. Using this fact and Eqs. (25), we see that

$$\hat{s}_2 \hat{s}_1 | v_1^e \rangle = Q_{11}^1 Q_{11}^2 | v_1^e \rangle, \tag{A7}$$

$$\hat{s}_2 \hat{s}_1 | v_1^o \rangle = Q_{22}^1 Q_{22}^2 | v_1^o \rangle. \tag{A8}$$

Since $\hat{s}_2 \hat{s}_1$ is proportional to the identity, we can conclude $Q_{11}^1 Q_{11}^2 = Q_{22}^1 Q_{22}^2$.

This discussion can be easily generalized to the *K*-state jumping scenario. We again calculate the value of $\int d^{nK} \vec{\tau} |\alpha_{nK}|^2$ in two ways when $\alpha_0 = 1$ and $\beta_0 = 0$. The first way is to note that, as before, in this case $\int d^{nK} \vec{\tau} |\alpha_{nK}|^2 = \int d^{nK} \vec{\tau} |||\tilde{\psi}_{nK}\rangle||^2 = 1$. The second way is to use the fact that in this case

$$\alpha_{nK} = \left(Q_{11}^1 \cdots Q_{11}^K\right)^n \prod_{j=0}^{n-1} e^{-\lambda_1^{e} \tau_{jK+1}} \cdots e^{-\lambda_K^{e} \tau_{(j+1)K}}, \quad (A9)$$

so that

$$\int d^{nK} \vec{\tau} |\alpha_{nK}|^2 = \left(\frac{Q_{11}^1 \cdots Q_{11}^K}{2^n \operatorname{Re} \lambda_1^e \cdots \operatorname{Re} \lambda_K^e}\right)^n.$$
(A10)

Thus $Q_{11}^1 \cdots Q_{11}^K = 2^n \operatorname{Re} \lambda_1^e \cdots \operatorname{Re} \lambda_K^e$. Next, we show that $Q_{11}^1 \cdots Q_{11}^K = Q_{22}^1 \cdots Q_{22}^K$. We prove this relationship for Q_{kk}^j by considering the full cycle jump operator $\hat{S}_K = \hat{s}_K \cdots \hat{s}_1$.

Assuming (as we can do without loss of generality) that the jump operator \hat{c} is traceless, we have $\hat{c}^2 \propto \hat{I}$ so that

$$\hat{S}_{K} \left| v_{1}^{e} \right\rangle = \left[g_{1}(\mu_{1}, \dots, \mu_{K}) + g_{2}(\mu_{1}, \dots, \mu_{K}) \hat{c} \right] \left| v_{1}^{e} \right\rangle.$$
(A11)

The exact nature of functions g_1 and g_2 depends on the proportionality constant relating \hat{c}^2 to the identity operator. Just as for Eq. (17), a system can undergo jumping dynamics iff $g_2 = 0$. This means that the full cycle jump operator \hat{S}_K is proportional to the identity and we conclude that $Q_{11}^1 \cdots Q_{11}^K = Q_{22}^1 \cdots Q_{22}^K$, using the trick from Eqs. (A7) and (A8).

APPENDIX B: GROEBNER BASIS

The task of finding jumping states for adaptive unravelings requires solving nonlinear equations. We now introduce important concepts from computational algebraic geometry and review some algorithms used to solve multivariate polynomial systems. Our presentation will rely on analogies with linear algebra. From now on we regard a polynomial as the finite sum of terms, where each term is a product of a coefficient and monomial.

Suppose we want to solve the system of nonlinear equations,

$$\{f_1(\vec{x}_n) = 0, f_2(\vec{x}_n) = 0, \dots, f_s(\vec{x}_n) = 0\},$$
 (B1)

where $\{f_1, f_2, \ldots, f_s\}$ are the polynomials with real rational coefficients and $\vec{x}_n = (x_1, x_2, \ldots, x_n)$ is the list of variables. Our goal is to find all the solutions to the set of equations in Eq. (B1) and a concept of *ideal* becomes useful. A collection of polynomials generates the ideal via $I = \langle f_1, f_2, \ldots, f_n \rangle = \{\sum_{i=1}^s h_i f_i : h_1, \ldots, h_s \in \mathbb{C}[\vec{x}_n]\}$, where *s* can be any finite index and $\mathbb{C}[\vec{x}_n]$ is a collection of all possible polynomials with complex coefficients with variables \vec{x}_n . Thus ideals are similar to vector spaces, which are formed from all possible scalar combinations of vectors. But instead of scalars, one uses all possible polynomial functions, h_k , defined on $\mathbb{C}[\vec{x}_n]$ to form an ideal. Ideals are important because the solution set to the newly created ideal and to the original system are the same.

Different sets of equations can have the same solution set and thus generate the same ideal. One of the main ideas in algebraic geometry is to pick a good set representing the ideal that has nice properties and yields an easy way to identify the solutions. Such a set is called a *Groebner basis*.

Gaussian elimination is the algorithm used to solve a system of linear equations. The extension to this algorithm used to solve a system of polynomial equations is known as Buchberger's algorithm and it is implemented in many packages for symbolic computations such as MATHEMATICA, MAPLE, and SAGE. The set of equations obtained as a result of these algorithms is known as the Groebner basis.

Before proceeding we explain how to check if a given set of equations produced by some software package is indeed a Groebner basis. To do this, we first need to know the *S polynomial*. Given two polynomials *f* and *g*, let \mathbf{x}^{α} be the least common multiple of leading terms of *f* and *g*, denoted by L(f) and L(g). Then the *S* polynomial is computed via $S(f,g) = f\mathbf{x}^{\alpha}/L(f) - g\mathbf{x}^{\alpha}/L(g)$. A set of polynomials G = $\{g_1, \ldots, g_n\}$ are a Groebner basis iff for all $i \neq j$ the remainder on division of $S(g_i, g_j)$ by *G* is 0. An alternative explanation proceeds as follows. In general, one can always write $S(g_i, g_j)$ as $S(g_i, g_j) = \sum_{i=1}^n a_i g_i + r$, where a_i are some polynomials and the degree of the remainder polynomial r is smaller than that of any polynomial in G. Then G is a Groebner basis iff the remainder polynomial r is 0 for all $i \neq j$. Division by a collection of polynomials is usually implemented in the software packages for computational algebraic geometry.

Now we explain how to use a Groebner basis to find *all* solutions to the system of nonlinear equations. Unlike linear algebra, where a vector space always has the same number of basis vectors, different Groebner bases can have different numbers of elements and drastically different properties for the same ideal. The collection of Groebner bases arise from different orderings of monomials in the system of polynomials. Two particular orderings are relevant for the current discussion: lexicographic order (Lex) and degree reverse lexicographic order (DRL). Lex order is an alphabetical order (write out monomial in full without any powers and order like the words in a dictionary from left to right). For DRL order, we first compare the total degree and then perform lexicographic ordering by reading the expressions from right to left. Both orderings allow different access to information about solutions to polynomial systems of equations. The details on different orderings can be found in Ref. [16].

Lex order is particularly useful because of the elimination theorem [16], which states that the Groebner basis computed with respect to Lex order will yield a set of polynomials that can be solved by back substitution. This means that if the system of nonlinear equations has a finite number of solutions and the Lex Groebner basis is given by $G = \{g_1, \ldots, g_t\}$, then $g_1 = g_1(x_1)$ is a polynomial function of one variable only and we can determine all values for x_1 by solving $g_1(x_1) = 0$. The next polynomial g_2 is mostly like a function of two variables, x_1 and x_2 , so that we can solve for x_2 given all possible values for x_1 . This sequential substitution allows one to solve for all variables \vec{x}_n . Some of the elements of *G* are just constraints that eliminate possible solutions. Such a procedure is used to solve a system of nonlinear equations in MATHEMATICA (command NSolve).

This algorithm comes with some complications. For some systems, the algorithm becomes unstable. This is so because the degree and size of coefficients in the Lex Groebner basis quickly become huge even for relatively small systems. And solving $g_1(x_1) = 0$ involves rounding, which introduces errors that quickly propagate through back substitution and yield wrong results.

The Groebner basis computed with respect to DRL ordering offers an alternative way to compute solutions to a polynomial system of equations that does not involve back substitution. We now review machinery from algebraic geometry needed to compute solutions using DRL ordering of monomials. The first construct involves the space of all polynomials $\mathbb{C}[\vec{x}_n]$ and ideal *I* defined above. Polynomials $\mathbb{C}[\vec{x}_n]$ can be classified into distinct categories (cosets) such that polynomials *f* and *g* belong to same coset iff *f* and *g* have 0 remainder with respect to ideal *I* or equivalently polynomials *f* and *g* have the same remainder with respect to *I*. We denote the reminder of *f* by [f]. Note that for checking this criteria it is sufficient just to consider the remainder with respect to the Groebner basis (same ordering as the one used for division algorithm) spanning *I*.

We can consider the space of cosets, i.e., the space of the remainders with respect to division by *I*. This is a quotient ring, denoted by $\mathcal{A} = \mathbb{C}$ [\vec{x}_n]. It also happens to be an algebra that has vector space structure and carries important information about solutions to the system of polynomials that spanned the ideal *I*.

We now point out the key steps in computing solutions to polynomials $\{f_1, f_2, \ldots, f_n\}$ and illustrate them using a simple example.

(i) We compute the Groebner basis, $G = \{g_1, \dots, g_t\}$, with respect to DRL ordering for the ideal $I = \langle f_1, f_2, \dots, f_n \rangle$.

(ii) We identify the leading term for every element of the Groebner basis $L(g_i)$.

(iii) We create set *B* from monomials that a not divisible by $L(g_i)$, i.e., monomials in *B* have a degree smaller than the degree of $L(g_i)$. This set is a basis for the space of remainders, A.

For example, suppose we want to solve a system of equations:

$$f_1 = x^2 - y^2 + xy = 0,$$

$$f_2 = x^2y + y - 1 = 0.$$
(B2)

The Groebner basis (DRL ordering and x > y) for this system is

$$g_1 = x^2 + xy - y^2,$$

$$g_2 = y^3 - xy^2 + y - 1,$$

$$g_3 = y^4 + xy + 2y^2 - x - 2y.$$
(B3)

Leading terms for these polynomials (with respect to chosen ordering) are $L(I) = \{x^2, xy^2, y^4\}$ and the set $B = \{1, x, y, xy, y^2, y^3\}$ is the basis for the quotient ring.

To extract the information about solutions from the quotient ring \mathcal{A} , we associate each polynomial f in $\mathbb{C}[\vec{x}_n]$ with a linear map $m_f : \mathcal{A} \to \mathcal{A}$ whose action is given by $m_f(g) = [fg]$, where g is some polynomial. Map m_f can be represented as a matrix. To do this, we consider the action of m_f on basis elements in the set B. And $m_f(b)$ for every $b \in B$ is an element of quotient ring \mathcal{A} and can be written as a column vector with respect to B.

For example discussed above, we show how to compute m_x . The action of m_x on 1 is given by

$$m_x(1) = [x] = (0, 1, 0, 0, 0, 0)^T,$$
 (B4)

i.e., m_x is the coset [x], which can be represented as a vector with respect to elements in basis set *B*. Map m_x will take some basis elements outside of *B*. In this case, we compute the remainder with respect to the Groebner basis and again compute a column vector with respect to *B*. For

example,

$$m_x(x) = [x^2] = [g_1 + y^2 - xy] = [y^2 - xy]$$

= (0,0,0, -1,1,0)^T. (B5)

Then the matrix form for m_x is

$$m_{x} = \begin{pmatrix} 0 & 0 & 0 & 1 & -1 & 0 \\ 1 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & -1 & 1 & 1 \\ 0 & -1 & 1 & 0 & 0 & -1 \\ 0 & 1 & 0 & 0 & 0 & -1 \\ 0 & 0 & 0 & 1 & 1 & 0 \end{pmatrix}.$$
 (B6)

We need to compute these matrices because they have a very nice property: eigenvalues of m_f are the values polynomial ftakes, when evaluated on the solution set. Thus eigenvalues of m_x are all possible values a variable x takes in the solution set. Moreover, we don't need to compute every possible matrix m_f because $m_{fg} = m_f m_g$. This is significant since computing m_f relies on costly symbolic calculations, which are slow, whereas matrix multiplication is fast.

Thus eigenvalues for all different maps m_{x_i} tell us all possible values each variable x_i can take. But to learn the full solution (how to combine different x_i to form \vec{x}_n) we also need information that is encoded in the eigenvector of m_f for some f. Let **v** be such an eigenvector and we normalize this vector so that the first component is 1 (assuming that 1 is the first element in B). If variable x_i is the *j*th element in set B, then the *j*th component of **v** is the value of x_i that yields one of the solutions. All values extracted in this way from one eigenvector come from one solution. However, this procedure does not work for every polynomial f. If m_f has degenerate eigenvalues then corresponding eigenvectors cannot be used to extract the solution. However, generic linear combination of the variables will yield a desired matrix. The details on this can be found in Refs. [17,19].

For three-state jumping, we computed the solutions using the Groebner basis with DRL order. The set of equations for cyclic three-state jumping has a lot of symmetry, e.g., one can map $\vec{r}_1 \mapsto \vec{r}_2 \mapsto \vec{r}_3 \mapsto \vec{r}_1$ and $\kappa_{12} \mapsto \kappa_{23} \mapsto \kappa_{13} \mapsto \kappa_{12}$ and the system will remain unchanged. As the result, all matrices m_{x_i} associated with unknowns in the system have degeneracy with respect to real eigenvectors and cannot be used to extract the solution. Instead we used matrix m_f , where $f = r_{11} + \kappa_{12}$. Here r_{11} is the first component of the Bloch vector \vec{r}_1 . The reason for this choice is the following. Matrices associated with κ_{12} , κ_{23} , and κ_{31} are nice because they have the least nonzero elements, which speeds up the calculation (symbolic part). However, any combination of κ_{12} , κ_{23} , and κ_{31} or the same coordinate of r_1 , r_2 , and r_3 does not break the degeneracy due to symmetry above, but $r_{11} + \kappa_{12}$ does.

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