

# Using a quantum computer to investigate quantum chaos

Rüdiger Schack\*

Department of Mathematics, Royal Holloway, University of London, Egham, Surrey TW20 0EX, United Kingdom

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We show that the quantum baker's map, a prototypical map invented for theoretical studies of quantum chaos, has an efficient realization in terms of quantum gates. Chaos in the quantum baker's map could be investigated experimentally on a quantum computer based on only three quantum bits.  
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Since the discovery that a quantum computer can in principle factor large integers in polynomial time [1,2], quantum information has become a major theoretical and experimental research topic, focusing on properties, applications, generation, and preservation of highly entangled quantum states [3]. Although it is not clear if a full-scale quantum computer will ever be realized [4,5], experiments with quantum gates are being performed at present [6–9]. It is important to devise applications for early quantum computers that are incapable of large-scale computations such as factoring.

Early quantum computers appear to be well suited to study the dynamics of simple quantum maps. The quantum baker's map [10], one of the simplest quantum maps used in quantum chaos research, has been extensively studied in recent years [11–16]. Hannay *et al.* [17] have proposed a realization of the baker's map in classical optics, which allows them to understand its quantization in terms of the relationship between ray and wave optics. Here we describe a genuine quantum system whose dynamics is governed by the quantum baker's map. As a consequence of recent progress in the field of quantum computing [6–9], an experimental realization of the quantum baker's map seems possible in the near future.

Any unitary operator can be approximated by a sequence of simple quantum gates [18–20]. The main result of this paper is that the quantum baker's map can be realized in terms of quantum gates in a particularly simple and efficient way. Similar to the quantum Fourier transform, simulating the quantum baker's map on a quantum computer is exponentially faster than a simulation on a classical computer.

The quantum baker's map displays behavior of fundamental interest even for a Hilbert space of small dimension. Numerical simulations [13] in  $D=16$  dimensional Hilbert space suggest that a rudimentary quantum computer based on as few as three quantum bits (qubits) (i.e., three two-state systems spanning  $D=8$  dimensional Hilbert space) could be used to study chaos in the quantum baker's map. In particular, it may be possible to find experimental evidence for hypersensitivity to perturbation, a proposed information-theoretical characterization of quantum chaos [13,21–23].

The classical baker's transformation [24] maps the unit square  $0 \leq q, p \leq 1$  onto itself according to

$$(q, p) \mapsto \begin{cases} (2q, \frac{1}{2}p) & \text{if } 0 \leq q \leq \frac{1}{2} \\ [2q - 1, \frac{1}{2}(p+1)] & \text{if } \frac{1}{2} < q \leq 1. \end{cases} \quad (1)$$

This corresponds to compressing the unit square in the  $p$  direction and stretching it in the  $q$  direction, while preserving the area, then cutting it vertically, and finally stacking the right part on top of the left part, in analogy to the way a baker kneads dough.

To define the quantum baker's map [10], we quantize the unit square following [11,25]. To represent the unit square in  $D$ -dimensional Hilbert space, we start with unitary “displacement” operators  $\hat{U}$  and  $\hat{V}$ , which produce displacements in the “momentum” and “position” directions, respectively, and obey the commutation relation [25]

$$\hat{U}\hat{V} = \hat{V}\hat{U}\epsilon, \quad (2)$$

where  $\epsilon^D = 1$ . We choose  $\epsilon = e^{2\pi i/D}$ . We further assume that  $\hat{V}^D = \hat{U}^D = 1$ , i.e., periodic boundary conditions. It follows [11,25] that the operators  $\hat{U}$  and  $\hat{V}$  can be written as

$$\hat{U} = e^{2\pi i \hat{q}}, \quad \hat{V} = e^{-2\pi i \hat{p}}. \quad (3)$$

The position and momentum operators  $\hat{q}$  and  $\hat{p}$  both have eigenvalues  $j/D$ ,  $j=0, \dots, D-1$ .

In the following, we restrict the discussion to the case  $D=2^L$ , i.e., the dimension of Hilbert space is a power of 2. For consistency of units, let the quantum scale on “phase space” be  $2\pi\hbar = 1/D = 2^{-L}$ . A transformation between the position basis  $\{|q_j\rangle\}$  and the momentum basis  $\{|p_j\rangle\}$  is effected by the discrete Fourier transform  $F'_L$ , defined by the matrix elements

$$(F'_L)_{kj} = \langle p_k | q_j \rangle = \sqrt{2\pi\hbar} e^{-ip_k q_j / \hbar} = \frac{1}{\sqrt{D}} e^{-2\pi i k j / D}. \quad (4)$$

There is no unique way to quantize a classical map. Here we adopt the quantized baker's map introduced by Balazs and Voros [10] and defined by the matrix

$$T' = F'_L{}^{-1} \begin{pmatrix} F'_{L-1} & 0 \\ 0 & F'_{L-1} \end{pmatrix}, \quad (5)$$

where the matrix elements are to be understood relative to the position basis  $\{|q_j\rangle\}$ . Saraceno [11] has introduced a quantum baker's map with stronger symmetry properties by using antiperiodic boundary conditions, but in this article we restrict the discussion to periodic boundary conditions as used in [10].

\*Electronic address: r.schack@rhbc.ac.uk

The discrete Fourier transform used in the definition of the quantum baker's map (5) plays a crucial role in quantum computation and can be easily realized as a quantum network using simple quantum gates. The following discussion of the quantum Fourier transform follows [2] closely. The  $D=2^L$  dimensional Hilbert space modeling the unit square can be realized as the product space of  $L$  qubits (i.e.,  $L$  two-state systems) in such a way that

$$|q_j\rangle = |j_{L-1}\rangle \otimes |j_{L-2}\rangle \otimes \cdots \otimes |j_0\rangle, \quad (6)$$

where  $j = \sum j_k 2^k$ ,  $j_k \in \{0,1\}$  ( $k=0, \dots, L-1$ ), and each qubit has basis states  $|0\rangle$  and  $|1\rangle$ .

To construct the quantum Fourier transform, two basic unitary operations or *quantum gates* are needed: the gate  $A_m$  acting on the  $m$ th qubit and defined in the basis  $\{|0\rangle, |1\rangle\}$  by the matrix

$$A_m = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad (7)$$

and the gate  $B_{mn}$  operating on the  $m$ th and  $n$ th qubits ( $m < n$ ) and defined by

$$B_{mn}|j_{L-1}\rangle \otimes \cdots \otimes |j_0\rangle = e^{i\phi_{mn}} |j_{L-1}\rangle \otimes \cdots \otimes |j_0\rangle, \quad (8)$$

where

$$\phi_{mn} = \begin{cases} \pi/2^{n-m} & \text{if } j_m = j_n = 1 \\ 0 & \text{otherwise.} \end{cases} \quad (9)$$

In addition we define the gate  $S_{mn}$  that swaps the qubits  $m$  and  $n$ .

The discrete Fourier transform  $F_L$  can now be expressed in terms of the three types of gates as

$$F_L = S \times (A_0 B_{01} \cdots B_{0,L-1}) \times \cdots \times (A_{L-3} B_{L-3,L-2} B_{L-3,L-1}) \\ \times (A_{L-2} B_{L-2,L-1}) \times (A_{L-1}), \quad (10)$$

where

$$S = \begin{cases} S_{0,L-1} S_{1,L-2} \cdots S_{L/2-1,L/2} & \text{for } L \text{ even} \\ S_{0,L-1} S_{1,L-2} \cdots S_{(L-3)/2,(L+1)/2} & \text{for } L \text{ odd} \end{cases} \quad (11)$$

reverses the order of the qubits. The quantum baker's map (5) is then given by

$$T = F_L^{-1} (I \otimes F_{L-1}), \quad (12)$$

where  $F_{L-1}$  acts on the  $L-1$  least significant qubits and  $I$  is the identity operator acting on the most significant qubit. The gates corresponding to the bit-reversal operator  $S$  can be omitted if the qubits in the tensor product (6) are relabeled after each execution of  $F_L$  or  $F_{L-1}$ .

In  $D=8=2^3$  dimensional Hilbert space, one iteration of the quantum baker's map is performed by the short sequence of gates

$$T = S_{02} A_0 B_{01}^\dagger B_{02}^\dagger A_1 B_{12}^\dagger A_2 S_{01} A_0 B_{01} A_1. \quad (13)$$

This implementation of the quantum baker's map can be viewed in two complementary ways. On the one hand, it shows that the quantum baker's map can be efficiently simulated on a quantum computer. A 30-qubit quantum computer could perform simulations that are virtually impossible on present-day classical computers. On the other hand, an iteration of the gate sequence (12) on an  $L$ -qubit quantum computer is a physical realization of the quantum baker's map. This opens up the possibility of an experimental investigation of chaos in a physical system in a purely quantum regime.

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