Algorithm for solving unconstrained unitary quantum brachistochrone problems

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We introduce an iterative method to search for time-optimal Hamiltonians that drive a quantum system between two arbitrary, and in general mixed, quantum states. The method is based on the idea of progressively improving the efficiency of an initial, randomly chosen Hamiltonian, by reducing its components that do not actively contribute to driving the system. We show that our method converges rapidly even for large-dimensional systems and that its solutions saturate any attainable bound for the minimal time of evolution. We provide a rigorous geometric interpretation of the iterative method by exploiting an isomorphism between geometric phases acquired by the system along a path and the Hamiltonian that generates it. Our method is directly applicable as a powerful tool for state preparation and gate design problems.

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I. INTRODUCTION

Inspired by the problem of finding the curve of fastest descent between two points, quantum brachistochrone problems (QBPs) aim to find the Hamiltonian that generates the timeoptimal evolution between two given quantum states. Such problems have been considered to obtain accurate minimumtime protocols for the control of quantum systems [1,2], clarify the role of entanglement and quantum correlations in time-optimal evolution [3,4], study the speed of Hermitian and non-Hermitian Hamiltonians [5–7], and improve bounds on the minimal time of evolution, known as quantum speed limits (QSLs) [8-10]. QBPs have a special role in quantum information theory and technology, where they are of fundamental importance to accurately perform tasks such as preparing a desired state of a system or implementing a specific gate, while satisfying the strict physical and fault tolerance requirements imposed by the locality of interactions and short decoherence times [11,12]. For this reason, their solutions have found applications in information processing [13–18], quantum state preparation [19–24], cooling [1,25], metrology [26–28], and quantum thermodynamics [29–32]. Moreover, QBPs give a physical interpretation to the complexity of quantum algorithms, which emerges from the minimization of the time required to obtain the desired unitary transformation [33].

Solving QBPs is generally hard, and accurate analytic and numerical solutions are only known for some special cases, such as the unconstrained unitary evolution of pure states and control problems of well-structured quantum systems [9,12,33–37]. There are methods that can be used to address a relatively large class of QBPs, but obtaining precise solutions becomes increasingly challenging as the dimension of the system grows and the constraints on the control Hamiltonian become more complex. For instance, the quantum brachistochrone equations proposed in Ref. [33] involve solving ordinary differential equations with boundary conditions, for which there are no efficient numerical methods when high-dimensional systems are considered [12,29,38–40]. One crucial open problem is the case of unconstrained unitary evolution between two mixed states. As opposed to the case of pure states, the solution to this QBP is not known, except for special cases with highly degenerate spectra. This is akin to the problem of finding tight bounds on the minimal time of evolution, known as QSL. These well-known bounds are attainable for pure states but are often loose for mixed states because of the complicated structure of the space of density operators [41–43], though we have recently proposed geometric methods to derive simple, efficient, and tight bounds for the time of evolution of mixed states [10,44].

In this paper, we take a similar approach to solving the complementary problem, that of finding the optimal unconstrained unitary evolution between two mixed states. That is, we look for the generator H of a unitary operator $U = \exp(-iHt)$ that takes a mixed state ρ to $\sigma = U \rho U^{\dagger}$, such that the transition time t is minimized ($\hbar \equiv 1$). We introduce an iterative method to search for the optimal time-independent Hamiltonian, while respecting an energetic constraint. The method progressively improves the efficiency of the Hamiltonian [45] that generates the evolution, by applying an algorithm that rapidly converges even for large-dimensional systems.

We investigate the efficacy of our method using QSLs. By comparing the achievable upper bound, provided by the iterative method, with the inviolable lower bound offered by QSLs, we demonstrate the strong synergy of the two results: When the two are found to be close, we can be sure that both results are close to optimal. We study the performance of our method with respect to the size of the system and its dependence on parameters such as the convergence threshold, before discussing direct applications of the algorithm, its potential combination with time-optimal gate design and Monte Carlo methods, and its geometric interpretation, juxtaposing it to that of Grover's quantum search algorithm.

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II. TIME-OPTIMAL EVOLUTION AND HAMILTONIAN EFFICIENCY

Let us begin by considering the QBP for a two-level system, defined by the initial state $\rho = (\mathbb{1}+p\Lambda_x)$ and the target state $\sigma = (\mathbb{1}+p\Lambda_y)/2$. Here, $p \in (0, 1]$ and $\mathbf{\Lambda} = (\Lambda_x, \Lambda_y, \Lambda_z)$ are Pauli matrices. The Bloch vectors for the two states are given by $\mathbf{r} = (p, 0, 0)$ and $\mathbf{s} = (0, p, 0)$, respectively. Any unitary $O(\boldsymbol{\varphi}) = e^{i\varphi_1}|s_1\rangle\langle r_1| + e^{i\varphi_2}|s_2\rangle\langle r_2|$ can be used to map $\rho \to \sigma$, where $|r_k\rangle$ and $|s_k\rangle$ are the eigenvectors of ρ and σ , respectively, and where $\boldsymbol{\varphi} = (\varphi_1, \varphi_2)$ represent the different geometric phases that the state gathers from $O(\boldsymbol{\varphi})$. The Hamiltonians $H(\boldsymbol{\varphi}) = i \ln O(\boldsymbol{\varphi})$ define different evolutions depending on the choice of $\boldsymbol{\varphi}$.

Let us focus on two possible choices for these phases, $\varphi_z = (\pi/4, \pi/4)$ and $\varphi_{xy} = (3\pi/4, -\pi/4)$. The associated Hamiltonians are given by $H(\varphi_z) = \Lambda_z \pi/4$ and $H(\varphi_{xy}) = (\Lambda_x + \Lambda_y)\sqrt{2}\pi/4$. These two Hamiltonians generate different paths on the Bloch sphere. Under energetic constraints, such as fixing the standard deviation $\Delta H_\rho = \sqrt{\text{tr}[\rho H^2] - \text{tr}[\rho H]^2}$ with respect to the initial state, the lengths of these paths meaningfully correspond to evolution times [46]. The path generated by $H(\varphi_z)$ connects ρ to σ with an arc of great circle, i.e., the geodesic with respect to the Fubini-Study (FS) metric [47], and thus it constitutes a solution to the considered QBP for any homogeneous energy constraint [8], whereas $H(\varphi_{xy})$ draws a *longer* path, which is not time optimal.

A heuristic explanation for the variable performance of different Hamiltonians is that the Bloch vectors h associated with them $H = h \cdot \Lambda$ have a different orientation with respect to rand s. In particular, when h is orthogonal to r, it generates a rotation on a plane that passes through the origin of the Bloch sphere, while when h is not perpendicular to r, it generates a rotation on a plane that does not. The slower evolution can be seen as due to a less efficient use of the resources encoded in constraints on the energy, which are *wasted* on parts of the Hamiltonian that not actively drive the system [45].

When pure states are considered, we refer to the the notion of Hamiltonian efficiency introduced in Ref. [48], given by

$$\eta(H,\rho) := \frac{\Delta H_{\rho}}{\|H\|_{op}},\tag{1}$$

with $\|\cdot\|_{op}$ being the operator norm. It is possible to consider different energetic constraints, such as $\|H\|_{HS} = \omega$, known as finite-energy bandwidth, or $\|H\|_{op} = \omega$, which corresponds to bounding the largest eigenvalue of the Hamiltonians. Different notions of efficiency can be obtained for any choice of speed and total energy measure. The efficiency measures how much of the energy available, quantified by $\|H\|_{op}$, is transferred to the system to drive the evolution, and thus *converted* into speed ΔH_{ρ} . If we assess the performance of the two Hamiltonians $H(\varphi_z)$ and $H(\varphi_{xy})$ considered above using η , we obtain $\eta(H(\varphi_z), \rho) = p$ and $\eta(H(\varphi_{xy}), \rho) = p/2$. The dependence on p (and thus on the purity of the states) implies that this notion of efficiency cannot be saturated in general. A generalization of η that can be saturated for arbitrary states is introduced in Sec. VIII.

This intuitive geometric argument can be generalized for dimension d > 2 by replacing the notion of orthogonality between vectors with commutation relations between states

and Hamiltonians. Given an operator ρ , the space of Hamiltonians [the Lie algebra $\mathfrak{su}(d)$] splits into a maximal dimensional *parallel* subspace, closed under multiplication, that commutes with ρ , and an orthogonal *perpendicular* subspace, every element of which does not. This allows us to decompose the Hamiltonian *H* into components H_{\parallel} and H_{\perp} which are elements, respectively, of these two subspaces, such that $[H_{\parallel}, \rho] = 0$ and $[H, \rho] = [H_{\perp}, \rho]$ [49]. This observation leads to the idea at the core of the iterative method that we will introduce in the next section, where efficient Hamiltonians for a QBP are achieved by requiring their parallel component to be vanishing at all times during the evolution.

III. ITERATIVE METHOD FOR EFFICIENT HAMILTONIANS

We can now consider the more general QBP defined by an arbitrary isospectral pair of initial and final states $\rho = \sum_k \lambda_k |r_k\rangle \langle r_k|$ and $\sigma = \sum_k \lambda_k |s_k\rangle \langle s_k|$ of finite dimension *d*. When nondegenerate states are considered, the operator

$$O(\boldsymbol{\varphi}) = \sum_{k} e^{i\varphi_k} |s_k\rangle \langle r_k| \tag{2}$$

represents the most general unitary that connects initial and final states ρ and σ , while φ_k represent the geometric phases gained evolving along the path generated by $i\ln O(\varphi)$. When degenerate states are considered, the phases φ_k are replaced by unitary operators $U_k \in SU(m)$ on the subspace associated with eigenvalues λ_k with multiplicity *m*. Since any unitary $O(\varphi)$ maps $\rho \to \sigma$, we can choose an arbitrary initial phase $\varphi^{(0)}$ to obtain the initial unitary $O^{(0)} = O(\varphi^{(0)})$. The Hamiltonian $H^{(0)} = i\ln O^{(0)}$ [50], canonically associated with $O(\varphi^{(0)})$, is then split into parallel and perpendicular components $\mathcal{M}_{\rho}[H^{(0)}] = H_{\parallel}^{(0)}$ and $H_{\perp}^{(0)} = H^{(0)} - H_{\parallel}^{(0)}$, such that $[H_{\parallel}^{(0)}, \rho] = 0$, via a map \mathcal{M}_{ρ} that depends on ρ , and which will be referred to as *mask*. We chose the mask to be

$$\mathcal{M}_{\rho}[H] = D(M \circ D^{\dagger}HD)D^{\dagger}, \text{ with } M_{ij} = \delta_{\lambda_i\lambda_j}, \quad (3)$$

where *D* diagonalizes ρ , i.e., its columns are given by the eigenstates $\{|r_k\rangle\}$ of ρ , and where \circ is the Hadamard, i.e., element-wise, product [51]. This mask projects the Hamiltonian onto the maximal subalgebra $\mathfrak{c} \subset \mathfrak{su}(d)$ that commutes with ρ . When degenerate states are considered, the eigenvalues' multiplicity defines the structure of the mask via $\delta_{\lambda_i \lambda_i}$.

We now notice that the unitary $U^{(0)} = \exp[iH_{\parallel}^{(0)}]$ can be composed with $O^{(0)}$ to obtain another unitary $O^{(1)} = O^{(0)}U^{(0)}$, formally equivalent to evolving with the time-dependent Hamiltonian $e^{-iH^{(0)}t}H_{\perp}^{(0)}e^{iH^{(0)}t}$ (as we detail in Sec. IV), which also drives $\rho \to \sigma$ since $[\rho, H_{\parallel}] = 0$. In general, the unitary $O^{(1)}$ is associated with a new geometric phase, and the Hamiltonian $H^{(1)} = i\ln O^{(1)}$ draws a different path, which might be shorter (or longer) than that generated by $H^{(0)}$. In a second iteration, we apply the mask to the new Hamiltonian to obtain $H_{\parallel}^{(1)} = \mathcal{M}_{\rho}[H^{(1)}]$ and define another gate $O^{(2)}$ analogously. By iterating this method, we obtain a sequence of Hamiltonians $\{H^{(i)}\}_{j=0}^{n}$ that drive $\rho \to \sigma$ via $\exp[-iH^{(j)}]$. The unitary at each step is related to the previous one by

$$O^{(j+1)} = O^{(j)} e^{-i\mathcal{M}_{\rho}[H^{(j)}]}, \text{ with } H^{(j)} = i \ln O^{(j)}.$$
(4)

Algorithm 1: Iterative method for efficient Hamiltonians.

Input: Initial state $\rho = \sum_{k=1}^{d} \lambda_k r_k\rangle \langle r_k $ and final state
$\sigma = \sum_{k=1}^{d} \lambda_k s_k\rangle \langle s_k $. Initial phase $\varphi^{(0)}$. Threshold
for convergence ε .

Output: Optimal Hamiltonian.

- 1: Initialize $O^{(0)} = \sum_{k} e^{i\varphi_{k}^{(0)}} |s_{k}\rangle \langle r_{k}|, H^{(0)} = i \ln O^{(0)},$ $H^{(0)}_{\scriptscriptstyle \parallel} = \mathcal{M}_{\rho}[H^{(0)}];$
- 2: while $||H_{\parallel}^{(j)}||_{\text{HS}} > \varepsilon ||H^{(j)}||_{\text{HS}}$ do 3: Set $O^{(j)} = O^{(j-1)}e^{iH_{\parallel}^{(j-1)}}, H^{(j)} = i\ln O^{(j)},$ $H^{(j)}_{\scriptscriptstyle \parallel} = \mathcal{M}_{\rho}[H^{(j)}];$
- 4 end while
- 5: **return** the final Hamiltonian $H^{(n)}$ of the sequence $\{H^{(j)}\}_{j=0}^{n}$.

A necessary and sufficient condition for the routine to reach a fixed point $H^{(n)}$, such that $H^{(n+1)} = H^{(n)}$, is given by $\mathcal{M}_{\rho}[H^{(n)}] = 0$, i.e., when the parallel part of the *n*th Hamiltonian vanishes under the action of the mask \mathcal{M}_{ρ} , which follows trivially from the fact that in this case $e^{-i\mathcal{M}_{\rho}[\dot{H}^{(n)}]} = \mathbb{1}$.

When implementing the method numerically, one has to fix a convergence threshold to stop the routine as soon as the parallel component $H_{\parallel}^{(n)}$ becomes small enough with respect to the full Hamiltonian $H^{(n)}$. We have chosen to quantify this threshold with a positive number $\varepsilon \ll 1$, such that the method stops when $\|H_{\parallel}^{(n)}\|_{\text{HS}} \leq \varepsilon \|H^{(n)}\|_{\text{HS}}$. Accordingly, the number of iterations n required for the method to converge implicitly depends on ε . Numerical evidence suggests that the sequence always converges toward a Hamiltonian $H^{(n)}$ that is fully perpendicular with respect to ρ along the whole evolution, in the sense that the parallel components eventually vanish within the precision defined by ε . This iterative method is summarized in its simplest form by Algorithm III, available in Ref. [52], and can be interpreted as an optimisation of energy cost associated with the different geometric phases φ , accomplished via the recursive suppression of ineffective components of the Hamiltonians.

There is also a geometric interpretation of our method analogous to that of Grover's famous quantum search algorithm [53]. The latter aims to find the unique input, encoded in a quantum state, of a function that produces a particular output value. Its action on an initialization state $|\psi^{(0)}\rangle$ can be interpreted as a sequence of rotations that converges with high probability to the desired state. Similarly, the iterative method introduced here can be seen as a sequence of rotations of some initialization unit vector $\hat{\boldsymbol{h}}^{(0)}$, associated with the Hamiltonian via $H^{(0)} = \mathbf{h}^{(0)} \cdot \mathbf{\Lambda}$, where $\mathbf{\Lambda} = (\Lambda_1, \dots, \Lambda_{d^2-1})$ forms a Lie algebra for SU(d). However, unlike Grover's algorithm, the vectors $\{\boldsymbol{h}^{(n)}\}$ do not span a two-dimensional real plane, but a high-dimensional subspace of \mathbb{R}^{d^2-1} .

IV. DERIVATION OF ALGORITHM

Let $O = \sum_{k} |s_k\rangle \langle r_k|$ and $i \ln O = H = H_{\parallel} + H_{\perp}$, where H_{\parallel} commutes with ρ . Consider the unitary U_t generated by the time-dependent Hamiltonian $O_t H_\perp O_t^{\dagger}$, where $O_t =$

 $\exp[-iHt]$, which satisfies the equation

$$\partial_t U_t = -iO_t H_\perp O_t^\dagger U_t. \tag{5}$$

To leading order in time, this will rotate the perpendicular part of the Hamiltonian to follow the system's evolution. Now consider the unitary $\widetilde{U}_t = O_t^{\dagger} U_t$, whose equation of motion is given by

$$\partial_t \widetilde{U}_t = -iH_{\perp} \widetilde{U}_t + (\partial_t O_t^{\dagger})U_t$$

= $-i(H_{\perp} - H)\widetilde{U}_t$
= $iH_{\parallel} \widetilde{U}_t.$ (6)

Equation (6) implies that $\widetilde{U}_t = \exp[iH_{\parallel}t]$, and thus that \widetilde{U}_t commutes with ρ . Accordingly, $U = U_1 = O_1 \widetilde{U}_1$ transforms $\rho \rightarrow \sigma$, just as well as O does. The procedure can be iterated until it reaches a fixed point, i.e., until $H = H_{\perp}$ and $H_{\parallel} = 0$.

The parallel part H_{\parallel} of the Hamiltonian is not uniquely defined in terms of the commutator with the initial density matrix. In order to fix a precise definition, one can consider the Hamiltonian's projection onto the maximal subalgebra $\mathfrak{c} \subset \mathfrak{su}(d)$ that commutes with ρ . Such subalgebra is unique and given by $\mathfrak{c} = U^{\dagger} \bigotimes_{i} \mathfrak{su}(d_{i})U$, where d_{i} is the degeneracy of the *j*th unique eigenvalue of ρ and U is the unitary that diagonalizes it. That mask defined in Eq. (3) achieves precisely this projection, in the sense that there is no component of H_{\perp} that commutes with ρ . Let us consider a fully nondegenerate initial state $\rho = \sum_{k} r_k |k\rangle \langle k|$. Let us imagine that $G \in \mathfrak{su}(d)$ is also an component of H_{\perp} , i.e., $G = \sum_{i \neq j} g_{ij} |i\rangle\langle j|$ in the eigenbasis of ρ , as prescribed by the mask in Eq. (3). For G to commute with ρ , we must have

$$[G, \rho] = \sum_{i \neq j \mid k} g_{ij} r_k(|i\rangle\langle j|k\rangle\langle k| - |k\rangle\langle k|i\rangle\langle j|)$$
(7)

$$=\sum_{i\neq j}^{n}g_{ij}(r_j-r_i)|i\rangle\langle j|=0, \qquad (8)$$

which vanishes for all the linearly independent components $|i\rangle\langle j|$ if and only if $g_{ij} \equiv 0$, since $r_i \neq r_i$ for hypothesis (nondegenerate spectrum). This means that G can only commute with ρ if it is trivial. Similar arguments can be made for the case of degenerate states, considering the form of the map $M_{ij} = \delta_{\lambda_i \lambda_j}.$

The iterative method introduced here involves removing the part of the Hamiltonian operator that commutes with the initial state ρ . A similar approach could be taken, where the part of the Hamiltonian that commutes with the final state σ is removed, and the unitary joining the two states is acted on from the left in each iteration, instead of from the right. It also possible to remove both components simultaneously, replacing Eq. (4) with $O^{(j+1)} = e^{i\hat{\mathcal{M}}_{\sigma}[H^{(j)}]}O^{(j)}e^{-i\mathcal{M}_{\rho}[H^{(j)}]}$. While we find this approach generally changes the convergence rate of the algorithm, it does not seem to perform significantly better than the algorithm outlined in the main text.

V. PERFORMANCE OF THE METHOD

Our iterative method can be directly applied to solve QBPs defined by the unconstrained time-optimal unitary evolution of any isospectral pair of density operators. A subclass of these problems with known solutions are those of

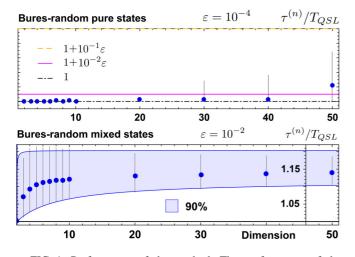


FIG. 1. Performance of the method. The performance of the method is here studied for Bures-random pure (top) and mixed (bottom) states, using the ratio $\tau^{(n)}/T_{\text{QSL}} \ge 1$ between the optimized time $\tau^{(n)}$ required to drive $\rho \rightarrow \sigma$ with the solution $H^{(n)}$, and the QSL T_{QSL} given in Ref. [10], evaluated for each pair of states ρ , σ , and optimized Hamiltonian $H^{(n)}$. The plotted points show the average performance for different Hilbert space dimensions d, with the *error* bars representing 99% (top) and 90% (bottom) confidence intervals. The shaded area (bottom) represents a fit of the 90% confidence interval for the average $\tau^{(n)}$. The convergence threshold is $\varepsilon = 10^{-4}$ for pure states (top) and $\varepsilon = 10^{-2}$ for mixed states (bottom). The sample size is 10^4 for each dimension. For pure states, the method returns solutions that converge on the QSL as the precision ε^{-1} is increased (at the expense of requiring more iterations).

unconstrained unitary evolution between pure states or between mixed states whose eigenvalues are all degenerate except one. To demonstrate the performance of our method, we have tested it on Bures-random pairs of pure states of dimension d = 2, ..., 100, successfully obtaining Hamiltonians that are time optimal and fully efficient with respect to the notion of efficiency introduced in Ref. [48]. We can be confident that the solutions we obtain are globally optimal by exploiting the fact that the QSL for unitary evolution of pure states is attainable [34]. Comparing the evolution time $\tau^{(n)}$ of the optimized Hamiltonian $H^{(n)}$ with the minimal evolution time T_{QSL} [54], we find $\tau^{(n)} \approx T_{QSL}$, within the precision imposed by ε , for all cases, as shown in Fig. 1.

A more challenging test was run on random pairs of mixed states of dimension d = 3, ..., 100 (for initial states with both degenerate and nondegenerate spectra), for which a general solution to the unconstrained unitary QBP is not known. Since the QSL for the unitary evolution of mixed states [10] is in general not tight, it is harder to benchmark the quality of the solutions provided by our method in this case. On the other hand, the natural synergy between this iterative method and the QSLs introduced in Ref. [10] can be used to gauge the performance of the former and the tightness of the latter. The *actual* minimal time of evolution τ_{\min} for a given choice of ρ and σ is bounded as

$$T_{\text{QSL}} \leqslant \tau_{\min} \leqslant \tau^{(n)},$$
 (9)

where T_{QSL} corresponds to the inviolable lower bound offered by QSLs of Ref. [10], and $\tau^{(n)}$ to an achievable upper bound,

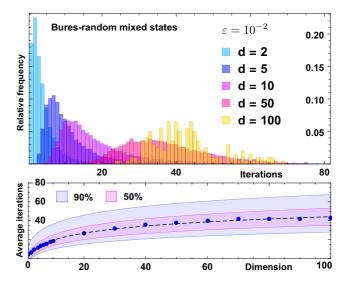


FIG. 2. Number of iterations required for convergence. The number *n* of iterations required for convergence is represented above for Bures-random mixed initial and final states ρ and σ , with Hilbert space dimension d = 2, ..., 100, as indicated in the plots. The histograms (top) show the relative frequency of the number of iterations required for convergence. The plotted points (bottom) show the average number of iterations required for convergence for the considered system dimensions, together with their 50% and 90% confidence regions. The black dashed line represents a logarithmic fit of the average number of iterations required to converge as a function of *d*. The sample size is 10^4 for each dimension $d \leq 50$ and 200 for d > 50, while the convergence threshold is chosen to be $\varepsilon = 10^{-2}$.

provided by the solutions obtained with our iterative method. The first inequality can be saturated for pure states and for states with d - 1 degenerate eigenvalues. For such case, the second inequality can be saturated up to the precision imposed by ε , which can be seen as an implicit trade-off between nand ε . By looking at the difference between $\tau^{(n)}$ and T_{QSL} for the given solution $H^{(n)}$, we can assess the quality of both results, which we find to coincide within the precision set by ε for some choices of ρ and σ , even when their degeneracy structure differs from that of pure states. The performance of the solutions $H^{(n)}$ obtained with this iterative method is stable under small perturbations of the initial state ρ , as described in Sec. VII. We now discuss the rate at which the algorithm converges.

VI. CONVERGENCE OF THE METHOD

The number *n* of iterations required for convergence generally grows with the dimension *d* of the system and the strictness of the precision set by ε . Recalling that the number of elements of the matrices associated with noncommuting density operators ρ and σ grows quadratically with *d*, one might expect conservatively that the average number of iterations \bar{n} would grow in the same way. However, \bar{n} grows logarithmically with *d*, as shown in Fig. 2, with a slower growth for the case of pure states and highly degenerate mixed states. Equivalently, for composite systems \bar{n} grows linearly with the number of constituent subsystems. The choice of initial phase vector $\boldsymbol{\varphi}^{(0)}$ also affects the number of

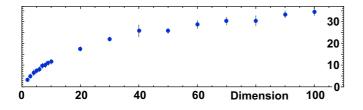


FIG. 3. Growth of the 20th percentile of *n*. The average 20th percentile of the number of iterations required for convergence is plotted along the Hilbert space dimension *d*, with the *error bars* representing the standard deviation. For each dimension we have sampled 10^2 pairs of initial and final states ρ and σ , while sampling 10^2 random initial phases for each pair. The slow growth of the lower end of the distribution of *n* suggests that Monte Carlo sampling methods are a feasible option to speed up the computation even for large *d*.

iterations required to converge, without noticeably affecting the performance of the end point $H^{(n)}$. In particular, running the algorithm backward from σ to ρ is equivalent to fixing a particular choice of initial phase vectors and produces the same optimized Hamiltonian.

Moreover, while \bar{n} grows with d and ε^{-1} , a good choice of initial geometric phase $\varphi^{(0)}$ can still lead to quick convergence, returning runs that can take less then 10 iterations even for d = 100 and $\varepsilon = 10^{-4}$. For these reasons, there is the potential to combine Algorithm III with Monte Carlo sampling methods [55-57], in order to deploy many parallel runs of the iterative method, for the same pair of initial and final states ρ and σ , with each run initialized with a different, randomly chosen geometric phase φ . Since the whole computation is stopped as soon as one of these runs converges to a solution, the number n of iterations is guaranteed to be smaller than the average one. We studied the dependence of n on the choice of initial geometric phases by running several instances of the method for the same pair of states ρ and σ , uniformly sampling each phase $\varphi_i^{(0)}$ in the interval [0, 2π]. As shown in Fig. 3, the slow growth of the 20th percentile of the number of iterations required for converges suggests that Monte Carlo sampling methods could combined with this method to speed up the computation.

VII. STABILITY UNDER PERTURBATIONS

The iterative method defined by Algorithm III is stable under small perturbations of the initial and final states ρ and σ . We considered convex and unitary perturbations, respectively given by

$$\rho \to \rho' = (1 - \delta)\rho + \delta \chi \text{ and}$$
 (10)

$$\rho \to \rho' = e^{iV\delta} \rho \ e^{-iV\delta}. \tag{11}$$

where χ is a random state of the same dimension as ρ and V is a random Hamiltonian of unit Hilbert-Schmidt norm. For each pair ρ , σ , the perturbed final state σ' is then obtained by applying any unitary O that maps $\{|r_k\rangle\}_{k=1}^d \rightarrow \{|s_k\rangle\}_{k=1}^d$ on the perturbed initial state ρ' . We look for the effect of perturbations calculating the relative deviations from the unperturbed solutions, given by $\|H^{(n)} - H'^{(m)}\|_{\text{HS}}/\|H^{(n)}\|_{\text{HS}}$, where $H'^{(m)}$

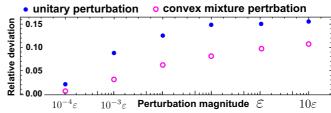


FIG. 4. Stability under perturbations. The relative deviation from the unperturbed solutions, given by $||H^{(n)} - H'^{(m)}||_{\text{HS}}/||H^{(n)}||_{\text{HS}}$, where $H'^{(m)}$ is the solution to the perturbed problem, is plotted against the perturbation strength δ , with $\varepsilon = 10^{-3}$. Numerical evidence suggests that the relative deviations grow slowly as a function of the perturbation strength and are negligible for $\delta \ll \varepsilon$.

is the solutions to the perturbed problem. Numerical evidence suggest that the relative deviations grow slowly as a function of the perturbation strength δ and are negligible for $\delta \ll \varepsilon$, as shown in Fig. 4.

VIII. HAMILTONIAN EFFICIENCY FOR DENSITY OPERATORS

In Sec. II, we considered a notion of Hamiltonian efficiency, introduced in Ref. [48] for the case of pure states, and we showed it cannot be saturated ($\eta = 1$) when mixed states are considered. Nevertheless, when it comes to estimate how much energy of the driving Hamiltonian is wasted along components of the Lie algebra that do not actively contribute to driving the system, it is possible to adapt such notion of efficiency to the case of density operators. A possible generalization is given by

$$\eta^{\star}(H,\rho) = \frac{\sqrt{\text{tr}[\rho^2 H^2] - \text{tr}[(\rho H)^2]}}{\sqrt{\text{tr}[\rho^2 H^2] - (\text{tr}[(\rho H)^2] - \text{tr}[\rho H]^2)}},$$
 (12)

which is a positive function smaller than 1, that is saturated for $H_{\parallel} = 0$. The numerator of η^* is proportional to the speed $\|\dot{\rho}\|_{\text{HS}}$ of the generator $\dot{\rho} = -i[H, \rho]$ [10,41,42] and reduces to ΔH_{ρ} when pure states are considered, for which $\eta^* =$ $1 \leftrightarrow \eta = 1$. Moreover, like ΔH_{ρ} , $\|\dot{\rho}\|_{\text{HS}}$ is also invariant under the addition of a parallel component H_{\parallel} with respect to ρ . Accordingly, the Hamiltonians generated by the iterative method are of unit efficiency for all the considered QBPs, which reflect the ability to converge to a fully perpendicular Hamiltonian.

When applying this notion of efficiency to the sequence of Hamiltonians $\{H^{(j)}\}_{j=0}^n$ obtained with this iterative method, it is possible to witness a transient nonmonotonic behavior for η^* , followed by a rapid monotonic ascent toward unit efficiency, as shown in Fig. 5. While we cannot rule out the existence of functionals that are strictly monotonic over the sequence $\{H^{(j)}\}_{j=0}^n$, this signature of nonmonotonicity can be seen in other, much simpler, figures of merit, such as the energy tr[$\rho H^{(j)}$] of the initial state ρ with respect to the Hamiltonians $H^{(j)}$ of the sequence.

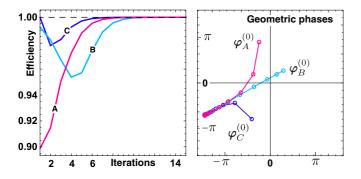


FIG. 5. Nonmonotonic solutions and their geometric phase trajectories. Algorithm III is here applied on the same choice of ρ and σ for three different choices of initial geometric phases $\varphi_A^{(0)}, \varphi_B^{(0)}$, and $\varphi_{C}^{(0)}$. The Hamiltonian efficiency $\eta^{\star}(H^{(j)}, \rho)$ of the Hamiltonians in the three different sequences is plotted along the iterations *i* for each run (left). Since each Hamiltonian (and unitary) has an associated geometric phase as prescribed by Eq. (2), we can plot the trajectory of the geometric phase for each run, such as as $\varphi_A^{(j)} = (\varphi_{A,1}^{(j)}, \varphi_{A,2}^{(j)}, \varphi_{A,3}^{(j)})$ for run A. By neglecting the first phase, which can be absorbed into a global phase, we plot the remaining components on a plane (right). All runs converge to the same solutions; however, while the efficiency of run A converges monotonically, the others converge nonmonotonically. The trajectories of their geometric phases do not give away the signature of nonmonotonicity, and straight paths can correspond to the slowest nonmonotonic descent toward the fixed point. Moreover, such trajectories can show richness and variety for different choices of ρ and σ and of initial geometric phases.

IX. DISCUSSION

In this paper, we have introduced an iterative algorithm to obtain efficient time-independent Hamiltonians that generate fast evolution between two *d*-dimensional isospectral states ρ

and σ . Such a method could be used to address fast driving, state preparation, and gate design problems. Recently, such problems have been experimentally tackled in Refs. [58–60] by means of shortcuts to adiabadicity and transitionless quantum driving [61–64]. Incidentally, these methods are closely linked to ours, with a key common element being the optimisation of geometric phases. Other possible applications of the iterative method introduced here could be found in its combination with quantum optimal control methods. Algorithm III can be used to identify a unitary O, associated with an efficient Hamiltonian H, to be implemented by means of time-optimal gate design methods, such as those experimentally realized in Ref. [2]. More generally, our method has broad application in further elucidating the geometric structure of quantum control, since it has been shown that quantum brachistochrone problems can be recast as those of finding geodesics in the space of unitary operators [12].

A challenging outlook is to extend and adapt this method to more general control problems. However, it is not clear that, for open quantum evolutions, simply replacing the Hamiltonian with a Liouvillian as the generator would preserve the complete positivity of the dynamics. Nevertheless, it is possible to adopt the current method for open dynamics by applying it to purification of ρ and σ or dilation of the dynamics. Another challenge that remains is to build in constraints on the form of the generators, naturally imposed by physical restrictions on the order and range of interactions. These can dramatically change the bounds on minimal evolution time, as recently shown in Ref. [32].

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