# Scrambling Dynamics and Out-of-Time-Ordered Correlators in Quantum Many-Body Systems

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This tutorial article introduces the physics of quantum information scrambling in quantum many-body systems. The goals are to understand how to precisely quantify the spreading of quantum information and how causality emerges in complex quantum systems. We introduce a general framework to study the dynamics of quantum information, including detection and decoding. We show that the dynamics of quantum information is closely related to operator dynamics in the Heisenberg picture, and, under certain circumstances, can be precisely quantified by the so-called out-of-time-ordered correlator (OTOC). The general behavior of the OTOC is discussed based on several toy models, including the Sachdev-Ye-Kitaev model, random circuit models, and Brownian models, in which the OTOC is analytically tractable. We introduce numerical methods, including exact diagonalization and tensor network methods, to calculate the OTOC for generic quantum many-body systems. We also survey current experimental schemes for measuring the OTOC in various quantum simulators.

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#### **I. INTRODUCTION**

In recent years, there have been remarkable developments in laboratory platforms for studying quantum physics. These systems range from ultracold atoms, trapped ions, and superconducting qubits to universal quantum computers, providing exciting opportunities to study quantum many-body physics that was previously out of reach. One research frontier concerns the long-time coherent quantum many-body dynamics in closed systems, which has drawn extensive research interests from multiple communities, such as condensed matter physics, atomic, molecular, and optical physics, quantum information science, and high-energy physics. Synergistic experimental and theoretical research has revealed a series of discoveries in the arena of quantum dynamics. Moreover, these platforms expand the scope of traditional condensed matter physics and demand new tools and frameworks to study quantum many-body systems that are far from equilibrium (see Ref. [1] for an overview on quantum simulators and the references therein).

In the simplest kind of quench experiment, one prepares an initial state, designs a Hamiltonian or unitary circuit to evolve the state, and measures the final state. While the freedom in the initial state and engineered dynamics largely depends on the specific experimental platform, this general class of experiments certainly raises questions regarding the general behavior of quantum many-body dynamics when the initial state is far from equilibrium. Consider a simple product initial state of qubits, with each qubit in  $|0\rangle$  or  $|1\rangle$ . Now, let the state evolve under a generic unitary operator. The general expectation is that the state will not remain a product state, but will instead become a complicated superposition of product states.

One way to track the complexity is to monitor the buildup of entanglement in the state. Entanglement can be quantified using the tool of entanglement entropy. Given a subregion A of the system, one can obtain the density matrix by tracing out the complement  $\overline{A}$ :

$$\rho_A = \operatorname{tr}_{\bar{A}} |\psi\rangle \langle\psi| \,. \tag{1}$$

The entanglement entropy of A is defined as

$$S(A) = -\mathrm{tr}\rho_A \log_2 \rho_A. \tag{2}$$

It is also straightforward to show that  $S(A) = S(\overline{A})$  when the total state is pure.

In a quench experiment starting from a product state, S(A) will begin at zero and then grow over time. This growth indicates that the density matrix is becoming more mixed, which corresponds to an increasingly featureless state of subsystem A. Based on statistical considerations, we expect the density matrix to approach a maximally mixed state if the evolving unitary is generic. In other words, at late times, the subsystem thermalizes by entangling with its environment. If we consider a system that conserves energy instead of a generic evolution, the latetime density matrix is expected to approach a thermal state with a temperature determined by the initial state's average energy. Once  $\rho_A$  thermalizes, it depends only on macroscopic quantities such as energy or charge, while microscopic information about the initial state is apparently lost [2-6].

Thermalization typically occurs within a short time scale, after which the complicated many-body wave function still evolves unitarily in a closed system. This evolution raises a pivotal question: beyond thermalization, is there a universal out-of-equilibrium process and a new timescale that come into play? One significant approach to investigate this is through scrambling dynamics, the primary focus of this tutorial. Owing to thermalization, two orthogonal initial states with the same energy or charge density would thermalize to the same local density matrix, and so one would be unable to distinguish them locally. However, since the dynamics is unitary, the two states remain orthogonal and are always distinguishable if global information about the two states is accessible. Thermalization and unitary dynamics suggest that the local features of the initial state become increasingly nonlocal under the dynamics and cannot be recovered by local probes. This phenomenon, where local information becomes distributed across a system, is what defines quantum information scrambling.

To be more concrete, let us consider a system of interacting qubits. One qubit is prepared in a single-qubit state  $|a\rangle$  to encode a secret message. At the outset, it is possible to retrieve the original state of this qubit through single-qubit state tomography. However, as time progresses, due to interactions among the qubits, the qubit initially in state  $|a\rangle$  no longer remains in that pure state but becomes entangled with others. Because of thermalization, any local state, which approaches the Gibbs state, no longer contains information about  $|a\rangle$ . Consequently, state  $|a\rangle$  becomes distributed among multiple qubits, making its recovery increasingly challenging, even in principle. Quantum information scrambling is quantified by the minimal subsystem required to recover the encoded state  $|a\rangle$  as time progresses. Initially studied in the context of black hole dynamics [7–10], quantum information scrambling has also been extended to general quantum many-body systems [11-13] and becomes a finer tool than thermalization to characterize nonequilibrium dynamics.

It is useful to contrast scrambling with thermalization. Thermalization and scrambling both result from interacting unitary dynamics. While related to thermalization, scrambling is a distinct process. Thermalization describes how a local region of a quantum system loses its initial information under unitary dynamics. In contrast, quantum scrambling concerns how the "lost" information flows to nonlocal degrees of freedom. These two processes are characterized by different length and time scales. The thermalization time of a local region is typically independent of the system size and is determined by the coupling energy scale. On the other hand, the scrambling time, which is roughly the time scale when the initial local information is fully shared among the system, typically depends on the system size.

In this tutorial, we delve into the topic of scrambling dynamics in quantum many-body systems. We argue that, just as a piece of metal can be characterized by its transport properties associated with electrons, a generic quantum many-body system can be characterized by its transport properties related to quantum information. Our tutorial aims to provide a deeper discussion by building on prior perspectives from related articles [14,15] that offer basic intuition.

The discussion is centered around a quantum information perspective with the overall aim of clarifying in a concrete and widely accessible setting the precise way in which out-of-time-order correlators measure information dynamics. We focus mostly on qubit models, including random circuit models, which are widely studied in the quantum information and condensed matter communities, and, to a lesser extent, in the high-energy physics and quantum gravity communities. The upshot of this approach is that we will be able to state very precisely the relationship between information dynamics and out-of-time-order correlators. The cost is that we must omit or be much more schematic about many topics, including semiclassical physics, field theory approaches, connections to black hole physics [9], and much else. One may reasonably conjecture that the basic connection between information dynamics and out-of-time-order correlators extends to these settings, but a considerably greater background is required to define these models and properly formulate the notion of information dynamics in them (e.g., see Refs. [16,17] for a discussion of AdS/CFT models in the same spirit as this tutorial). For the same reason, the main body of the tutorial does not cover the connection between scrambling dynamics and quantum chaos, despite a substantial body of literature on this subject. Interested readers are directed to a brief discussion in the Epilogue (Sec. VIII) and the reference therein.

The rest of the tutorial is structured as follows. In Sec. II, we examine the fundamental setup of scrambling dynamics using an Alice-Bob communication protocol. In Sec. III A, we explore how to quantify scrambling dynamics through entanglement entropy measures, offering basic insight through random unitary dynamics. In Sec. IV, we examine the Hayden-Preskill protocol, a specific instance of the general setup, and demonstrate that the scrambling dynamics can be quantified by the outof-time-ordered correlator (OTOC). In Sec. V, we link scrambling and the OTOC to operator dynamics and provide an overview of the OTOC in systems with few-body interactions. In Sec. VC, we delve into the behavior of the OTOC in systems with short-range interactions through several toy models. In Sec. VI, we survey numerical methods for calculating the OTOC in general systems. Finally, in Sec. VII, we examine various experimental approaches for measuring the OTOC.

# **II. BASIC SETUP OF SCRAMBLING DYNAMICS**

We begin by specifying our prototype quantum manybody system. For the sake of concreteness, we consider a system of N qubits. Each qubit has a basis that is spanned by  $|0\rangle$  and  $|1\rangle$ . On each two-level system, a complete basis of operators can be defined, which consists of the identity and the Pauli operators. These operators are represented by matrices as follows:

$$I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad \sigma^{x} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \sigma^{y} = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix},$$
$$\sigma^{z} = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}.$$
(3)

The total Hilbert space is the tensor product of local Hilbert spaces and has a dimension of  $2^N$ . In some parts of this tutorial, we generalize the situation to include qudits with

a local Hilbert space dimension of q, or Majorana fermion systems. However, for now, let us continue to work with qubits.

The dynamics of a given initial state  $|\psi\rangle$  in the system is described by a unitary time evolution operator given by

$$|\psi(t)\rangle = U(t) |\psi\rangle, \qquad (4)$$

where U(t) is currently an arbitrary family of unitary matrices that acts on the total Hilbert space. However, we assign more structure to U(t) later on. The simplest quench experiment involves selecting an initial state  $|\psi\rangle$ , choosing a dynamics U(t), and selecting a set of observables to measure in the final state  $U(t) |\psi\rangle$ .

Throughout this tutorial, we often use tensor network diagrams to provide a visual representation of equations for clarity. We now introduce the graphic notation below. In general, a node with legs represents a tensor that is a multidimensional array, and each leg represents an index of the tensor. For instance,  $\bigcirc$  with one leg represents a single-qubit state vector,  $\bigcirc$  with two legs represents a single-qubit operator or a matrix in general, and  $\square$  has four legs and thus represents a tensor such as a two-qubit operator. In particle, we use a line to represent the identity operator  $\delta_{ab}$  and a line with a dot for the normalized identity operator  $\delta_{ab}/\sqrt{d}$ , where d is the dimension of the index. It can also be interpreted as the maximally entangled Einstein-Podolsky-Rosen (EPR) state, i.e.,

In a tensor network diagram, the nodes are connected together by joining their legs and summing over the shared indices. For example,

represents a matrix-vector multiplication or an operator acting on a state, and the result with one open leg is a vector.

# A. Unitary dynamics as a classical communication protocol: the significance of commutators

We first consider unitary dynamics as an intuitive classical communication protocol to illustrate some of the essential aspects of scrambling dynamics, and in the next section, we reformulate it in full quantum terms. Consider the scenario illustrated in Fig. 1, where Alice owns one

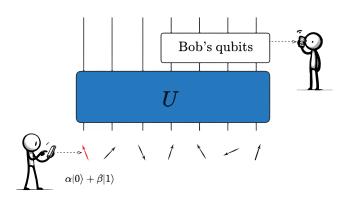


FIG. 1. Alice and Bob try to communicate through a strongly interacting system of N qubits. The time evolution of the system is described by a unitary operator U. Alice has full control of the first qubit  $q_A$  and Bob has access to a set of qubits in the system, but not all of them.

of the *N* qubits denoted by  $q_A$ , which she has full control over. Bob owns a set of qubits denoted by *B*. Starting with an initial state  $|\psi\rangle$ , Alice wishes to send a classical bit  $a \in 0, 1$  to Bob. Depending on the value of *a*, Alice either flips her qubit by applying the  $\sigma_{q_A}^x$  operator to the state or does nothing. The system then evolves for a time *t*. Finally, Bob makes a measurement  $O_B$  of his qubits to attempt to learn whether Alice flipped the spin or not.

The expectation value of Bob's measurement given that Alice does not flip her qubit is

$$\langle O_B \rangle_0 = \langle \psi | U^{\dagger} O_B U | \psi \rangle = \langle \psi(t) | O_B | \psi(t) \rangle.$$
 (7)

Alice can also flip her qubit, and then the expectation value of Bob's measurement is

$$\langle O_B \rangle_1 = \langle \psi | \sigma_{q_A}^x U^{\dagger} O_B U \sigma_{q_A}^x | \psi \rangle$$
  
=  $\langle \psi(t) | \sigma_{q_A}^x(-t) O_B \sigma_{q_A}^x(-t) | \psi(t) \rangle,$  (8)

where  $\sigma^{x}(-t) = U\sigma^{x}U^{\dagger}$  is a Heisenberg operator. The difference between the expectation values  $\langle O_B \rangle_0$  and  $\langle O_B \rangle_1$  is

$$\langle O_B \rangle_0 - \langle O_B \rangle_1 = \langle \psi(t) | \sigma^x(-t) [\sigma^x(-t), O_B] | \psi(t) \rangle, \quad (9)$$

where we have used the operator identity  $\sigma^{x}(-t)\sigma^{x}(-t) = I$ .

Let us pause to understand the physics. Whenever the difference is small, Alice and Bob need to run the experiment many times to see the difference and communicate as the measurement outcome of each run is probabilistic. In this case, very little information is transmitted from Alice to Bob per run of the experiment since what Bob measures is nearly independent of what Alice did.

Using the Cauchy-Schwarz inequality (we provide a list of useful inequalities in Appendix A 3), we can bound the

difference by

$$|\langle O_B \rangle_1 - \langle O_B \rangle_0|^2$$
  

$$\leq \langle \psi(t) | [\sigma_{q_A}^x(-t), O_B]^{\dagger} [\sigma_{q_A}^x(-t), O_B] | \psi(t) \rangle$$
  

$$\leq \| [\sigma_{q_A}^x(-t), O_B] \|_{\infty}^2, \qquad (10)$$

where  $||O||_{\infty}$  is the operator norm defined as the square root of the largest eigenvalue of the positive Hermitian operator  $O^{\dagger}O$ . The first inequality is from the Cauchy-Schwarz inequality, and the second is from the operator norm's definition. (Definitions of various matrix norms of operators are given in Appendix A 1.) Therefore, the difference in Bob's measurement between Alice flipping her qubit or not is bounded by the operator norm of the commutator  $[\sigma^x(-t), O_B]$ . This statement is independent of the initial state  $|\psi\rangle$  that Alice and Bob choose as the medium to attempt to transmit the information.

The bound has a very intuitive interpretation. At t = 0,  $\sigma^x$  only has support on Alice's qubit and does not overlap with Bob's qubits. Therefore, the commutator is zero initially, and no matter what Bob does to his qubits, he cannot tell whether Alice flips her qubit. He can do no better than random guessing when trying to determine Alice's bit. As *t* increases,  $\sigma^x(-t)$  starts to grow as a Heisenberg operator, acting on more qubits. Whenever the supports of  $\sigma^x(-t)$  and  $O_B$  start to overlap, their commutator becomes nonzero, and Bob has a chance to tell whether Alice flipped her qubit. The operator norm of the commutator  $|[\sigma_{q_A}^x(-t), O_B]|_{\infty}$  quantifies how quickly  $\sigma q_A^x(-t)$ spreads in the system and starts to overlap with  $O_B$ . If this operator norm is small then the overlap is small, and Bob's measurement cannot distinguish between the two cases of Alice flipping her qubit or not.

#### **B.** Bound on commutators

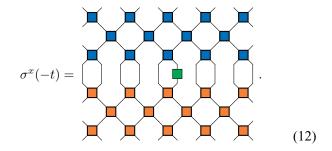
A natural first question is whether fundamental bounds exist on the norm of the commutators given t and r, which is the separation between  $q_A$  and B. There are many possible behaviors for the commutator of local operators in a quantum many-body system. For example, one might expect very different behavior between integrable, noninteracting, and strongly interacting models, and between localized and delocalized models.

One is probably familiar with at least one such constraint, namely, the limitation on communication imposed by the speed of light. In the modern language of quantum field theory, this is called microcausality. It states that, given any two physical local operators W(x) and V(y)located at space-time points x and y, their commutator must vanish if x and y are "spacelike separated,"

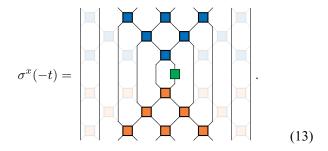
$$x, y$$
 spacelike separated  $\rightarrow [W(x), V(y)] = 0.$  (11)

In other words, if y is outside of the "light cone" of space-time point x then the corresponding operators must exactly commute. Crucially, this is an operator statement and hence a state-independent bound on information propagation. It is a fundamental property of any unitary Lorentz-invariant local quantum field theory.

There is somewhat analogous property for many lattice models that do not have relativistic causality built-in microscopically. As discussed earlier, the commutator is zero when  $O_B$  is outside the support of the Heisenberg operator  $\sigma^{x}(-t)$ . Therefore, the support  $\sigma^{x}(-t)$  that grows as a function of time serves as an emergent "light cone" for the lattice models, which describes how fast information can propagate in these systems. To provide an explicit example of the emergent light cone, let us consider the unitary time evolution operator with a tensor network structure built from local two-qubit unitary gates, as sketched in Fig. 2. This unitary operator describes the time evolution of a spin chain with nearest-neighboring interaction, without relativistic causality built in but only locality. Given the brickwork structure of U, the tensor network representation of a Heisenberg operator, say  $\sigma^{x}(-t) = U\sigma^{x}U^{\dagger}$ , is



Each blue or orange block represents a local unitary matrix, and the green block represents the  $\sigma^x$  at t = 0. The blue block and orange block are unitaries conjugated to each other. As a result, if a blue block and an orange block are directly connected, they are replaced by straight lines, representing identity operators. The tensor network after this transformation is



The remaining blocks (nonshaded region) form a linear light cone, visualizing the growth of the Heisenberg operator over time. The effective speed of light is  $a/\Delta t$ , where a is the lattice constant and  $\Delta t$  is the time scale associated with one layer of the unitary circuit. This is

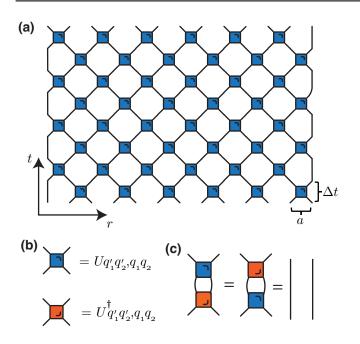


FIG. 2. (a) The tensor network representation of a unitary circuit with spatial locality built in, the so-called brickwork circuit. (b) Each blue block with four legs (two in and two out) represents a two-qubit unitary gate  $U_{q'_1q'_2,q_1q_2}$ . (c) Tensor network diagrams for the identities  $u^{\dagger}u = I$  and  $uu^{\dagger} = I$ .

a simple but remarkable result, showing that an effective linear light cone can emerge from the locality in systems without microscopic relativistic causality. The commutator  $[\sigma^x(-t), O_B]$  is strictly zero when  $O_B$  is outside the emergent light cone.

There is one caveat to this approach to obtain the speed limit for a static Hamiltonian. For instance, consider a Hamiltonian describing nearest-neighboring interactions between qubits:

$$H = \sum_{r} H_{r,r+1}.$$
 (14)

One can Trotterize the time evolution operator  $\exp(-iHt)$  to the local tensor structure shown in Fig. 2. Each local unitary block takes the form  $u = \exp(-iH_{r,r+1}\Delta t)$ . The tensor structure in Fig. 2 approaches  $\exp(-iHt)$  in the limit  $\Delta t \rightarrow 0$ . Then the velocity  $a/\Delta t$  from the tensor structure goes to infinity, which is not a meaningful bound. Nevertheless, for discrete models with a local Hamiltonian and a finite local Hilbert space dimension, one can establish a much tighter linear light cone with finite speed of the commutator using a more sophisticated approach originally due to Lieb and Robinson, which we discuss in Appendix E. This bound, now usually referred to as the Lieb-Robinson bound [18], states that, for two local operators W(r) and V(r') at (spatial) positions r and r', respectively, we have the following

bound on the operator norm of the commutator:

$$\frac{\|[W(r,t), V(r',0)]\|_{\infty}}{\|W\|_{\infty} \cdot \|V\|_{\infty}} \le ae^{\lambda(t-|r-r'|/v_{\mathrm{LR}})}$$
(15)

with a,  $\lambda$ , and  $v_{LR}$  depending on the microscopic parameters of the Hamiltonian. This inequality, combined with Eq. (10), bounds the difference of the signal that Bob measures at time t to determine whether Alice flips her qubit at time t = 0. Observe that if the distance  $|r' - r| \ge v_{LR}t$ , the bound, although not zero, is exponentially small, indicating that it is almost impossible for Bob to tell whether what Alice did. This establishes an approximate light cone with a finite speed  $v_{LR}$ , called the Lieb-Robinson velocity, for local systems without relativistic causality.

## C. Two central goals: detection and recovery

The Lieb-Robinson bound is independent of state  $|\psi\rangle$ and is universally applicable. Its importance lies in proving that information cannot travel super ballistically in quantum many-body systems with short-ranged interaction. However, in many cases, the bound can be quite loose, just like the physical speed of light is a loose bound on how fast an object moves in our universe. Moreover, it does not provide a way to calculate  $v_{LR}$ , but only proves that it is finite. Therefore we need an operational approach to calculate how fast information spreads for a given system—the information light cone (we provide a precise definition of the information light cone in the next section). For now, we can interpret the information light cone as the minimal set of qubits, outside which Bob either cannot distinguish Alice's actions at t = 0 or needs exponentially many measurements to do so. This set only contains Alice's qubit  $q_A$  at t = 0, but grows over time. One of the central goals in studies of scrambling dynamics is to calculate the linear size of this region as a function of t. The Lieb-Robinson bound provides an upper bound  $R(t) \leq v_{LR}t$  for local Hamiltonians.

A key second goal is to determine Bob's optimal measurement to capture the signal sent by Alice, or to find the optimal  $O_B$  to maximize the difference between  $\langle O_B \rangle_0$  and  $\langle O_B \rangle_1$ . In noninteracting or weakly interacting systems, the excitation created by Alice can remain coherent for a long time and propagates with a group velocity determined by the underlying medium. Imagine a wave packet of an electron or a magnon moving through a metal or a magnet. In this case, Bob can easily tell whether Alice made the excitation by performing a local measurement, since the signal (the wave packet) remains local for a long time. On the other hand, in strongly interacting quantum systems, the physics of excitations is typically very different. In fact, such a system often cannot sustain any coherent excitation for very long, unless that excitation has a special reason for being protected, such as a sound mode or a Goldstone

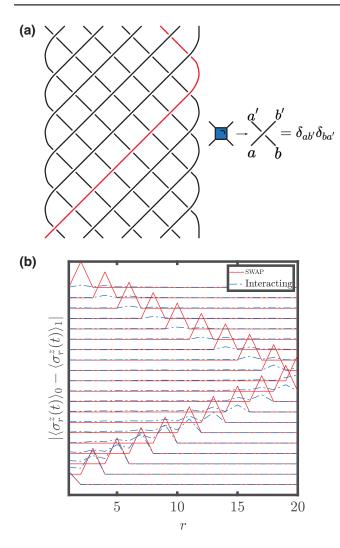


FIG. 3. (a) Quantum unitary circuit made from a SWAP gate. The circuit is constructed by replacing the generic two-qubit gate in Fig. 2 with a noninteracting SWAP gate  $\times$ . In this case, Alice's qubit state propagates through the system coherently and remains on a single qubit (the red line). The signal travels ballistically because of the light-cone structure of the circuit model shown in Eq. (13). (b) The maximal signal Bob can obtain from single-qubit measurements. The *x* axis labels the qubit. Each line represents the signal at a different time with an offset. In the case of circuits from SWAP gates, the signal (red curves) remains maximal and propagates through the system. On the other hand, when we perturb the SWAP gate by multiplying it by another unitary gate close to identity, the signal (blue curves) decays quickly.

mode associated with a broken symmetry. Figure 3 illustrates this difference between a noninteracting system and a strongly interacting system. We use the circuit in Fig. 3 to represent the noninteracting case, which is built from the two-qubit SWAP gate. Since the SWAP gate simply swaps the two qubits it acts on, Alice's state on the first qubit propagates coherently through the system and always remains on a single qubit that Bob can measure. On the other hand, once the local gate of the circuit is perturbed away from the SWAP gate, the signal generally decays as time increases.

The lack of coherent excitations is a consequence of quantum thermalization. For a strongly interacting system, a time-evolved state becomes as complicated as allowed, consistent with macroscopic constraints, such as a fixed total energy. As a result, the state looks thermal locally when  $t > \tau$ , where  $\tau$  is a relaxation time. This is to say, for any local operator O, we have

$$\langle \psi(t) | O | \psi(t) \rangle \underset{t \gg \tau}{\approx} \frac{\operatorname{tr}(e^{-\beta H}O)}{\operatorname{tr}(e^{-\beta H})},$$
 (16)

where  $\beta$  depends only on the average energy  $\langle \psi | H | \psi \rangle$  of the initial state  $|\psi\rangle$ . Technically, a finite system will eventually revive to its initial state after a very long (double-exponential) time scale [19,20], which we do not consider here.

Let us return to the communication protocol between Alice and Bob. Since  $|\psi_0\rangle$  and  $|\psi_1\rangle = \sigma_{q_A}^x |\psi_0\rangle$  only differ by a single-qubit flip, they have the same energy density. Therefore, we expect the two states to thermalize and look the same locally after a time  $\tau$ . Now imagine that Bob is trying to tell what Alice did by performing local measurements in a region some distance away from Alice's qubit  $q_A$ . At early time, the information light cone has not arrived at Bob's region yet, and, by definition, Bob can do nothing to tell the difference between the two states. Later, when the information light cone reaches Bob's region, the time is well past the relaxation time  $\tau$ , so Bob still cannot tell the difference between the two states since they now look the same locally. Therefore, unlike in the noninteracting case, Bob cannot tell what Alice did using only local measurements when the system thermalizes.

It is not plausible that information has simply stopped spreading in the strongly interacting system, but it becomes inaccessible to local measurements. Said differently, it may be very difficult to transmit information in the Alice-Bob communication protocol coherently. However, information is still spreading and Bob needs an approach to recover it. We know at least one approach; Bob can perform the measurement using the Heisenberg operator  $\sigma_0^z(-t)$ , which is highly nonlocal. Then it is equivalent to measuring  $\sigma^z$ at t = 0, and Bob can easily tell what Alice did. However, this approach might not be optimal since  $\sigma_0^z(-t)$  can be very complex and its support contains Alice's qubit as well.

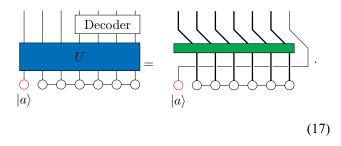
The discussion so far was framed in terms of classical information—whether Bob can tell if Alice flips her qubit or not. A stronger version is about transmitting quantum data. One can ask the following question. Given that Alice initially prepares an arbitrary quantum state on her qubit  $q_A$ , is it possible for Bob to recover that quantum state on one of his qubits, denoted  $q_B$ , following some decoding procedure, after the system is evolved by some time t?

The takeaway of this section is that information spreading is intimately related to commutators of local operators and its speed is upper bounded by the possibly very loose Lieb-Robinson bound. This section also raises two central questions regarding quantum information dynamics in strongly interacting quantum many-body systems.

- (a) How is the information propagation detected?
- (b) How is the information recovered?

#### **III. QUANTUM INFORMATION FORMULATION**

In this section, we formulate the general problem of information in fully quantum terms. Recall that Alice initializes one qubit,  $q_A$ , into an arbitrary quantum state. The medium that Alice and Bob share then evolves by the unitary U. Finally, Bob's goal is to apply some decoding operation to isolate the quantum information  $|a\rangle$  that Alice originally encoded in qubit  $q_A$ ,



In general, Bob's decoding can be a very complex quantum operation. We discuss examples of Bob's decoding in Sec. IV below, but let us first understand how to determine which qubits Bob needs to control to decode Alice's information in principle. In other words, let us understand how to track where the quantum information is before the decoding. Remarkably, this task can be accomplished just by following a special kind of entanglement with an auxiliary reference system, as we next explain [see also Sec. IV(b) below].

#### A. Entanglement spreading

Consider a system containing N qubits whose dynamics is described by a unitary matrix U(t). We pick two orthogonal initial states,  $|\psi_0\rangle$  and  $|\psi_1\rangle$ , that differ by a spin flip at Alice's qubit,

$$|\psi_0\rangle = |0\rangle_{q_4} \otimes |\psi\rangle, \qquad |\psi_1\rangle = |1\rangle_{q_4} \otimes |\psi\rangle, \qquad (18)$$

where  $|0\rangle$  and  $|1\rangle$  are two orthogonal states in some basis on Alice's qubit  $q_A$ , and  $|\psi\rangle$  is the state of the remaining system. These two states correspond to the two possible initial states in the Alice-Bob communication protocol in Sec. II.

Next, introduce a new auxiliary qubit called reference R. The unitary dynamics U does not act on the reference

qubit, which one may think of as sitting in an isolated box. Before isolating the reference, however, the reference qubit is entangled with the system through Alice's qubit  $q_A$ . The initial composite state of the system and the reference is

$$\begin{split} |\Psi\rangle &= \frac{1}{\sqrt{2}} (|0\rangle_R |\psi_0\rangle_{\text{SYS}} + |1\rangle_R |\psi_1\rangle_{\text{SYS}}) \\ &= \frac{1}{\sqrt{2}} (|0\rangle_R |0\rangle_{q_A} + |1\rangle_R |1\rangle_{q_A}) |\psi\rangle \,. \end{split}$$
(19)

The reference R and Alice's qubit  $q_A$  form a Bell pair and are maximally entangled. The time evolution of the state is given by

$$\begin{split} |\Psi(t)\rangle &= \frac{1}{\sqrt{2}} (|0\rangle_R U |\psi_0\rangle_{\text{SYS}} + |1\rangle_R U |\psi_1\rangle_{\text{SYS}}) \\ &= \frac{1}{\sqrt{2}} (|0\rangle_R |\psi_0(t)\rangle_{\text{SYS}} + |1\rangle_R |\psi_1(t)\rangle_{\text{SYS}}), \quad (20) \end{split}$$

or, graphically,

0370

with the black dot representing the EPR state  $(|00\rangle + |11\rangle)/\sqrt{2}$ .

Given the initial state, we may probe the information dynamics by tracking the entanglement between the system and the reference as a function of time. Initially, the reference is only entangled with the first qubit  $q_A$ . This entanglement can be diagnosed using mutual information between *R* and  $q_A$ . First, define the von Neumann entropy of a set of qubits *A*,

$$S(A) = -\operatorname{tr}(\rho_A \log_2 \rho_A), \qquad (22)$$

where  $\rho_A$  is the reduced density matrix of A. Then the mutual information of A with B is

$$I(A:B) = S(A) + S(B) - S(AB).$$
 (23)

Using subadditivity and the triangle inequality (see Appendix A 3), one can show that

$$0 \le I(A:B) \le 2\min(S(A), S(B)).$$
 (24)

For *R* and  $q_A$ , it is straightforward to show that at t = 0,  $S(R) = S(q_A) = 1$  and  $S(R \cup q_A) = 0$  from the fact that they form a Bell pair initially. Therefore, initially,  $I(R : q_A) = 2$ , which is the largest possible value, indicating maximal entanglement between  $q_A$  and *R*.

Correspondingly, at t = 0, the mutual information between R and any other set of qubits is zero.

Starting from the initially localized entanglement, one should expect the entanglement with reference R to expand out across the system in some fashion. One possibility is that the entanglement is carried in some coherent wave packet throughout the system, remaining localized in space at any given time. This can occur under the right conditions in noninteracting systems, for instance, in the SWAP circuit shown in Fig. 3. However, with strong interactions, the entanglement seems likely to spread and become more complex [21]. In other words, while at time zero the reference is entangled with one qubit; as time progresses, the reference will instead become entangled with a complex collection of many qubits. This process can be quantified by the mutual information between the reference qubit and certain regime B in the state.

For example, one can choose *B* to include the first *l* qubits  $q_1 \cdots q_l$  in the system. Then, the mutual information between *R* and *B* is

$$I(R:q_1 \cdots q_l) = S(R) + S(q_1 \cdots q_l) - S(R + q_1 \cdots q_l).$$
(25)

At t = 0, the reference is entangled with  $q_A$  contained in the first *l* qubits. Therefore,  $I(R : q_1 \cdots q_l) = 2$  for all *l*. As time increases, fixing l,  $I(R : q_1 \cdots q_l)$  decreases once the entanglement leaks out the first l qubits. This behavior is illustrated in Fig. 4, where  $I(R:q_1\cdots q_l)$  for  $1 \leq 1$  $l \leq N$  are shown for the mixed-field Ising chain, which is a prototype nonintegrable quantum many-body model for studying thermalization [22]. The Hamiltonian is introduced in Eq. (109) in Sec. VI below on the numerical methods. The result, adapted from Ref. [17], describes the dynamics of 22 spins initiated in a high-energy pure state. Note the parallel curves for l < N/2. This is the first concrete example of the ballistic propagation of information in a strongly interacting system in this tutorial. Also, note that the late-time value of the mutual information increases with l and approximately stays at the maximal value of 2 for l > N/2. This is a signature of a strongly interacting system, which we explain in Sec. III C below using random unitary matrices.

#### **B.** Communicating quantum information

How is entanglement spreading related to the Alice-Bob communication protocol? The description above might have already hinted at some similarities. We now make the connection more precise. We consider a region B of the system and denote its complement as  $\overline{B}$ . The mutual information between the reference R and B is

$$I(R:B) = S(R) + S(B) - S(RB),$$
 (26)

which is between 0 and 2. Since state  $|\Psi\rangle$  defined on R, B, and  $\overline{B}$  is pure,  $S(RB) = S(\overline{B})$ . (See Appendix A 2).

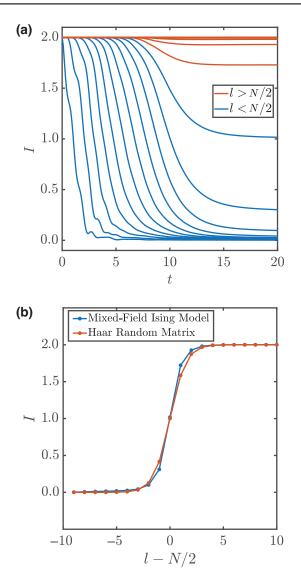


FIG. 4. (a) Entanglement spreading characterized by the mutual information  $I(R: q_1 \cdots q_l)$  between the reference qubit and first l qubits in the system of N qubits. Each curve represents the mutual information for *l* ranging from 1 to *N*. Here N = 22. For l < N/2, each curve remains at the maximal value 2 until the information leaks out and then decays to zero. The result is adapted from Fig. 2 of Ref. [17]. The parallel curves are a signature of ballistic propagating quantum information. As l exceeds half of the system, the mutual information does not decay, indicating that one can recover the initial entanglement between Rand the first qubit if one has access to more than half of the system. (b) The late-time saturation value of the mutual information in (a). It undergoes a sharp transition from 0 to 2 as l passes N/2, indicating that one can recover the initial entanglement between R and the first qubit if one has access to more than half of the system. The late-time value also agrees with the random matrix calculation in Eq. (39) below.

Therefore,

$$I(R:\bar{B}) + I(R:B) = 2.$$
 (27)

The mutual information can also be cast into a form of relative entropy as follows:

$$I(R:B) = \operatorname{tr}(\rho(RB)(\log \rho(RB) - \log \rho(R) \otimes \rho(B))).$$
(28)

Therefore, I(R : B) is a measure of the difference between the density matrix  $\rho(R) \otimes \rho(B)$  and  $\rho(RB)$ . These density matrices depend on the two orthogonal states  $|\psi_0(t)\rangle$  and  $|\psi_1(t)\rangle$  in the Alice-Bob communication protocol,

$$\rho(RB) = \frac{1}{2} \begin{pmatrix} \rho_0(B) & \rho_{01}(B) \\ \rho_{10}(B) & \rho_1(B) \end{pmatrix},$$

$$\rho(R) \otimes \rho(B) = \frac{1}{4} \begin{pmatrix} \rho_0(B) + \rho_1(B) & 0 \\ 0 & \rho_0(B) + \rho_1(B) \end{pmatrix},$$
(29)

where

$$\rho_{0}(B) = \operatorname{tr}_{\bar{B}} |\psi_{0}(t)\rangle \langle \psi_{0}(t)|,$$
  

$$\rho_{1}(B) = \operatorname{tr}_{\bar{B}} |\psi_{1}(t)\rangle \langle \psi_{1}(t)|,$$
  

$$\rho_{01}(B) = \rho_{01}(B)^{\dagger} = \operatorname{tr}_{\bar{B}} |\psi_{0}(t)\rangle \langle \psi_{1}(t)|.$$

Now we discuss the implication of the minimum and maximum of I(R : B) on the Alice-Bob communication protocol. From the quantum relative entropy, I(R : B) is zero only when  $\rho_{RB} = \rho_R \otimes \rho_B$ , leading to the conditions  $\rho_0(B) = \rho_1(B)$  and  $\rho_{01}(B) = 0$ . As a result,

$$\langle \psi_1 | O_B | \psi_1 \rangle = \langle \psi_0 | O_B | \psi_0 \rangle, \qquad \langle \psi_1 | O_B | \psi_0 \rangle = 0,$$
(30)

for arbitrary operator  $O_B$  within region B, indicating that no operators in B can distinguish the two states. To make the implication of these conditions on states  $|\psi_0\rangle$  and  $|\psi_1\rangle$ manifest, we perform a singular value decomposition on both states between regimes B and  $\overline{B}$ ,

$$\begin{aligned} |\psi_0\rangle &= \sum_n \lambda_0^{(n)} |\psi_0^{n,B}\rangle |\psi_0^{n,\bar{B}}\rangle ,\\ |\psi_1\rangle &= \sum_n \lambda_1^{(n)} |\psi_1^{n,B}\rangle |\psi_1^{n,\bar{B}}\rangle . \end{aligned} \tag{31}$$

The conditions  $\rho_0(B) = \rho_1(B)$  and  $\rho_{01}(B) = 0$  are equivalent to

$$\lambda_{0}^{(n)} = \lambda_{1}^{(n)}, \quad \langle \psi_{0}^{n,B} | \psi_{1}^{m,B} \rangle = \delta_{mn}, \quad \langle \psi_{0}^{n,\bar{B}} | \psi_{1}^{m,\bar{B}} \rangle = 0.$$
(32)

An arbitrary superposition of  $|\psi_0\rangle$  and  $|\psi_1\rangle$  has the same density matrix in region *B* and is thus the same to Bob.

In the opposite limit, I(R:B) takes the maximal value 2 if and only if  $I(R:\overline{B}) = 0$ . This leads to the

same condition as in Eq. (31) with *B* replaced with  $\overline{B}$ . As a result,  $\rho_0(B) = \sum_n \lambda^{(n),2} |\psi_0^{n,B}\rangle \langle \psi_0^{n,B}|$  and  $\rho_1(B) = \sum_n \lambda^{(n),2} |\psi_1^{n,B}\rangle \langle \psi_1^{n,B}|$  act on orthogonal states in the Hilbert space of *B*, and  $|\rho_1(B) - \rho_0(B)| = 2$ , indicating maximal difference between the two states in region *B*. In this case, in principle, the optimal operator  $O_B$  that differentiates the two states can be constructed as

$$O_{B} = \sum_{n} |\psi_{0}^{n,B}\rangle \langle \psi_{0}^{n,B}| - |\psi_{1}^{n,B}\rangle \langle \psi_{1}^{n,B}|, \qquad (33)$$

so that  $\langle \psi_0 | O_B | \psi_0 \rangle = 1$  and  $\langle \psi_1 | O_B | \psi_1 \rangle = -1$ .

What happens when I(R : B) is small but finite? By the quantum Pinsker inequality, we have

$$I(R:B) \ge \frac{1}{2\ln 2} \operatorname{tr} |\rho(RB) - \rho(R) \otimes \rho(B)|^2.$$
(34)

This can be further applied to upper bound any connected correlation between R and A,

$$I(R:B) \ge \frac{\left(\langle O_R O_B \rangle - \langle O_R \rangle \langle O_B \rangle\right)^2}{2\ln 2 \|O_R\|_{\infty}^2 \|O_B\|_{\infty}^2}.$$
 (35)

Applying this inequality to all Pauli operators  $O_R$  leads to

$$\frac{1}{\|O_B\|_{\infty}} |\langle \psi_0| O_B |\psi_0\rangle - \langle \psi_1| O_B |\psi_1\rangle| \le \sqrt{2\ln 2I(R:B)},$$
$$\frac{1}{\|O_B\|_{\infty}} |\langle \psi_0| O_B |\psi_1\rangle| \le \sqrt{2\ln 2I(R:B)}.$$
(36)

Therefore, the smallness of I(R : B) prevents Bob from distinguishing the two states.

# C. Quantum mutual information from random unitary dynamics

As we have seen from comparing the noninteracting and strongly interacting cases (Figs. 3 and 4), the quantum mutual information between reference R and a certain region of system B, I(R : B), depends on the unitary operator U that governs the dynamics. To gain some intuition for I in a strongly interacting system, we next consider a simple toy model where U is taken to be a random unitary operator drawn from the Haar ensemble.

It is important to understand that a random unitary matrix is a generic matrix acting on the Hilbert space and does not obey a local structure sketched in Fig. 2. Nevertheless, random unitary dynamics is a good starting point for thinking about quantum information dynamics. It can be used to mimic the local unitary dynamics in the late-time regime where the initial entanglement between the reference and the system is fully scrambled over all degrees of freedom. We first start with the setup in Eq. (20). By averaging over the Haar ensemble, we obtain an analytical universal formula to describe the late value of the mutual information I(R, B) as the size of *B* increases. The formula shows that I(R, B) drastically increases from 0 to 2 as the size of *B* exceeds half of the system, which agrees with the numerical result on the mixed-field Ising model shown in Fig. 4(a).

We then generalize the setup to an initial mixed state  $\rho$ and demonstrate that the late-time mutual information is dependent on the thermal entropy of  $\rho$ . Specifically, when  $\rho$  is fully mixed, the mutual information remains significantly below its maximal value unless *B* encompasses the entire system. Moreover, the setup involving a fully mixed state is crucial for the Hayden-Preskill protocol, which we introduce in the subsequent section.

## 1. Pure state

To proceed, we consider the Rényi version of the mutual information,

$$I^{(2)}(R:B) = S^{(2)}(R) + S^{(2)}(B) - S^{(2)}(RB),$$
(37)

where  $S^{(2)}$  stands for the Rényi entropy [see the definition in Eq. (A7) in Appendix A 2]. We start with the case without the memory, where the combined state of reference *R* and the system is pure. Recall that the time-evolved state is given in Eq. (20),

$$|\Psi(t)\rangle = \frac{1}{\sqrt{2}} (|0\rangle_R U_{\text{Haar}} |\psi_1\rangle_{\text{SYS}} + |1\rangle_R U_{\text{Haar}} |\psi_0\rangle_{\text{SYS}}).$$

Of course, the Rényi entropy of any particular U is in principle a complicated function of U, but what is analytically tractable is an average of the Rényi entropy over all U sampled from the Haar measure. The averaged Rényi entropy captures well the Rényi entropy of a particular U for large systems because the system size strongly suppresses the variance due to quantum typicality [23].

The Rényi entropy averaged over Haar random unitaries can be calculated using the identities in Eqs. (A17) in Appendix A 4. We have

$$2^{-S(B)} = \mathbb{E}(\operatorname{tr}_{B}(\operatorname{tr}_{R\bar{B}}|\Psi(t)\rangle\langle\Psi(t)|)^{2})$$
  
=  $\frac{1}{2^{2N}-1}\left(\left(\frac{1}{2}-\frac{1}{2^{N}}\right)2^{N+l}+\left(1-\frac{1}{2^{N+l}}\right)2^{2N-l}\right)$   
 $\approx 2^{l-N-1}+2^{-l},$  (38)

where  $\mathbb{E}$  indicates averaging over the Haar ensemble. Since there is no notion of locality associated with  $U_{\text{Haar}}$ , the results depend only on the size l of region B, but not its location in the system. We have taken the limit that  $N \to \infty$  in the last line. The entropy S(RB) is the same as that of the complement of *B*,  $S(\overline{B})$ , and can be obtained by replacing *l* with N - l in S(B).

Putting these results together, we obtain the Rényi mutual information as

$$I^{(2)}(R:B) = 1 + \log_2\left(2 - \frac{3(1 - 2^{2l - 2N})}{2 + 4^{l - N/2}}\right), \quad (39)$$

where *l* is the number of qubits in region *B*. The mutual information is 1 when l = N/2, i.e., the region *B* occupies half the system. It increases to its maximal value of 2 as *l* exceeds N/2 and decreases to its minimal value of 0 exponentially when *l* is less than N/2. This result indicates that, when the initial information is fully scrambled, any portion of the system less than half does not contain the initial information. On the other hand, any portion larger than half is maximally entangled with the reference spin *R* and can be used to recover the initial information. This is in sharp contrast with the nonscrambling SWAP circuit shown in Fig. 3 where the information is located at a specific qubit at a given time.

The change of the mutual information at l = N/2 can be understood as follows. The system contains two disconnected regions with l < N/2, but the reference spin cannot be simultaneously maximally entangled with two nonoverlapping regions; otherwise, it would violate the monogamy of entanglement. As a result, the mutual information drastically reduces when l < N/2.

This result is also valid in the late-time regime of local unitary dynamics or even Hamiltonian dynamics. As shown in Fig. 4(b), for the mixed-field Ising model, the late-time value of mutual information for different regions obeys the random unitary calculation, drastically increasing from 0 to 2 when the region considered exceeds half of the system. In the large-N limit, I increases from 0.033 to 1.967 in a window of six qubits around N/2.

#### 2. Mixed state

This setup can be generalized to the case where the two initial states of the system are mixed. Similar to the previous case, we pick the two initial states

$$\rho_0 = |0\rangle \langle 0| \otimes \rho, \qquad \rho_1 = |1\rangle \langle 1| \otimes \rho, \qquad (40)$$

where  $\rho$  is the density matrix of the system, excluding the first qubit. Then we introduce the purification of the density matrix  $\rho$  by including an additional auxiliary system called memory. The purification of  $\rho$ , denoted  $|\sqrt{\rho}\rangle$ , is a pure state living in the Hilbert space of the system (excluding the first qubit) and the memory. It has the property that tracing out the degrees of freedom in the memory gives

back the density matrix  $\rho$ ,

$$\rho = \operatorname{tr}_{\mathrm{MEM}} |\sqrt{\rho}\rangle \langle \sqrt{\rho} | \,. \tag{41}$$

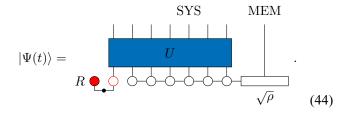
Alice's qubit still forms a Bell pair with the reference. The entire state, including the reference, the system, and the memory, is

$$|\Psi\rangle = \frac{1}{\sqrt{2}} (|0\rangle_R |0\rangle_{q_A} + |1\rangle_R |1\rangle_{q_A}) |\sqrt{\rho}\rangle.$$
 (42)

This state is very similar to Eq. (19). The difference is that  $|\sqrt{\rho}\rangle$  contains degrees of freedom of the memory that the time evolution operator does not act on. The time evolution of the state is

$$\begin{split} |\Psi(t)\rangle &= \frac{1}{\sqrt{2}} (|0\rangle_R \, U_{\text{SYS}} \otimes I_{\text{MEM}} \, |0, \sqrt{\rho}\rangle \\ &+ |1\rangle_R \, U_{\text{SYS}} \otimes I_{\text{MEM}} \, |1, \sqrt{\rho}\rangle). \end{split}$$
(43)

Graphically,



A similar calculation of the averaged Rényi entropy now yields

$$2^{-S_B} \approx \left(\frac{\mathrm{tr}\rho^2}{2} - \frac{1}{2^N}\right) 2^{l-N} + 2^{-l},$$

$$2^{-S_{RB}} \approx \left(\mathrm{tr}\rho^2 - \frac{1}{2^{N+1}}\right) 2^{l-N} + 2^{-l-1}.$$
(45)

Therefore, the mutual information is

$$I^{(2)}(R:B) = 1 + \log_2\left(2 - \frac{3(1 - 2^{2l - 2N})}{2 + (2^{-s} - 2^{-N+1})4^{l - N/2}}\right),$$
(46)

where  $s(\rho)$  is the Rényi entropy of  $\rho$  in the initial state. It reduces back to Eq. (39) when the initial state is pure, i.e., s = 0. When the initial state is mixed and has finite entropy, I(l) drastically increases from 0 to its maximal value 2 at  $l \sim (N + s)/2$ . The behavior of I(l) for l near (N + s)/2 is independent of s in the large-N limit and is already shown in Fig. 4(b). This result suggests that one now needs more than (N + s)/2 qubits to recover the initial information.

#### 3. Maximally mixed state

In the special case when the initial state is maximally mixed,  $s(\rho) = N - 1$  (the first spin is entangled with the reference spin). In this case,  $I = \log_2(1 + 3 \times 2^{2l-2N})$  and equals 2 when l = N, i.e., the entire system. The mutual information  $I^{(2)}$  drastically decreases as the accessible region contains fewer qubits. These results also apply to the late-time regime of unitary dynamics. Denote the part of the system that Bob does not have access to as *E*, and the number of qubits in this region as |E|. We have

$$I^{(2)}(R: SYS - E) = \log_2(1 + 3 \times 4^{-|E|}).$$
(47)

This behavior is plotted in Fig. 5.

In this special case, the Rényi mutual information can be used to bound the von Neumann entropy because the Rényi entropy  $S^{(2)}(SYS - E)$  and  $S^{(2)}(R)$  are maximal and equal their von Neumann counterparts. Since  $S^{(2)}(R \cup (SYS - E)) \leq S(R \cup (SYS - E))$ , we have

$$I(R: SYS - E) \le I^{(2)}(R: SYS - E)$$
  
= log<sub>2</sub>(1 + 3 × 4<sup>-|E|</sup>). (48)

Combining with I(R : SYS) = 2, this indicates that recovering the initial information requires accessing the entire system when the initial state is maximally mixed, in contrast with half of the system for the pure initial state.

Another closely related diagnosis of scrambling in this special setup is the tripartite mutual information [12,24] given by

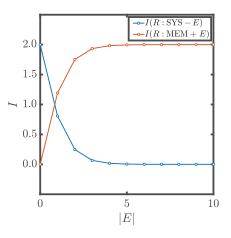


FIG. 5. In a fully mixed initial state setup, after applying the Haar random unitary to the system, the mutual information between the reference and the system excluding a few qubits, I(R : SYS - E) (blue curve), decays exponentially with the number of qubits excluded. On the contrary, the mutual information between the reference and the memory plus a few qubits in the system, I(R : SYS + E) (red curve), saturates to the maximum.

$$I_3(R:E:SYS - E) = I(R:E) + I(R:SYS - E)$$
  
-  $I(R:SYS),$  (49)

which quantifies how much information is not encoded in E and SYS – E but in their union. The last term is 2 for all times. For a scrambled system, the first two terms are zero, the tripartite mutual information takes the minimal value -2.

# IV. HAYDEN-PRESKILL PROTOCOL: DETECTING AND RECOVERING THE INFORMATION

Having understood where the information resides at late times, we now turn our attention to understanding how to retrieve this information. We start with the fully mixed case as this was the setup first studied in the Hayden-Preskill protocol [7], which was originally proposed for black hole dynamics, but is applicable to generic quantum many-body systems.

The previous section introduces MEM as a mathematical tool to purify the mixed state. Now let us also consider MEM as a physical system that someone may have access to, but it does not evolve by the unitary dynamics. In the black hole problem, the memory is supposed to describe previously emitted Hawking radiation that is entangled with the remaining system. The maximally mixed case then corresponds to the entropic midpoint of the black hole's evolution, where the remaining black hole is maximally entangled with the emitted Hawking radiation and is described by a fully mixed state. The question posed in Ref. [7] was: if another qubit is thrown into the black hole then how long does one have to wait until the information in that qubit is available again in the subsequent radiation? The result of the prior section is that, when Bob has access to the memory (the early radiation), then the information in the fresh qubit quickly becomes available again.

As shown in Fig. 5, when the accessible region of Bob is SYS – E, the mutual information drastically reduces to zero as the number of qubits in E increases, indicating that Bob needs the entire system to recover the information. What happens when Bob also has access to the qubits in the memory? The memory is only coupled to the system by the initial entanglement. The unitary dynamics never mix the system and memory. As a result, one expects that the memory does not contain any information about the reference, indicated by zero mutual information. On the other hand, since the quantum state on the composite system, including the reference, system, and memory, is pure, we have

$$I(R : SYS - E) + I(R : MEM + E) = 2.$$
 (50)

Thus, according to Eq. (48),

$$I(R: MEM + E) \ge 2 - \log_2(1 + 3 \times 4^{-|E|});$$
 (51)

I(R : MEM + E) saturates to the maximal value 2 exponentially fast (Fig. 5, red curve). In other words, remarkably, Bob can recover the initial information provided that he has access to the full memory, which does not contain any information about the reference [I(R : MEM) = 0 for all time], and a few qubits *E* in the system.

#### A. Detecting the information front

The Hayden-Preskill protocol also provides an operational way to measure information dynamics. Let us still consider the setup with a fully mixed initial state. The unitary operator U governing the dynamics may now have more structure, such as locality, instead of being Haar random. Under time evolution, the mutual information between reference R and system SYS is always maximal, and the mutual information between the reference and MEM is always 0. Initially, R is maximally entangled with the first spin. As discussed previously, when the unitary operator is Haar random, the entanglement instantly spreads over the entire system. The mutual information between the reference to any subregion of the time-evolved system, even excluding a few qubits, becomes almost 0.

Now suppose that the unitary operator has a local structure built in, so that the entanglement spreads in a time-dependent manner, as shown in Fig. 4. In order to capture the profile of the entanglement spreading, a natural approach is to trace the mutual information between the reference and the first *n* qubits in the output  $I(R : q_1 \cdots q_l)$ , which increases monotonically with l. The front of the spreading at a given time is determined by the largest *l* for which  $I(R:q_1\cdots q_l) < 2$ , because it implies that the information has begun to leak out the first l qubits. This approach is a conceptually straightforward application of the definition, but it is challenging for experiments since  $I(R: q_1 \cdots q_n)$  involves the entanglement entropy of a nonlocal region. Alternatively, inspired by the Hayden-Preskill setup, one can measure the mutual information between the reference and the memory plus one qubit in the system  $I(R : \text{MEM} \cup q_n)$ . A finite value indicates that the initial information has reached the *n*th spin. Therefore, the front of the entanglement spreading can be determined by the largest *n* that  $I(R : \text{MEM} \cup q_n) > 0$ .

The mutual information used in the second approach can be estimated by local measurement, as we discuss below, and is much more accessible in experiments. Now let us look into the mutual information  $I(R : \text{MEM} \cup q_n)$  more closely. From the definition of I, we have

$$I(R : MEM \cup q_n)$$
  
=  $S(R) + S(MEM \cup q_n) - S(R \cup MEM \cup q_n)$   
=  $S(R) + S(MEM \cup q_n) - S(SYS - q_n).$  (52)

In the second equality, we used the fact that  $R \cup \text{MEM} \cup q_n$  and  $\text{SYS} - q_n$  are complementary regions in a pure

state. The first term, S(R), is always 1 because the reference qubit is not affected by the dynamics. The last term,  $S(SYS - q_n)$ , attains its maximal value N - 1 because the system stays fully mixed under the unitary dynamics, as does any subsystem. Therefore, the mutual information depends only on the entanglement entropy of MEM  $\cup q_n$ ,

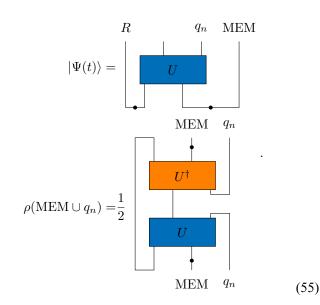
$$I(R: \text{MEM} \cup q_n) = 2 - N + S(\text{MEM} \cup q_n),$$
$$N - 2 \le S(\text{MEM} \cup q_n) \le N.$$
(53)

Initially,  $S(\text{MEM} \cup q_n)$  takes the minimal value for all qubits  $q_n$  except the first one. As time increases, its deviation from the minimal value signals that the information has reached the *n*th spin. Since the von Neumann entropy upper bounds the Rényi entropy, we have

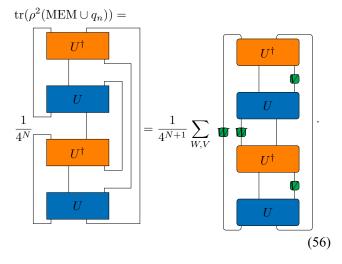
$$I(R: \text{MEM} \cup q_n) \ge 2 - N + S^{(2)}(\text{MEM} \cup q_n).$$
 (54)

As a result,  $S^{(2)}$  can be used as an indicator for the front of quantum information propagation. This result is particularly nice, since  $S^{(2)}$  can be related to correlation functions that are more accessible than the von Neumann mutual information, as we show below. See also the experimental schemes in Sec. VII below.

Since in the case we consider here, the qubits  $q_{2\sim N}$  are in a fully mixed state, we can choose the simplest purification where the memory contains N - 1 auxiliary qubits that form N - 1 EPR pairs with the N - 1 spins in the system. The time-evolved state and the density matrix  $\rho$  (MEM  $\cup q_n$ ) is



From the density matrix, we can obtain the purity  $\operatorname{tr}(\rho^2(\operatorname{MEM} \cup q_n)) = e^{-S^{(2)}(\operatorname{MEM} \cup q_n)}$  as



The operators V and W are summed over local Pauli operators (including the identity) on  $q_1$  and  $q_n$ . The second equality uses the completeness relation of Pauli operators in Eq. (A16) in Appendix A 2 b. When  $W_1$  or  $V_n$  equals the identity, the trace contributes 1 to the sum. Separating these terms from the others, we get

$$\operatorname{tr}\rho^{2}(\operatorname{MEM} \cup q_{n}) = \frac{1}{2^{N+2}} \bigg( 7 + \sum_{W_{1} \neq I, \ V_{n} \neq I} \frac{1}{2^{N}} \operatorname{tr}(W_{1}(-t)V_{n}W_{1}(-t)V_{n}) \bigg).$$
(57)

The purity becomes a sum of correlators between timeevolved local Pauli operators. The Rényi entropy is just  $-\log_2 \operatorname{tr}(\rho^2)$ . From Eq. (57), we can bound the mutual information by the correlators,

$$P(R: MEM \cup q_n) \ge 4 - \log_2 \left(7 + \sum_{W_1 \neq I, V_n \neq I} \frac{1}{2^N} tr(W_1(-t)V_nW_1(-t)V_n)\right).$$
(58)

We emphasize that this inequality applies to any unitary U. Each term in the summation has a maximum value 1, in which case the right-hand side takes the minimal value 0. This happens when the Heisenberg operator  $W_1(-t)$  commutes with the operator  $V_n$  for all V and W. When  $W_1(-t)$ and  $V_n$  start to overlap, the correlator decreases from 1. As a result,  $I(R : \text{MEM} \cup q_n)$  is nonzero, indicating that the information has reached  $q_n$ . At late time, all the terms decay to 0 and the right-hand side becomes  $4 - \log_2 7$ , consistent with Eq. (51) when |E| = 1.

Because the kind of correlator appearing on the righthand side of Eq. (58) can be used to diagnose the information propagation, it deserves a name. In the literature, it is referred to as the out-of-time-ordered correlator, usually written as

$$F(t) = \frac{1}{2^N} \operatorname{tr}(W_1(t) V_n W_1(t) V_n).$$
(59)

The time argument is changed from -t to t for simplicity. Larkin and Ovchinnikov first introduced the OTOC in the context of superconductivity [25]. It has gained extensive renewed interest recently due to the connection to scrambling dynamics discussed here as well as quantum many-body chaos in the semiclassical regime (see the short discussion in Sec. VIII below).

Several remarks are in order. First, we emphasize that the OTOC is only an indicator of information propagation and scrambling. When  $W_1(t)$  and  $V_n$  do not commute, it indicates that the information front has reached the qubit n, but does not mean that one can recover Alice's initial action by measuring the *n*th qubit. In fact, we have seen from the random unitary calculation that it requires all the qubits in the system, or the entire memory plus a few qubits with nonzero commutators, to recover the information. Second, the OTOC as an indicator of the information propagation is tied to the simple initial state where  $\rho(q_2 \sim N)$  is fully mixed, namely, the Hayden-Preskill setup. When the initial state is not fully mixed, a precise relation between the mutual information, which is the defining characterization of information propagation, and simple correlation functions such as the OTOC is not established. One can even show that the OTOC overestimates how fast information propagates for some other initial states.

# B. Recovering the information: many-body teleportation

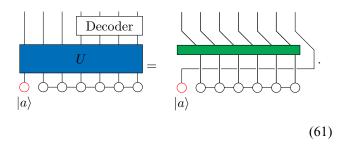
This section discusses the second question of quantum information dynamics on how to recover the initial information. In Sec. III B, we showed that maximal mutual information between the reference and Bob's qubits indicates that Bob can distinguish Alice's action on the initial state using the operator constructed in Eq. (33). However, in practice, the optimal operator  $O_B$  is nonlocal and challenging to construct. It will be ideal if the initial state of Alice's qubit  $q_A$  reappears on one of Bob's qubits after Bob follows a specific decoding protocol on his qubits. This is called many-body teleportation [7,26–28].

### 1. Requirement for many-body teleportation

Many-body teleportation works if

$$U(|a\rangle_{q_A} \otimes |\psi\rangle) \stackrel{\text{Bob's decoding}}{\Longrightarrow} |\phi\rangle \otimes |a\rangle_{q_B} \qquad (60)$$

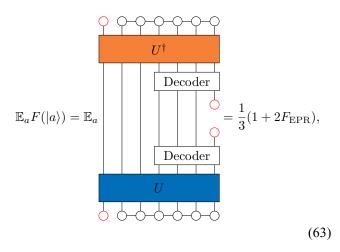
for any state  $|a\rangle$  to be teleported, where  $q_A$  is Alice's qubit,  $q_B$  is one of Bob's qubits,  $|\psi\rangle$  is the initial state on the qubits except  $q_A$ , and  $|\phi\rangle$  is the final state on the qubits except  $q_B$ . This means that Alice would be able to teleport a bit of quantum information to Bob through a strongly interacting medium described by U, which fully scrambles her information into the entire system. The key point here is that Bob only owns part of the qubits, so he cannot trivially unscramble the information by applying  $U^{\dagger}$ . Graphically, the condition for successful teleportation is



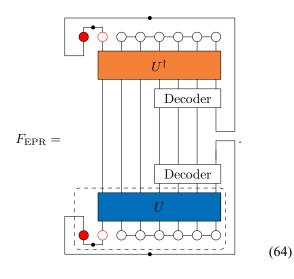
Denote the quantum state after the decoding  $|\Psi_{out}\rangle$ ; the fidelity of teleporting state  $|a\rangle$  to qubit  $q_B$  is defined as

$$F(|a\rangle) = \langle \Psi_{\text{out}} | (|a\rangle \langle a|)_{q_B} \otimes I | \Psi_{\text{out}} \rangle.$$
(62)

When the fidelity averaged over Alice's state  $\mathbb{E}(F(|a\rangle))$  is 1, it indicates that the system is able to teleport any quantum state with perfect fidelity. The averaged fidelity can be obtained by sampling Alice's state from the action of a random unitary  $u_a$  on a basis state  $|0\rangle$ ,



where



Recall the boxed setup with reference *R* used in Eqs. (20) and (43). Here  $F_{\text{EPR}}$  is just the fidelity that reference *R*, initially forming an EPR pair with Alice's qubit  $q_1$ , forms an EPR pair with one of Bob's qubits after the unitary transformation and decoding;  $F_{\text{EPR}} = 1$  indicates perfect teleportation. In this case, the mutual information between the reference and Bob's qubit  $q_B$  is 2. The necessary condition for perfect teleportation is that the mutual information between the reference and all of Bob's qubits I(R : B) is 2, as discussed in Sec. III A. The purpose of the decoder is to concentrate the entanglement with reference to a single qubit  $q_B$ . Even given maximal mutual information I(R : B), it is still in general challenging to design the decoding protocol.

#### 2. Conventional teleportation

Before discussing the decoder for the Hayden-Preskill protocol, it is useful to review conventional quantum teleportation. Alice has a qubit encoding the state to be teleported  $|a\rangle$ . Alice has an additional qubit that forms an EPR state with Bob's qubit. To teleport state  $|a\rangle$  to Bob's qubit, Alice first measures her two qubits in the Bell basis. The measurement projects the two qubits into one of four Bell states:

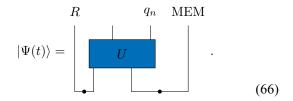
$$|I\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle), \qquad |X\rangle = \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle),$$
$$|Y\rangle = \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle), \qquad |Z\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle).$$
(65)

Those states obey that  $I \otimes \sigma^x |X\rangle = I \otimes \sigma^y |Y\rangle = I \otimes \sigma^z |Z\rangle = |I\rangle$ . After the measurement, Alice tells Bob the Bell state  $|s\rangle$  she obtained. Each of the four possibilities appears with probability 1/4. Based on Alice's message, Bob applies the corresponding Pauli operator to his qubit

or does nothing if the Bell state is  $|I\rangle$ . Then state  $|a\rangle$  appears on Bob's qubit. Taking all the four measurement outcomes into account, the final state is a mixed state  $\frac{1}{4}I_{Alice} \otimes |a\rangle \langle a|_{Bob}$ , in which Alice's two qubits are in the fully mixed state and Bob has Alice's original state. One can also show that if we introduce an addition reference qubit forming an EPR pair with Alice's first qubit and perform the same protocol, the reference will form an EPR pair with Bob's qubit in the final state with fidelity  $F_{EPR} = 1$ .

#### 3. Many-body teleportation

Now we are ready to discuss the decoding protocol for the Hayden-Preskill protocol [26]. Recall the setup for the Hayden-Preskill protocol. The system contains Nqubits. The first one forms an EPR pair with reference R; the remaining N-1 qubits form N-1 EPR pairs with another N - 1 auxiliary qubits in the memory, which Bob owns. In addition, Bob also owns a set of qubits E in the system, so that  $B = MEM \cup E$ . In Sec. IV, we showed that I(R:B) approaches 2 exponentially fast as |E| increases when the system is time evolved into the late-time regime and fully scrambled; therefore, a decoding protocol is in principle possible. The question now is how to design a decoding protocol on Bob's protocol so that R forms an EPR pair with one of Bob's qubits. The quality of the decoding protocol can be characterized by  $F_{\text{EPR}}$ . In the Hayden-Preskill protocol, the state before the decoding is

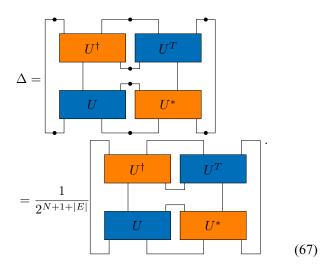


Yoshida and Kitaev [26] found two decoding protocols for this state, one probabilistic and one deterministic. The probabilistic decoding protocol goes as follows.

- (1) Bob takes another two qubits,  $q'_1$  and R', and prepares them in an EPR state.
- (2) Bob applies the unitary operator  $U^*$  to MEM  $\cup q'_1$ .
- (3) Bob performs the Bell measurement on each qubit in *E* and its partner in MEM, with which it forms an EPR pair initially.
- (4) The entire protocol is repeated, including preparing the state in Eq. (66) until the outcome of all the Bell measurements is the EPR state  $|I\rangle$ .

After these step, one of Bob's new qubits, R', and the reference qubit R would have high fidelity to form an EPR pair. This implies that in the case without the reference, any state injected into  $q_1$  initially has a high fidelity to reappear on R', Bob's other new qubit. This protocol is probabilistic

because in step (3) Bob needs to postselect the EPR pairs from the Bell measurements. To understand this decoder, let us calculate the probability of successful postselection and the fidelity  $F_{\text{EPR}}(R, R')$  given successful postselection. The probability of successful postselection is



Note that the diagram is the same as that in Eq. (56) for calculating  $\rho^2(\text{MEM} \cup E)$ . Using  $S^{(2)}(R) = 1$  and  $S^{(2)}(R \cup \text{MEM} \cup E) = N - |E|$ , we have

$$\Delta = 2^{N-1-E} \operatorname{tr}^2(\rho_{\text{MEM}\cup E}) = 2^{-I^{(2)}(R:\text{MEM}\cup E)}.$$
 (68)

The probability of the postselection is directly related to the Rényi mutual information between R and Bob's qubit before the decoding. Given the successful postselection, the fidelity that R and R' form an EPR pair is

$$F_{\rm EPR} = \frac{1}{4\Delta} = 2^{I^{(2)}(R:\rm MEM\cup E)-2}.$$
 (69)

When the unitary U is fully scrambling, such as a Haar random unitary, the Rényi mutual information can be obtained from Eq. (47) as

$$I^{(2)}(R: MEM \cup E) = 2 - I^{(2)}(R: SYS - E)$$
  
= 2 - log<sub>2</sub>(1 + 3 × 4<sup>-|E|</sup>). (70)

Substituting this into the equation for  $F_{\text{EPR}}$ , we get

$$F_{\rm EPR} = \frac{1}{1 + 3 \times 4^{-|E|}}.$$
 (71)

The fidelity  $F_{\text{EPR}}$  approaches 1 exponentially fast as |E| increases, indicating perfect teleportation fidelity given successful postselection for fully scrambling unitary time evolution. In general, since  $\text{tr}^2 \rho^2 (\text{MEM} \cup E)$  can be written as the sum of OTOCs, as shown in Eq. (56), the fidelity

 $F_{\rm EPR}$  is also directly related to OTOCs as

$$F_{\rm EPR} = \left(\frac{1}{4^{|E|}} \frac{1}{2^N} \sum_{W_1, V_E} \operatorname{tr}(W_1(-t)V_E W_1(-t)V_E)\right)^{-1},\tag{72}$$

where  $V_1$  and  $W_E$  are summed over all local operators, including the identity, in the first qubit and E, respectively. In the fully scrambled regime, the OTOC is 1 if either  $V_1$ or  $W_E$  is the identity and 0 otherwise, and we get Eq. (71) back. We see that the OTOC not only provides a tool to detect information propagation, but is also directly related to the fidelity of information recovery.

The above decoding protocol is probabilistic. Bob has to postselect the state on  $E \cup E'$  to be the EPR state after the Bell measurement. The probability  $\Delta$  is directly related to the Rényi mutual information between the reference and Bob,  $\Delta = 2^{-I^{(2)}(R:\text{MEM}\cup E)}$ . The optimal fidelity  $F_{\text{EPR}}$ requires the minimal  $\Delta$ , which is 1/4 for teleporting a single qubit and  $4^{-n}$  for teleporting multiple qubits. To overcome the small successful postselection probability, Yoshida and Kitaev [26] also designed a deterministic decoder that does not require postselection but only unitary transformation on Bob's qubits. The general idea is to perform a Grover search on Bob's qubits. After a series of unitary transformations, the states in which *E* does not form an EPR pair with *E'* are canceled due to destructive interference.

The above discussion assumes that the reference, system, and memory form a closed system and that the dynamics is unitary. It would also be interesting to study how dissipation and measurement affect the recovery fidelity [29,30].

## V. MICROSCOPIC PHYSICS OF OTOCS AND OPERATOR GROWTH

## A. Relating the OTOC to operator dynamics

In previous sections, we established that OTOCs can be used to track the front of the information dynamics, and they are directly related to the fidelity to recover the initial information. In this section, we discuss the microscopic physics of OTOCs regarding the growth of Heisenberg operators [31,32].

We consider a quantum many-body system whose dynamics is governed by a unitary operator U(t). We use  $W_0(t) = U^{\dagger}(t)W_0U$  to represent the time-evolved Heisenberg operator, which is located at the origin at t = 0, and use  $V_r$  as a static local operator at site r. Then the OTOC between W(t) and  $V_r$  can be written as

$$F(r,t) = \frac{1}{\mathrm{tr}I} \mathrm{tr}(W_0^{\dagger}(t)V_r^{\dagger}W_0(t)V_r).$$
(73)

The space and time dependence of F(r, t) has a very intuitive picture based on operator growth, sensitive to whether the support of W(t) overlaps with that of  $V_r$ . The connection between the OTOC and operator growth can be made more explicit by introducing the squared commutator

$$C(r,t) = \frac{1}{\text{tr}I} \text{tr}([W_0(t), V_r)]^{\dagger}[W_0(t), V_r])$$
  
=  $\frac{1}{\text{tr}I} ||W_0(t), V_r||_2^2,$  (74)

which is proportional to the Frobenius norm of the commutator  $[W_0(t), V_r]$  and thus always positive. One can easily show that

$$C(r,t) = -F(r,t) - F^{*}(r,t) + \frac{1}{\text{tr}I} [\text{tr}(W_{0}^{\dagger}(t)V_{r}^{\dagger}V_{r}W_{0}(t)) + \text{tr}(V_{r}^{\dagger}W_{0}^{\dagger}(t)W_{0}(t)V_{r})].$$
(75)

The last two terms are local observables that typically relax to a constant quickly due to thermalization. Therefore, C(r, t) and F(r, t) have the same behavior after a thermalization time. In particular, if both W and V are unitary, the sum of the last two terms is 2. Furthermore, when W and V are Hermitian, F(r, t) is real. Therefore, when W and V are unitary and Hermitian, e.g., Pauli operators, C(r, t) = 2 - 2F(r, t).

The squared commutator manifestly depends on the "size" of the Heisenberg operator W(t), the number of degrees of freedom that W(t) acts on. At t = 0, W(t) acts only on a simple site and commutes with any  $V_r$  that is far away, so C(r, t) = 0. (In a fermionic system, if the operators W and V are both fermionic, one should consider a squared anticommutator instead of a squared commutator.) As time increases, W(t) becomes more and more nonlocal and starts to overlap with  $V_r$ , indicated by the increase of C(r, t). Varying  $V_r$  for different r, C(r, t) remains small if  $V_r$  is outside the support of W(t) and large otherwise. Therefore, C(r, t) probes the size of the Heisenberg operator W(t) at a given time. This is consistent with the discussion in Sec. III A that the OTOC tracks the light cone of information dynamics for the infinite-temperature state.

To obtain a more precise understanding, it is useful to think about the growth of the Heisenberg operator W(t) in a complete basis of operators  $\{S\}$ . The basis obeys the following normalization and completeness conditions:

$$\frac{1}{\mathrm{tr}I}\mathrm{tr}(\mathcal{S}^{\dagger}\mathcal{S}') = \delta_{\mathcal{S}\mathcal{S}'}, \qquad \frac{1}{\mathrm{tr}I}\sum_{\mathcal{S}}\mathcal{S}_{ab}^{\dagger}\mathcal{S}_{cd} = \delta_{ad}\delta_{bc}.$$
 (76)

The conventional choice of the operator basis for qubit systems is the Pauli strings, which are products of Pauli operators or the identity operator on every site,

$$S = \prod_{r=1}^{N} \sigma_r^{(s)},\tag{77}$$

where s = 0, 1, 2, 3 stands for  $I, \sigma^x, \sigma^y$ , and  $\sigma^z$ . Note that there are in total  $4^N$  different Pauli strings for N qubits.

The Heisenberg operator W(t) can be expanded in the basis

$$W(t) = \sum_{S} \alpha(S, t) S.$$
(78)

Without loss of generality, we fix the norm of W(t) to be

$$\frac{1}{\mathrm{tr}I}\mathrm{tr}(W^{\dagger}W) = 1.$$
(79)

Since the time evolution is unitary, the normalization stays the same and  $\sum_{S} |\alpha(S, t)|^2 = 1$  for all time. Hence,  $|\alpha(S, t)|^2$  can be interpreted as a probability distribution of the operator.

We also define a complete local operator basis at site r, denoted  $S_r$ . In qubit systems, the local basis includes three Pauli operators and the identity operator. The generalization to qudit systems and fermion systems is discussed in Appendix B. In the general case, the local Hilbert space dimension is q, and there are  $q^2$  local operators in the local operator basis.

Using the completeness relation, one can show that the average OTOC between W(t) and all  $S_r$  is

$$\frac{1}{q^2} \sum_{\mathcal{S}_r} \frac{1}{\text{tr}I} \text{tr}(W^{\dagger}(t)\mathcal{S}_r^{\dagger}W(t)\mathcal{S}_r) = \sum_{\mathcal{S}_r=I} |\alpha(\mathcal{S},t)|^2, \quad (80)$$

which measures the probability that W(t) acts on site *r* as the identity operator, according to the probability distribution  $|\alpha(S, t)|^2$ . Similarly, the average squared commutator is

$$\frac{1}{q^2} \sum_{S_r} \frac{1}{\text{tr}I} \| [W(t), S_r] \|_2^2 = 2 \sum_{S_r \neq I} |\alpha(S, t)|^2, \qquad (81)$$

which is complementary to the average OTOC. Equation (80) establishes a precise connection between the OTOC and operator probability. In the summation in Eqs. (80) and (81), the term with  $S_r = I$  does not have dynamics and can be separated from the other terms. Therefore we define the following average OTOC and

squared commutator:

$$\mathcal{F}(r,t) = \frac{1}{q^2 - 1} \frac{1}{\operatorname{tr} I} \sum_{\mathcal{S}_r \neq I_r} \operatorname{tr}(\mathcal{W}^{\dagger}(t) \mathcal{S}_r^{\dagger} \mathcal{W}(t) \mathcal{S}_r)$$
$$= \frac{q^2}{q^2 - 1} \sum_{\mathcal{S}_r = I_r} |\alpha(\mathcal{S}, t)|^2 - \frac{1}{q^2 - 1}, \qquad (82a)$$

$$C(r,t) = \frac{1}{q^2 - 1} \sum_{S_r \neq I_r} \frac{1}{\text{tr}I} \| [W(t), S_r] \|_2^2$$
$$= \frac{2q^2}{q^2 - 1} \sum_{S_r \neq I_r} |\alpha(S,t)|^2.$$
(82b)

Starting with a simple local operator, the time evolution of the operator probability distribution can be very different between generic interacting systems and noninteracting systems. For instance, in noninteracting fermionic systems, a single-particle operator always remains single particle. On the contrary, in interacting systems, a local operator tends to become as complex as possible in the late time. Maximal complexity for an initial traceless operator implies uniform distribution over the operator basis S except for the identity. The identity operator is special since it is static under unitary time evolution. Therefore, in systems with L qudits with local dimension q, we have

$$\lim_{t \to \infty} |\alpha(\mathcal{S}, t)|^2 = \frac{1}{q^{2L} - 1} (1 - \delta_{\mathcal{S}, I}).$$
(83)

Based on Eqs. (82a) and (82b), the late-time operator probability distribution determines the late-time value of the average OTOC

$$\lim_{t \to \infty} \mathcal{F}(r,t) = -\frac{1}{q^{2L} - 1} \approx 0$$
(84)

and

$$\lim_{t \to \infty} \mathcal{C}(r, t) = 2\left(1 + \frac{1}{q^{2L} - 1}\right) \approx 2.$$
(85)

One can obtain the same late-time values of the OTOC between W(t) and a single local operator  $S_r$  using the Haar random unitary as the time evolution operator. These late-time values suggest that an operator becomes maximally complex and can be used to distinguish generic interacting many-body systems from noninteracting systems. We also note that the discussion above assumes no symmetry present. We briefly discuss how symmetries impact scrambling dynamics in Sec. VIII below.

#### **B.** Scrambling dynamics in geometrically local systems

In Haar random unitary dynamics, there is no notion of space and time since  $\mathcal{F}(r, t)$  reaches its final value in

a single step for every r. In contrast, a physical manybody system only contains few-body interactions, such as spin-spin interactions or interaction terms involving four fermionic operators. A generic physical Hamiltonian of Nsites is

$$H = \sum_{b} J_b(t) H_b, \tag{86}$$

where  $H_b$  acts on a set of sites by b and  $J_b$  is the coupling strength that can generally be time dependent. For example, in spin chains with the nearest-neighboring interaction, b labels each bond. In general, one can regard a quantum many-body system as a hypergraph, in which each site defines a vertex and each term in the Hamiltonian  $H_b$ defines a hyperedge  $e_b$  to be the set of vertices involved in  $H_b$ . This hypergraph completely determines the connectivity of the system. One can also generalize this description to unitary circuits, which are a tensor product of few-body unitaries,

$$U = \prod_{b} U_{b}.$$
 (87)

From the OTOC one can define a time scale  $t^*$  called the scrambling time, at which  $\mathcal{F}(r, t)$  relaxes to the final value of about 0 for all sites r, given a Heisenberg operator W(t)that is initially localized at one site. A natural question in this general setup is how the scrambling time  $t^*$  depends on the system size. The time scale is determined by the hypergraph's connectivity and the coupling strength  $J_b$ . While the complete answer to this question is not known, there exists extensive literature tackling specific regimes. We also note that, although this general definition of a physical quantum many-body system seems obscure and unnecessarily complex from a conventional point of view, there currently exist experimental schemes allowing for tuning the connectivity between microscopic degrees of freedom [33], making quantum many-body systems with general graph structure a physically relevant topic.

To make our discussion tangible, let us restrict to cases where  $H_b$  acts only on two sites, describing a spin-spin interaction, for instance. In this case, the hypergraph reduces to a graph. One can imagine arranging the N qubits in a chain. The Hamiltonian can be written as

$$H = \sum J_{rr'}^{\alpha\beta}(t)\sigma_r^{\alpha}\sigma_{r'}^{\beta}.$$
 (88)

The strongest connectivity occurs when the graph is a complete graph where each qubit connects to every qubit with approximately the same interaction strength J (exactly the same coupling between all pairs makes the model integrable and not scrambling). In such all-to-all connected models, it is typically found that the scrambling time scales logarithmically with N,  $t^* \sim \log N$ , provided the couplings  $J_b$  are normalized to give an extensive-in-N energy. Furthermore, the system is approximately permutation invariant in this case, so all qubits are equivalent. Therefore, C(r, t) does not have spatial dependence. Many calculations [11,34–36] have shown that in the large-N limit, C(r, t) in the early time takes an exponential growth form

$$C(r,t) \sim \frac{1}{N} \exp(\lambda t),$$
 (89)

where  $\lambda$  is called the Lyapunov exponent related to the coupling strength J. Here  $\mathcal{C}(r, t)$  becomes  $\mathcal{O}(1)$  when  $t \sim$  $\log N$ , setting the scrambling time scale. It is suggested that this is the fastest scrambling time scale for physical systems with a proper normalization of  $J_{rr'}$  [8,37]. Significant research interest has been put into imposing bounds on the Lyapunov exponent to understand how fast quantum many-body systems can process information. The chaos bound [11] concerns a finite-temperature version of the OTOC and shows that  $\lambda < 2\pi T$  in a quite general setting, leading to extensive studies on understanding and refining the bound and on operator growth in general, especially at finite temperature, e.g., [31,32,38-41]. We briefly comment on the notion of OTOCs at finite temperature and their connection to many-body quantum chaos in Sec. VIII below.

Moving away from all-to-all connected models, one should expect the scrambling time to increase as one reduces the connectivity since it will take longer for a local perturbation to spread over the system. One approach to reducing the connectivity is to require that  $J_{rr'}$  decreases as a function of the distance |r - r'| (see also the discussion about scrambling on sparse graphs [42-44]). As a result,  $\mathcal{C}(r, t)$  generally develops a space-time profile, from which one can define an information light cone by considering a contour of  $\mathcal{C}(r, t)$ . The contour specifies a function r(t)that describes how fast information propagates. The information propagation largely depends on how the interaction decays in space. An interesting case arises when  $J_{rr'}$  decays algebraically as a function of |r - r'|,  $J_{rr'} \sim 1/|r - r'|^{\alpha}$ . One can show that, as the interaction becomes more and more short ranged, C(r, t) deviates from the fast scrambling behavior [45]. As  $\alpha$  increases, the asymptotic form of r(t) undergoes a series of transitions from exponential  $r \sim$  $\exp(t^{\eta})$  ( $\eta > 0$ ) to algebraic  $r \sim t^{\xi}$  ( $\xi > 1$ ) and eventually, when  $\alpha > 1.5$ , becomes linear  $t \sim r$  [46,47], indicating that information transport slows down from superballistic to ballistic. When  $\alpha \to \infty$ , the interaction becomes short ranged and the usual Lieb-Robinson bound in Eq. (15) applies, restricting r(t) to be at most ballistic [48–50].

Another line of research investigates how the presence of quenched disorder and localization affects the information light cone [51-60]. Based on a phenomenological model called the weak link model [57], it was shown that starting from a model with short-ranged interaction, increasing the disorder strength impedes the information propagation and drives a series of transitions of the light-cone function from linear  $r \sim t$  to algebraic  $r \sim t^{\xi}$  ( $\xi < 1$ ) and eventually becomes logarithmic  $r \sim \log t$  when the system becomes many-body localized. Similar transitions are found in quasiperiodic systems as well [61,62].

As this high-level overview makes explicit, information scrambling in quantum many-body systems displays a rich set of behaviors dependent on both connectivity and disorder summarized in Fig. 6. In the next section, we focus on the case with short-ranged interaction where the information light cone is linear and discuss the behavior of OTOCs in more detail. Furthermore, we explicitly show how OTOCs are calculated in a brickwork random quantum circuit.

# C. Scrambling dynamics in systems with short-ranged interaction

For a generic system with short-ranged interaction and no disorder, local operators spread ballistically, which has been shown in numerous systems including field theories [63,64], free and integrable models [65,66], interacting spin chains [67,68], and circuit models [69–72]. In this case, we have  $r(t) \sim v_B t$ , where  $v_B$  is called the butterfly velocity. The typical behavior of a ballistically expanding C(r, t) is sketched in Fig. 7. Fixing r and varying t, C(r, t)grows from 0 to the saturation value, revealing how the operator becomes complicated locally once the operator front reaches point r. On the other hand, fixing t and varying r, C(r, t) decays from the saturation value to 0, as r exits the light cone. These plots are often used in the literature to characterize the behavior of C(r, t).

In the local system, the Lieb-Robinson bound, already mentioned in Sec. II imposes substantial restrictions on the form of C(r, t). Because of the Lieb-Robinson bound,

$$C(r,t) \le \| [W(t),\sigma_r] \|_{\infty} \le c e^{\lambda_{\mathrm{LR}}(t-r/v_{\mathrm{LR}})}.$$
 (90)

Recall that  $\lambda_{LR}$  is a state-independent loose bound on the behavior of C(r, t). From the bound, it is natural to guess that

$$C(r,t) \sim \exp(\lambda(t - r/v_B))$$
(91)

when *C* is far from saturation. This is indeed the correct form of many models in the large-*N* or semiclassical limit based on field theory calculations [63,73]. However, in the random unitary circuit model [69,70], the tail of the OTOC has a different behavior,

$$C(r,t) \sim \exp\left(-\frac{(x-v_B t)^2}{4Dt}\right).$$
 (92)

In contrast to large-N and semiclassical calculations, this ballistically expanding wave has a diffusively broadened

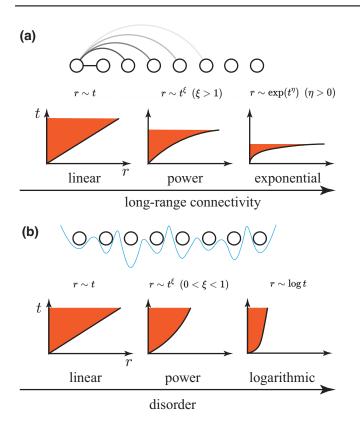


FIG. 6. Control information propagation by long-range coupling and disorder. (a) As the long-range connectivity of a system increases, the shape of the information light cone changes from a ballistic light cone to a algebraic light cone, and finally to an exponential one. (b) Increasing the strength of disorder in a system causes a slowdown in information propagation. As the disorder strength increases, the shape of the light cone changes from ballistic to algebraic and finally to logarithmic.

wavefront, meaning that the scale over which *C* varies as a function of  $u = r - t/v_B$  goes like  $\sqrt{D_B t}$ . Note that the random circuit model also has a version with a large number  $N_{dof}$  on each site, but while  $v_B$  and  $D_B$  depend on this number, the holographic form is never obtained. Furthermore, the Lieb-Robinson bound also applies to the noninteracting system, where the squared commutator can be calculated exactly. In this case, we obtain another different behavior

$$C(r,t) \sim \exp\left(\lambda \frac{(t-r/v_B)^{3/2}}{t^{1/2}}\right).$$
 (93)

The typical behavior of each of the three classes of OTOC is illustrated in Fig. 8. One should not expect the noninteracting limit to be generic, but the spectrum of multiple different universality classes allowed by the Lieb-Robinson bound is certainly raised. There seem many possibilities.

To understand the generic behavior of the OTOC, let us look at the constraint on the functional form of the OTOC imposed by the Lieb-Robinson bound. The Lieb-Robinson bound implies that

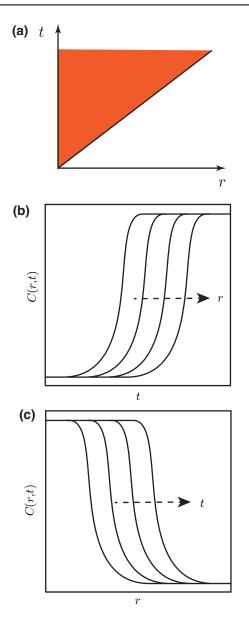


FIG. 7. Schematic sketch of the squared commutator C(r, t) in local clean system. (a) Illustrative contour plot demonstrates ballistic expansion of C(r, t). (b) Fixing r, C(r, t) increases as a function of t. (c) Fixing t, C(r, t) decreases as a function of r.

(i) C(r, t) decays exponentially or faster with r, fixing t;
(ii) C(r, t) grows exponentially or slower with t, fixing r;
(iii) C(vt, t) decays exponentially or faster with t for v > v<sub>LR</sub>.

These constraints are most restrictive outside the light cone where C(r, t) is still small. A general form [68,74] that satisfies these constraints is

$$C(r,t) = \exp\left(\frac{\lambda(t-r/v_B)^{1+p}}{t^p}\right).$$
 (94)

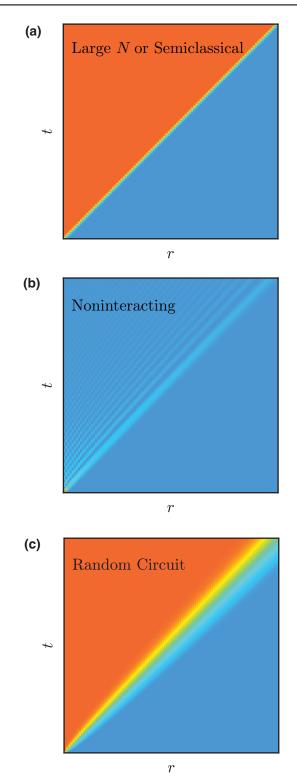


FIG. 8. Schematic illustration of the squared commutator with a ballistic propagating front in quantum many-body systems. (a) In large-N or semiclassical systems, the squared commutator C(r, t) increases rapidly after the wavefront passes, featuring a sharp wavefront. (b) In noninteracting systems, the squared commutator exhibits a drastically different behavior, prominently peaked near the wavefront. (c) In random circuits, C(r, t) exhibits similarities to the large-N or semiclassical case, but the wavefront broadens diffusively.

This growth form unifies the three classes mentioned above into a single framework by including one additional parameter p, called the broadening exponent. The large-N and semiclassical result fits the form with p = 0 (no broadening). The random circuit result fits the form with p = 1 in d = 1 (diffusive broadening). The noninteracting fermion result fits with p = 1/2 in d = 1.

The physics of the general growth form and the broadening exponent p are as follows. Given the general shape in Eq. (94), the contours  $C(r, t) = \theta$  obey

$$r_{\theta} = v_{B}t + \left(\frac{t^{p}}{\lambda}\log\frac{1}{\theta}\right)^{1/(1+p)}.$$
(95)

Hence, no matter the value of  $\theta$ , asymptotically, one has

$$\lim_{t \to \infty} \frac{r_{\theta}}{t} = v_B. \tag{96}$$

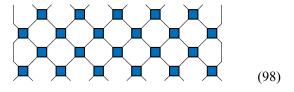
However, at any finite *t*, the contour has an extra subballistic time dependence going like  $t^{p/(p+1)}$ , which is due to the wavefront broadening. As a result, the spatial distance between two contours at a given *t* is

$$\delta r = r_{\theta_1} - r_{\theta_2} \sim t^{p/(1+p)}. \tag{97}$$

This is the key difference between the large-N or semiclassical models and the other models such as noninteracting systems and random circuit models. In the large-N or semiclassical models,  $\delta r$  does not grow with t as  $t \rightarrow \infty$ . Although the noninteracting system is not scrambling and should not be expected to be generic, the difference between the large-N or semiclassical models and random circuit models, both strongly interacting and scrambling, requires understanding.

### D. Random circuit model

We now provide a concrete calculation of the OTOC in a random quantum circuit, a prototypical many-body model for studying entanglement generation [75] and operator dynamics [69,70]. See also the review by Fisher *et al.* [76]. The random circuit contains alternating even and odd layers of random two-qubit gates



with each block  $\square$  an independent Haar random unitary of dimension  $4 \times 4$  for qubits or  $q^2 \times q^2$  in general for local dimension q. There are also generalizations of the random unitary circuit to respect charge conservation

[77,78], dipolar conservation [79], or other kinetic constraints, which are important for studying the interplay between transport phenomena and scrambling. In these random circuit models, the random average of the local unitary operator usually maps the quantum many-body models to statistical models that are easier to handle, while retaining the universal aspects of the quantum many-body dynamics.

Let us focus on the random circuit model without any structure except for the brickwork structure. The idea to calculate the OTOC is to track the Haar-averaged time evolution of the operator probability distribution,

$$|\alpha(S)|^{2} = \frac{1}{q^{2N}} |\text{tr}(S^{\dagger}W(t))|^{2},$$
(99)

where S is an operator string defined in Appendix B 1 and *t* is discrete.

In the random unitary circuit, each Haar random unitary can be averaged out independently. Consider applying a single local unitary gate u to site r and r + 1 to operator W. The updated operator probability is given by

$$|\alpha(\mathcal{S})|^2 = \frac{1}{q^{2N}} \mathbb{E} |\mathrm{tr}(\mathcal{S}^{\dagger} U^{\dagger} W U)|^2.$$
(100)

It is instructive to consider only two sites first. It is straightforward to show that the Haar average leads to

$$|\alpha'(\mathcal{S})|^{2} = \sum_{\mathcal{S}'} \left[ \delta_{I,S} \delta_{I,S'} + \frac{1}{q^{4} - 1} (1 - \delta_{I,S}) (1 - \delta_{I,S'}) \right] \times |\alpha(\mathcal{S}')|^{2}.$$
(101)

Note the remarkable feature that the updated operator probability depends only on the operator probability before the update, but not the amplitude. This simplification is due to the Haar random local dynamics, but it is not true in general many-body systems. In addition, the identity operator stays the identity operator, as expected from the unitary time evolution. The nonidentity operators because of the scrambling nature of the Haar random unitary. To proceed, each operator string is mapped to a binary string. On each site, the identity operator is mapped to 0 and the others are mapped to 1. The probability of the binary string P(S) is the sum of the probability of operators mapping to the same string. Then the transition rules are given by

$$P(00, t+1) = P(00, t),$$
(102a)

$$P(01, t+1) = \frac{1}{q^2 + 1} (P(11, t) + P(10, t) + P(01, t)),$$
(102b)

$$P(10, t+1) = \frac{1}{q^2 + 1} (P(11, t) + P(10, t) + P(01, t)),$$
(102c)

$$P(11, t+1) = \frac{q^2 - 1}{q^2 + 1} (P(11, t) + P(10, t) + P(01, t)).$$
(102d)

Since, in the Haar random circuit, each unitary is independent, the transition rules above also hold locally when a local unitary u acts on sites r and r + 1. In this case, the binary string probabilities only change locally on sites r and r + 1 according to the transition rules given above. Thus, the unitary dynamics after random average becomes stochastic dynamics described by a master equation.

To gain an analytical handle on the effective stochastic dynamics, it is useful to consider the probability that the last 1 in the binary string ends at r,  $P_{end}(r)$ . Based on Eqs. (102), applying the local unitary  $u_{rr+1}$  updates  $P_{end}$  as

$$P'_{\rm end}(r) = \frac{1}{q^2 + 1} (P_{\rm end}(r, t) + P_{\rm end}(r+1)), \qquad (103a)$$

$$P'_{\text{end}}(r+1) = \frac{q^2}{q^2+1}(P_{\text{end}}(r,t) + P_{\text{end}}(r+1)).$$
 (103b)

Note that  $(P_{end}(r, t) + P_{end}(r + 1))$  is conserved, as expected. The unitary  $u_{r,r+1}$  redistributes them so that the operator has a higher probability of ending at r + 1, leading to expansion. Recall that the unitary circuit acts on the even and odd bonds alternatively. The layer of even (odd) t acts on even (odd) bonds. To take into account the combined effect of even and odd layers of the unitary circuit, it is more convenient to track the sum of  $P_{end}$  on even bonds, denoted  $P_{end}(b)$ , after each even layer. Because of Eqs. (103),  $P_{end}(b)$  for even b fully specifies  $P_{end}(r)$  after the even unitary layer. One can show that  $P_{end}$  obeys

$$P_{\text{end}}(b, t+2) = \left(\frac{q^2}{q^2+1}\right)^2 P_{\text{end}}(b+2, t) + \left(\frac{1}{q^2+1}\right)^2 P_{\text{end}}(b-2, t) + \frac{2q^2}{(q^2+1)^2} P_{\text{end}}(b, t).$$
(104)

This equation describes a biased random walk. In the continuum limit, the equation becomes a biased diffusion equation

$$\partial_t P_{\text{end}}(r,t) = v_B \partial_r P_{\text{end}}(r,t) + D \partial_r^2 P_{\text{end}}(r,t),$$
 (105)

where

$$v_B = \frac{q^2 - 1}{q^2 + 1}, \qquad D = \frac{2q^2}{(q^2 + 1)^2}.$$
 (106)

Note that in the circuit model, an upper bound of velocity is 1, which is set by the geometry of the circuit and the naive light cone of the Heisenberg operator shown in Eq. (13). This upper bound can be regarded as the Lieb-Robinson velocity of the circuit model. Here the butterfly velocity obtained is *smaller* than the upper bound. For q = 2,  $v_B = 3/5$ . This is because of the return probability in the biased diffusion equation. As  $q \rightarrow \infty$ , the butterfly velocity approaches 1 and *D* decreases to 0.

The solution of the biased diffusion equation with initial condition  $\delta(r)$  is

$$P_{\rm end}(r,t) = \frac{1}{\sqrt{4\pi Dt}} \exp\left(-\frac{(r-v_B t)^2}{4Dt}\right).$$
 (107)

To obtain the squared commutator, it is reasonable to assume that the operator on every site behind the endpoint of the operator is in local equilibrium. As a result, from Eq. (82b),

$$C(r,t) = \frac{2q^2}{q^2 - 1} \sum_{S_r \neq I} |\alpha(S)|^2 \sim 2 \int_r^\infty P_{\text{end}}(r',t) dr'$$
  
= 1 + erf $\left(-\frac{(r - v_B t)}{\sqrt{4Dt}}\right).$  (108)

It exhibits ballistic expansion and diffusive broadening of the wavefront. When  $t > r/v_B$ , it quickly saturates to the final 2, indicating scrambling. The tail behavior of C(r, t)can be obtained by expanding the error function in the limit that  $r - v_B t \gg \sqrt{Dt}$ . This leads to the growth form given in Eq. (92), which is in contrast with the growth form obtained in the semiclassical or large-*N* and AdS/CFT models. Physically, the Gaussian tail of the squared commutator obtained in the random circuit model relies on two factors. First, the endpoint of the operator undergoes a random walk biased towards expansion. Second, the operators behind the endpoint immediately reach the local equilibrium because of the Haar random unitary.

#### **VI. NUMERICAL METHODS**

In this section, we discuss some existing numerical methods to calculate the OTOC in many-body systems, including commenting on their applicability and limitations. The numerical methods can be roughly divided into two categories, exact diagonalization and tensor networks. There are many other wholly or partially numerical approaches to calculating OTOCs, such as the truncated Wigner approximation in the semiclassical limit [80,81], but these are among the most general purpose. To be concrete, we consider the following prototypical spin chain Hamiltonian called the mixed-field Ising model with nearest-neighboring Ising interaction,

$$H = \sum_{i} (J\sigma_{i}^{z}\sigma_{i+1}^{z} + H_{x}\sigma_{i}^{x} + H_{z}\sigma_{i}^{z}).$$
(109)

We suggest interested readers try the different methods introduced below on this Hamiltonian.

#### A. Exact diagonalization and the Krylov space method

We first discuss the exact diagonalization—based method to compute the OTOC in the system described by Hamiltonian *H*. In this case, it is more convenient to consider the OTOC in the form of F = tr(W(t)VW(t)V), where W(t)is the time-evolved Heisenberg operator and *V* is a local probe operator at site *r*. The most straightforward method to compute OTOC is to perform a full exact diagonalization on *H* to obtain the eigenvectors  $|n\rangle$  as well as the eigenvalues  $\epsilon_n$ . Then the OTOC can be evaluated as

$$F(t) = \frac{1}{2^N} \sum_{m,n,p,q} e^{i\epsilon_{nm}t} W_{nm} V_{ml} e^{i\epsilon_{pq}t} W_{pq} V_{qn}, \qquad (110)$$

where  $\epsilon_{mn} = \epsilon_m - \epsilon_n$ . The matrices *W* and *V* are in the eigenstate basis. They only need to be computed once to calculate the OTOC at different times. Although this is the simplest and numerically exact method, it suffers severely from limited scalability. The bottleneck is the full diagonalization of *H* and storing the eigenstates, which can be done for up to 15 qubits—it takes about 40G memory to store the whole set of eigenstates.

One can avoid exact diagonalization by implementing the Heisenberg time evolution using the Krylov method [82]. The Krylov basis can be generated by applying the Hamiltonian to the operator iteratively k times, where k + 1 sets the effective Hilbert space dimension. Specifically, we have

$$W^{(n)} = HW^{(n-1)} - W^{(n-1)}H, \qquad 1 \le n \le k.$$
(111)

The Krylov method is also numerically exact. It does not require diagonalizing the Hamiltonian, but it requires regenerating the basis by applying the Hamiltonian to the operator multiple times at each step, which can be time consuming. This method is limited by storing the large matrix of W(t). Initially, W is a local operator in the form of a sparse matrix. The complexity of W(t) increases with time, and in the late time, W(t) becomes dense and requires a huge memory to store. As a result, this operator-Krylov method, similar to the naive exact diagonalization method, can also only work for up to 15 spins.

The difficulty of storing the full Heisenberg operator can be circumvented by calculating the OTOC in the Schrödinger equation and using the typicality of random states. Using the property of Haar random state  $|\psi\rangle$ , we know that

$$\mathbb{E}_{\psi}(\langle \psi | O | \psi \rangle) = \frac{1}{2^{N}} \text{tr} O.$$
(112)

Therefore, we can replace the trace in the OTOC by sampling over Haar random states,

$$F = \mathbb{E}_{\psi} \langle \psi | W(t) V W(t) V | \psi \rangle$$
  
=  $\mathbb{E}_{\psi} \langle \psi | e^{iHt} W e^{-iHt} V e^{iHt} W e^{-iHt} V | \psi \rangle.$  (113)

Based on quantum typicality, the sample-to-sample fluctuation is suppressed by the large Hilbert space. As a result, one random state is sufficient to capture the OTOC. See also Ref. [83]. One can evolve  $V|\psi\rangle$  and  $|\psi\rangle$  back and forth in time efficiently using the Krylov method. We dub this method the "Krylov-State" method. Figure 9(a) shows the OTOC from this method and the exact diagonalization based on Eq. (110), which agree well with each other. The main bottleneck of this method is to store the entire quantum state, a dense vector in the Hilbert space instead of a dense matrix as that in the previous methods. As a result, the maximal applicable system size is doubled, i.e., about 30 qubits [84,85].

All three methods introduced in this section are numerically exact for calculating the OTOC. They apply to arbitrarily long time, but are limited to quite a small system size. Among the three, the Krylov-State method applies to systems with up to 30 qubits and is thus significantly better than the Krylov-Operator method and the most straightforward the exact diagonalization (ED) method.

#### **B.** Tensor-network methods

In this section, we discuss another class of methods that have been used in the literature to calculate the OTOC, which is based on time-evolving matrix product states (MPSs) and matrix product operators (MPOs) [67,68,86-88] and Sec. 8.4 of Ref. [89]. Unlike the exact diagonalization and Krylov methods, these methods can be used to study huge systems containing a few hundred spins; the numerical cost increases linearly with L. The main bottleneck here is the time scale. These methods always produce a result for any time, but the result becomes more and more inaccurate as the state and operator are evolved longer and longer. In the following, we summarize the key concept of these methods and discuss their application to calculate the OTOC, assuming that readers have a basic familiarity of the density matrix renormalization group and time-evolution block decimation [90].

The central object in these methods is the matrix product state. In the MPS, the wave-function amplitude is a product of a series of matrices defined for each spin,

$$\langle \{s\} | \psi \rangle = A_1^{s_1} A_2^{s_2} \cdots A_L^{s_L}.$$
(114)

Fixing s to  $\uparrow$  or  $\downarrow$ ,  $A_i^s$  is a matrix associated with the *i*th spin, and s is the physical index. The dimension of the matrix  $\chi$  is called the bond dimension. On the first site and the last site,  $A^s$  is a vector so that the product

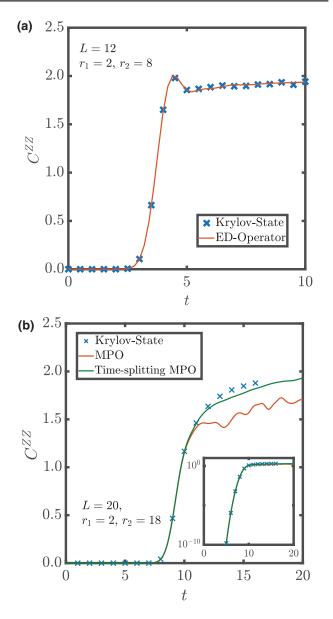


FIG. 9. Comparison of the squared commutator between  $\sigma_{r_1}^z$ and  $\sigma_{r_2}^z$  calculated using different methods. (a) Comparison between the Heisenberg time evolution based on exact diagonalization (ED-Operator) and the Schrodinger time evolution based on the Krylov method and quantum typicality (Krylov-State). Only one random initial state is used in the Krylov-State method. The squared commutators from the two methods agree well with each other. (b) Comparison between the Krylov-State method and MPO methods. The Krylov-State method is numerically exact. The result from the MPO method, which involves truncation errors, agrees perfectly with the exact result up to  $C^{zz} \sim 1$ . The accuracy of the MPO results can be enhanced by time splitting. The inset shows the same figure on a log-linear scale, demonstrating that the three methods agree when C is small. The maximal bond dimension of both MPO methods is set to 32.

results in a number, the wave function amplitude of the spin configuration  $\{s\}$ . The tensor network representation

of an MPS is



Analogously, one can also write down the matrix product form of an operator, called MPO, in which case, matrix *A* has two physical indices. The tensor network representation of an MPO is similar to that of an MPS:



The MPS requires storing 2L matrices, the 2 arising from the physical indices  $\uparrow$  and  $\downarrow$ , in total  $2L\chi^2$  complex numbers. Therefore, an MPS is a very compact approach to represent a wave function that contains  $2^L$  numbers. As expected, an MPS of a small bond dimension  $\chi$  can only represent a tiny portion of the Hilbert space, characterized by low entanglement entropy. The entanglement entropy of an MPS with bond dimension  $\chi$  is maximal log  $\chi$ , a constant, while the entanglement entropy of a random state in the Hilbert space scales linearly with *L*. Fortunately, the entanglement entropy of the ground state of one-dimensional gapped systems exhibits area law, i.e., does not scale with system size, and can thus be captured accurately.

A large toolbox for manipulating the MPS or MPO has been developed over the years [89,90]. The key idea behind studying dynamics using the MPS or MPO is to evolve the state and operator while preserving the matrix product form. However, unlike searching for the ground state, the complexities of the simple initial states and initial operators increase over time, and in late time, the states exhibit volume-law entanglement entropy. As a result, at some point during the evolution, usually a short time scaling as one over the interaction strength, the state cannot be captured faithfully by an MPS with a reasonable bond dimension. In practice, the quantities of physical interest are local observables and correlation functions, instead of the entire quantum state. To obtain the physical observable  $\langle O(t) \rangle$  at a given time, one can repeat the simulation with increasing bond dimension, and the result can be trusted if it converges with the bond dimension.

There are two major approaches to evolving an MPS or MPO: time-evolving block decimation (TEBD) and the time-dependent variational principle (TDVP). In TEBD, the local unitary gate is directly applied to the MPS and increases the bond dimension of the MPS. Then the MPS is truncated so that the bond dimension stays the same, leading to the truncation error. Furthermore, the truncation error in TEBD does not necessarily respect conserved quantities, such as the energy and charge, which start deviating from their initial values after a short time scale. This issue can be fixed by using the TDVP method to evolve the state instead. In short, TDVP generates an effective Hamiltonian for the local tensor in the MPS, which depends on other tensors. The effective Hamiltonian is used to evolve the tensor unitarily. As a result, the total energy is conserved by construction during the time evolution, and TDVP can take into account other conserved quantities. Although respecting conservation laws does not necessarily ensure correct local observable, correlation function or transport behavior, this is the first step towards approaching the correct nonequilibrium state beyond a short time. This is particularly important in the late-time regime where hydrodynamics emerges and conserved quantities play a crucial role.

#### 1. MPS-based method

Now, we discuss applying these methods to calculate the OTOC. Like methods using exact diagonalization, one can also calculate the OTOC in either the Schrödinger or Heisenberg picture. In Schrödinger's picture, similar to the exact diagonalization discussed before, the trace in the OTOC F(t) is replaced by the average over an ensemble of states. However, the ensemble of random states is not feasible in tensor network methods because the random states have large entanglement that the MPS cannot capture. Instead, in practice, an ensemble of random product states is used. Fix the initial state as  $\psi$ . Either TEBD or TDVP can be used to generate the two states

$$|\psi_1\rangle = e^{iHt} W e^{-iHt} V |\psi\rangle, \qquad |\psi_2\rangle = V e^{iHt} W e^{-iHt} |\psi\rangle.$$
(117)

Then the OTOC is just the overlap between the two states followed by averaging over the initial states. The two methods, dubbed the TEBD-MPS and TDVP-MPS methods, respectively, which differ by the time evolution method, are compared in detail in ref. [86] and benchmarked with the numerically exact Krylov-State method. As expected, the TDVP-MPS method produces more accurate results as time passes a short time scale. Since the numerical costs of the TDVP-MPS and TEBD-MPS methods are comparable, TDVP is the better approach to calculate the OTOC in the Schrödinger picture. In the state-based method, the primary source of error is the incapability of representing the complicated state using an MPS with a finite bond dimension.

#### 2. MPO-based method

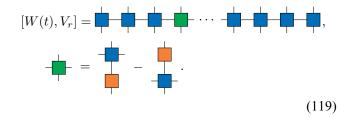
During the time evolution of a state, the entanglement of the states increases uniformly across the whole system. On the other hand, because of the emergent light-cone structure in the Heisenberg evolution of a local operator shown in Eq. (13), the Heisenberg operator W(t) is almost a product of identity operators  $I \otimes \cdots \otimes I$  outside the light cone and becomes complicated inside the light cone. For this reason, the entanglement entropy of an operator (see Appendix A 2 b for a brief introduction on the operator entanglement entropy), which can be regarded as a state with doubled physical degrees of freedom, develops a light-cone profile as well, staying small outside the light cone and becoming large within the light cone regardless of the time scale. Therefore, the part of the Heisenberg operator outside the light cone can be captured accurately using an MPO with a small bond dimension. This intuition leads to the method of calculating the OTOC in the Heisenberg picture using an MPO. In the MPO-based method, the main step is to evolve the Heisenberg operator, which is similar to evolving the MPS by treating the MPO as an MPS with doubled physical degrees of freedom. The time evolution operator to evolve the operator state is

$$W(t) = UWU^{\dagger} \to U \otimes U^* |W\rangle.$$
(118)

This can also be viewed as evolving the operator state  $|W\rangle$  using the "super-Hamiltonian"  $H \otimes I - I \otimes H^*$ . The time evolution can be implemented using either TEBD or TDVP. The OTOC can be calculated once the MPO form of the Heisenberg operator W(t) is obtained. In this approach, it is more convenient to calculate the OTOC in the form of a squared commutator

$$C(r,t) = \frac{1}{2^{L}} \text{tr}([W(t), V_{r}]^{\dagger} [W(t), V_{r}]),$$

where  $V_r$ , a local operator on site r, is scanned over the system. Given the MPO of W(t), the MPO of the commutator  $[W(t), V_r]$  is obtained efficiently by modifying the local tensor at site r as



The bond dimension remains the same because  $V_r$  is a local operator represented by the orange tensor with only physical indices. The squared commutator can be obtained by contracting the MPO of the commutator with its Hermitian conjugate copy. However, noting that C(r, t) is just the square of the Frobenius norm of the commutator, this last step can be replaced by evaluating the norm of the MPO directly by treating it as an MPS. Evaluating the norm of an MPS is a standard procedure in the tensor network simulation. Calculating the norm then squaring it is much more accurate than calculating the overlap between the MPO and its Hermitian conjugate copy, especially when C(r, t) is small. Figure 9(b) shows the squared commutators from the MPO method and the Krylov-State method for 20 qubits. The MPO results are accurate when the squared commutator is small, but starts to deviate from the exact Krylov-State result from the Krylov-State method as the squared commutator reaches about 1.

As shown in Ref. [68], because of the light-cone structure of the Heisenberg operator, the MPO method is very accurate in capturing the tail of the OTOC even with a bond dimension as small as 4, and can be used to obtain the butterfly velocity. This method also clearly demonstrates a broadened wavefront, agreeing with the general growth form for p > 0, in contrast to the behavior in the Sachdev-Ye-Kitaev (SYK) chain. Obtaining the asymptotic value of p is quite subtle and requires accessing a wider spacetime region and larger bond dimension. The TEBD-MPO method was also compared to the TDVP-MPS method by Hémery et al. [86], who showed that both methods are accurate in capturing the tail of the OTOC and tracking the light cone. The TEBD-MPO method has two advantages over the TDVP-MPS method. First, it does not require sampling over the initial states and is thus free of statistical error. Second, at given time t, the commutator of W(t) and  $V_r$  and therefore C(r) for all r can be obtained efficiently using Eq. (119). On the other hand, in the TDVP-MPS method, in order to get C(r, t) for all sites r, initial states  $V_r |\psi\rangle$  with different r need to be evolved individually in addition to sample  $|\psi\rangle$ .

One can also evolve the Heisenberg operator using the TDVP method based on the super-Hamiltonian  $H \otimes I - I \otimes H^*$ . The quantity made explicitly conserved in the TDVP-MPO method, instead of the energy, is the expectation value of the super-Hamiltonian

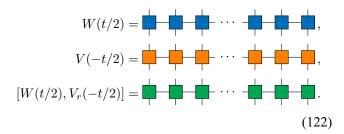
$$\langle W | \mathbb{H} | W \rangle = \operatorname{tr}(W^{\dagger}(t)HW(t)) - \operatorname{tr}(W(t)HW^{\dagger}(t)), \quad (120)$$

which is just zero for a Hermitian operator W(t). In practice, the TDVP-MPO method does not significantly increase the accuracy of the TEBD-MPO method. It is less accurate to capture the exponentially small value of the OTOC because it does not directly take advantage of the light-cone structure. However, the TDVP-MPO method is still a very useful and convenient method to use when the system has long-range interaction [47], in which case the circuit from Trotterization does not admit the structure shown in Fig. 2. It is also possible to generalize the TDVP-MPO method to conserve energy or/and charge, which remains to be explored.

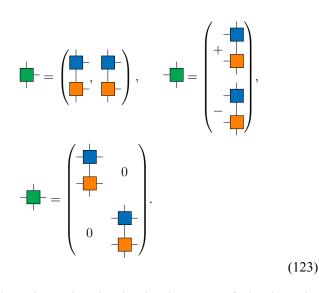
A few remarks are in order. First, the MPO-based method utilizes the light-cone profile of the entanglement structure of the Heisenberg operator. It becomes exact if the operator entanglement entropy is bounded, such as in the noninteracting system. Second, in a chaotic system, because of the rapid generation of operator entanglement entropy within the light cone, the tensor network based

$$\frac{1}{2^{L}} \operatorname{tr}([W(t/2), V_{r}(-t/2)]^{\dagger}[W(t/2), V_{r}(-t/2)]).$$
(121)

expression is equivalent, [W(t/2),This since  $V_r(-t/2)$ ] =  $e^{-iHt/2}$  [W(t),  $V_r$ ] $e^{iHt/2}$ . This time-splitting method divides the entanglement growth within the light cone into two operators and reduces the truncation error. The cost is that each  $V_r$  needs to be evolved individually in order to obtain the full space-time profile of C(r, t). In addition, the fast construction of the commutator in Eq. (119) is no longer applicable, since now the commutator is between two time-evolved Heisenberg operators, both taking MPO form with bond dimension  $\chi$ . As a result, the MPOs of W(t/2)V(-t/2) and W(t/2)V(-t/2) have bond dimension  $\chi^2$ , and the MPO of their difference, the commutator, has bond dimension  $2\chi^2$ . Graphically,



Fixing the physical indices to be  $\alpha$  and  $\beta$ , the local tensor of the MPO of the commutator is constructed from the local tensor of the two Heisenberg operators



The minus sign in the local tensor of the last site implements the difference between the two terms in the

commutator. Once the MPO of the commutator is constructed, its Frobenius norm and thus the OTOC C(r, t)can be calculated in the same way as before. As shown in Fig. 9(b), the time-splitting approach significantly enhances the accuracy of the MPO in the late-time regime. Finally, in all the tensor network–based methods applied to a system with about 100 spins, because of the targeted intermediate time scaling of about 100/*J*, where *J* is the coupling constant, one should carefully check the convergence of the result with the Trotter time step. In general, the time step *dt* should be reduced to about 0.005/*J* to avoid accumulating the Trotter error.

### VII. EXPERIMENTAL SCHEMES

As we enter an era of quantum simulation, marked by the existence of multiple experimental platforms with unprecedented power to control and detect quantum manybody physics far from equilibrium, there is a surge of interest in measuring scrambling dynamics and OTOCs. The essential step required to measure an OTOC is to write it as an observable or a combination of observables, including specifying the initial states, the time evolution operators, and the measurements. In this section, we briefly discuss several such experimental schemes, including those already implemented in the laboratory. This discussion focuses on the essence of several OTOC measurement protocols, and it does not delve into the experimental details of various concrete experiments. These schemes fall into two main categories, one requiring the rewinding time and the other requiring measurements averaging over random states.

### A. Rewinding time

The experimental schemes falling into this category require engineering both the forward evolution operator U(t) and the backward time evolution operator U(-t), and is closely related to the Loschmidt echo [91,92]. The first scheme that we describe is based on an interferometric protocol [93]. This scheme requires introducing a reference qubit that is initialized in state  $|+\rangle = (|0\rangle + |1\rangle)/\sqrt{2}$ . The initial state of the system and the reference qubit is

$$|\Psi_0\rangle = |+\rangle |\psi\rangle. \tag{124}$$

To measure OTOC  $\langle W^{\dagger}(t) V_r^x W(t) V_r^x \rangle$ , one first applies the controlled-*V* gate to the state (assuming that *V* is unitary), acting on site *r* and controlled by the reference qubit, then evolves the state by a butterfly unitary circuit  $U_{\text{butterfly}}$ , and applies the controlled-NOT (CNOT) gate. The butterfly unitary circuit  $U_{\text{butterfly}}$  does not act on the ancillary qubit and is of the form

$$U_{\text{butterfly}} = U^{\dagger}(t)WU(t) = W(t), \qquad (125)$$

where U(t) is an arbitrary unitary circuit and W is a local unitary gate. The butterfly circuit is just another name for the Heisenberg operator of a local unitary operator. After these steps, the state is

$$\begin{aligned} |\Psi\rangle &= \frac{1}{\sqrt{2}} (|0\rangle \, U_{\text{butterfly}} \, |\psi\rangle + |1\rangle \, \sigma_r^x U_{\text{butterfly}} \sigma_r^x \, |\psi\rangle) \\ &= \frac{1}{\sqrt{2}} (|0\rangle \, W(t) \, |\psi\rangle + |1\rangle \, \sigma_r^x W(t) \sigma_r^x \, |\psi\rangle). \end{aligned} \tag{126}$$

The last step is to measure  $\sigma^x$  on the reference spin. The expectation value, which can be obtained by running the circuit and measuring repetitively, is

$$\langle \Psi | \sigma_R^x | \Psi \rangle = \operatorname{Re} \langle \psi | W^{\dagger}(t) V_r^{\dagger} W(t) V_r | \psi \rangle \qquad (127)$$

in the form of an OTOC. The imaginary part of the OTOC can be obtained by measuring  $\sigma^{y}$  instead of  $\sigma^{x}$ . This protocol works for an arbitrary initial state of the system, but requires introducing the reference spin and the CNOT gate and is used in Ref. [94].

For certain initial states, introducing the reference qubit is not necessary. Consider an initial state that is an eigenstate of a local operator, for instance,  $\sigma_r^x$ . This means that the local qubit of the state is either  $|+\rangle$  or  $|-\rangle$ . The idea is to measure whether the qubit at site *r* remains an eigenstate of  $\sigma_r^x$  after evolving by the butterfly circuit. This is quantified by its expectation value, given by

$$\langle \psi(t) | \sigma_r^x | \psi(t) \rangle = \langle \psi | W^{\dagger}(t) \sigma_r^x W(t) | \psi \rangle.$$
(128)

If the support of W(t) does not overlap site *r* then the local qubit remains an eigenstate; otherwise, the expectation value decays. Formally, one can write the above equation in the form OTOC  $\langle \psi | W^{\dagger}(t)\sigma_r^x W(t)\sigma_r^x | \psi \rangle$ , using the fact that state  $|\psi\rangle$  is an eigenstate of  $\sigma_r^x$ . This scheme is carried out in Refs. [95–97]

In the previous example, the OTOC is measured with respect to a pure state. To directly measure the OTOC for the infinite-temperature ensemble, the most straightforward setup is to prepare double copies of the system, such as a ladder. The initial state is a product of Bell states across each rung  $|\Psi\rangle = \prod |0\rangle_i |0_{i'}\rangle + |1\rangle_i |1_{i'}\rangle$ . The Hamiltonian of the full system  $\mathbb{H} = H \otimes I - I \otimes H^*$ . Then the OTOC can be cast into an observable in this double copied system

$$\frac{1}{\operatorname{tr} I} tr(W^{\dagger}(t)V^{\dagger}W(t)V) = \langle \Psi | V^{\dagger}e^{i\mathbb{H}t}W \otimes W^{\dagger}e^{i\mathbb{H}t}V | \Psi \rangle.$$
(129)

Replacing the initial state  $|\Psi\rangle$  with the thermal field double state  $|\sqrt{\rho}\rangle$  generalizes this scheme to measure the OTOC at finite temperature for a specific generalization [98]. The double copy setup is also essential to directly realize the many-body teleportation protocol discussed in Sec. IV, where one copy represents the system and the other copy represents the memory qubit owned by Bob. The many-body teleportation protocol has been implemented in ion trap [99] and superconducting qubit [100] experiments.

It is also possible to measure the OTOC at infinite temperature without introducing the second copy in NMR experiments [101–103]. In this setup, the initial state is a high-temperature mixed state in the presence of the magnetic field along the z direction. The initial state is given by

$$\rho_0 = \frac{1}{\text{tr}I} \left( I + \epsilon \sum_r \sigma_r^z \right). \tag{130}$$

The density matrix is then evolved by unitary time evolution in the form of the butterfly circuit, which can be engineered by a sequence of pulses. In the evolved mixed state, the total magnetization is given by

$$\langle \sigma^{z} \rangle = \operatorname{tr}(U_{\rho}(t)\rho_{0}U_{\rho}^{\dagger}(t)\sigma^{z})$$
$$= \frac{\epsilon}{\operatorname{tr}I}\operatorname{tr}\left(W(t)\sum_{r}\sigma_{r}^{z}W^{\dagger}(t)\sum_{r}\sigma_{r}^{z}\right)$$
(131)

in the form of an OTOC. Note that the identity part of the initial mixed state does not contribute to the expectation value. In an NMR experiment, a natural choice of W is  $\exp(i\theta \sum_r \sigma_r^z)$  from a pulse, which is a nonlocal operator. Then result  $\langle \sigma^z \rangle_{\theta}$  is a periodic function in  $\theta$  with period  $2\pi$ , denoted  $\langle \sigma^z \rangle_{\theta}$ . Using the Lehmann representation, one can show that

$$\int_{0}^{2\pi} d\theta \int d\omega \omega^2 S(\theta) e^{-i\omega\theta} \sim \frac{1}{\operatorname{tr} I} \operatorname{tr} ([S^z(t), S^z][S^z(t), S^z]).$$
(132)

This squared commutator approximately counts the number of spins within the support of the local Heisenberg operator  $\sigma^{z}(t)$ . In practice, one can perform a discrete Fourier transformation on a finite number of measurement results with a discrete value of  $\theta$ . We note that it is also possible to measure the OTOC between two local operators using selective pulses for an ensemble of small molecules [101].

#### **B.** Randomized measurement

The key ingredient for the above protocols is the butterfly circuit  $U_{\text{butterfly}}$ , which requires implementing both U(t)and  $U^{\dagger}(t)$ , namely, effectively rewinding the time, and can be challenging for generic U. An experimental protocol that bypasses this requirement is to exploit statistical correlation from randomized measurement [104], which has been implemented in NMR [105] and ion traps [106]. In its simplest form, this protocol considers the product of two expectation values

$$F(\psi) = \langle \psi | W(t) | \psi \rangle \langle \psi | V^{\dagger} W(t) V | \psi \rangle.$$
(133)

When the initial state is averaged over the Haar ensemble, using the Haar random average formula in Eq. (A17b) in Appendix A 4, the above quantity leads to

$$\mathbb{E}_{\psi}F(\psi) = \frac{1}{2^{N}(2^{N}+1)} \operatorname{tr}(W(t))\operatorname{tr}(V^{\dagger}W(t)V) + \frac{1}{2^{N}(2^{N}+1)} \operatorname{tr}(W(t)V^{\dagger}W(t)V). \quad (134)$$

Typically, the first term is  $\mathcal{O}(1)$ , while the second term, which is the important piece, is  $\mathcal{O}(2^{-N})$ . However, when *W* is a traceless operator, the first term vanishes and only the second term, the OTOC, survives. This scheme can also be extended to obtain the leading finite-temperature correction of the OTOC as well. In this protocol, the major challenge is to prepare the Haar random initial states. Alternatively, one can also estimate the OTOC by sampling over the local unitary that acts on each qubit and averaging the initial states in the computational basis in a specific manner. In this approach, the quantity that is being considered is

$$F_{\vec{n}} = \langle \psi_{\vec{n}} | W(t) | \psi_{\vec{n}} \rangle \langle \psi_0 | V^{\dagger} W(t) V | \psi_0 \rangle.$$
(135)

Compared with Eq. (133), the two expectation values are measured from two different states  $|\psi_{\vec{n}}\rangle$  and  $|\psi_0\rangle$ . The state  $|\psi_{\vec{n}}\rangle$  is generated by acting with local unitaries on a state  $|n\rangle$  in the computational basis

$$\begin{aligned} |\psi_{\vec{n}}\rangle &= u_1 \otimes \cdots \otimes u_N |\vec{n}\rangle \\ &= u_1 |n_1\rangle \otimes u_2 |n_2\rangle \otimes \cdots \otimes u_N |n_N\rangle, \end{aligned} \tag{136}$$

where  $u_i$  is drawn from the Haar ensemble or any ensemble producing the same averaged result as the Haar ensemble involving four copies of the state (such an ensemble is called a 2-design) and  $n_i \in \{0, 1\}$ . The state  $|\psi_0\rangle$  is generated from the all-zero state  $|0\rangle$  using the same unitaries. The random average of the unitaries can be performed independently on each qubit using Eqs. (A17). To gain insight into the result, we first average  $F_{\vec{n}}$  over  $u_1$ . We define the following operators acting on the first qubit:

$$O_1 = \langle u_2^{n_2} | \cdots \langle u_N^{n_N} | W(t) | u_2^{n_2} \rangle \cdots | u_N^{n_N} \rangle,$$
  

$$\tilde{O}_1 = \langle u_N^0 | \cdots \langle u_2^0 | V^{\dagger} W(t) V | u_2^0 \rangle \cdots | u_N^0 \rangle.$$
(137)

Averaging over  $u_1$  leads to

$$\mathbb{E}_{u_1}(F_{\vec{n}}) = \frac{2 - \delta_{n_1,0}}{6} \operatorname{tr}(O_1) \operatorname{tr}(\tilde{O}_1) + \frac{2\delta_{n_1,0} - 1}{6} \operatorname{tr}_{n_1}(O_1\tilde{O}_1).$$
(138)

The trick here is to use a weighted sum over  $n_1$  to cancel the first term. One can show that

$$\sum_{n_1} \left( -\frac{1}{2} \right)^{n_1} \mathbb{E}_{u_1} F_{\vec{n}} = \frac{1}{4} \operatorname{tr}(O_1 \tilde{O}_1).$$
(139)

Consecutively averaging over other local unitaries  $u_i$  and summing over  $n_i$  in a similar fashion leads to the OTOC

$$\sum_{\vec{n}} \left( -\frac{1}{2} \right)^{\sum n_i} \mathbb{E}_{u_1 \cdots u_N}(F_{\vec{n}}) = \frac{1}{4^N} \operatorname{tr}(W(t)V^{\dagger}W(t)V).$$
(140)

One can also extend this protocol to directly measure the operator probability distribution defined in Sec. V [107].

# **VIII. EPILOGUE**

In this tutorial, we have covered the definition of quantum scrambling dynamics in generic quantum many-body systems, which is distinct from thermalization. While thermalization concerns the local density matrix, scrambling dynamics manifests in nonlocal degrees of freedom. We have shown that scrambling dynamics can be understood by treating unitary dynamics as a quantum communication protocol. Scrambling, by definition, is characterized by the time-dependent mutual information between a referee qubit, which is initially entangled with one of the qubits, and a partition of the system. It also affects other nonlocal dynamical properties of the system, such as the entanglement entropy and entanglement spectrum [108]. Importantly, we have shown that one does not need full access to the system to recover an initial qubit state prepared in one of the qubits in the many-body system, even when the system is fully scrambled. In the Hayden-Preskill setup, which involves a maximally mixed initial state, the mutual information can be converted to the OTOC. We have demonstrated the important role of OTOCs in quantum information and dynamics and discussed their general behavior in local systems featuring ballistic information dynamics. We have presented several toy models where the OTOC can be calculated exactly, as well as numerical tools to compute OTOCs in generic systems. We have also surveyed recent exciting experimental progress in detecting information dynamics. However, it is not practical to cover all aspects of this large field in this tutorial article. There are several related interesting topics that we did not discuss, and we would like to briefly mention them here.

#### A. Finite temperature

First, this tutorial largely considered scrambling at infinite temperature. Physically, this amounts to saying that Alice and Bob were using a medium at high temperature compared to the scales in the intrinsic Hamiltonian. Mathematically, this means that we considered OTOCs where the expectation was taken in the maximally mixed state. This regime is ideal for understanding the basics of information dynamics since we do not have to deal with any static correlations in the quantum state. However, it is interesting to understand information away from infinite temperature, especially in the context of black hole physics.

Given a quantum state  $\rho$ , one can define an OTOC in this state by  $F_{\rho} = \text{tr}(\rho W(t)VW(t)V)$ . When  $\rho$  is the maximally mixed state, this recovers the simple trace expression we considered for most of this tutorial. However, if Alice and Bob wish to use a medium at noninfinite temperature to convey quantum information then one expects so-called thermal OTOCs to be relevant where  $\rho \propto e^{-\beta H}$  is a thermal equilibrium state. However, the physics of scrambling is more complicated in this case. This is because there are multiple versions of  $F_{\rho}$  called "regulated" OTOCs,

$$\tilde{F}_{\rho} = \text{tr}(\rho^{q_1} W(t) \rho^{q_2} V \rho^{q_3} W(t) \rho^{q_4} V), \qquad (141)$$

where  $q_i \in [0, 1]$  and  $\sum_i q_i = 1$ . This class of objects corresponds to displacing the operators W(t) and V in imaginary time (when  $\rho \propto e^{-\beta H}$  is the thermal state).

The case  $q_1 = 1$ ,  $q_{i>1} = 0$  is the usual OTOC in state  $\rho$ . The case  $q_i = 1/4$  is a particularly natural choice from a mathematical point of view; this is the form of the OTOC to which the chaos bound [11] applies. In early examples, different regularizations gave equivalent characterizations of the scrambling dynamics. However, it was later discovered that the butterfly velocity can depend on the choice of regularization [109–111]. This raises the question of which regularization is most relevant for information spreading. A finite-temperature version of some of our information calculations was given in an appendix of Ref. [12]. A notion of "perfect size winding" [112] has also been shown to be related to optimal many-body teleportation at finite temperature. Both of these cases are related to the (1/2, 0, 1/2, 0) and (1/4, 1/4, 1/4, 1/4) OTOCs, but there is still work to do to elucidate the general structure of information spreading at finite temperature. There are several proposals to measure the OTOC at finite temperature for different regularizations [98,104,113]; recently, the OTOC with regularization (1/2, 0, 0, 1/2) has been measured experimentally in a small system [114].

## **B.** Symmetries

Symmetries and conservation laws can strongly affect scrambling dynamics. The basic intuition is that an initial state cannot scramble as much in the presence of conserved quantities due to the restricted Hilbert space. For instance, the early growth rate of the OTOC is suppressed by temperature [11] and/or chemical potentials [115,116]. Moreover, in extended systems, conserved quantities usually lead to a slow mode associated with their transport, which can significantly slow down the scrambling dynamics and cause a prolonged power-law tail of the OTOC in the late-time regime [53,77–79,117].

An interesting question regarding symmetry is how it affects the information recovery fidelity in the Hayden-Preskill protocol discussed in Sec. IV. It boils down to studying the late-time value of the OTOC appearing in Eq. (72), which we repeat here:

$$F_{\rm EPR} = \left(\frac{1}{4^{|E|}} \frac{1}{2^N} \sum_{W_1, V_E} \operatorname{tr}(W_1(-t)V_E W_1(-t)V_E)\right)^{-1}$$
(142)

with  $W(-t) = UWU^{\dagger}$ . This equation holds for any unitary operator U acting on qubit systems. In the presence of conserved quantity Q, we have [U, Q] = 0. Because of the block diagonalized structure U, the late-time value of OTOCs cannot be estimated using the Haar random unitary such as in Eq. (84). Instead of scaling with  $1/2^N$ , one can show that in systems with conserved charge or energy, the late-time value of the OTOCs scales as O(1/poly(N))[116,118]. Larger symmetry groups can even lead to a finite late-time value of OTOCs and thus significantly suppress the recovery fidelity [119,120].

#### C. More on wavefront broadening

Another direction concerns conjectured universality in the OTOC structure in locally interacting systems. In Sec. V C we indicated that semiclassical or large-N models and random circuit models gave two distinct classes of OTOC behavior, and we argued that the random circuit behavior was generic, i.e., that finite-N corrections would qualitatively change the large-N form of the OTOC. It is important to better understand the universality classes that can arise, especially in Hamiltonian systems at noninfinite temperature. In the literature, OTOCs are sometimes analyzed along a ray within the light cone in terms of a velocity-dependent Lyapunov exponent [74,121],

$$C(x,t) \sim e^{\lambda(v)t},\tag{143}$$

where v = |x|/t. For instance,  $\lambda(v) = \lambda(1 - v/v_B)^{p+1}$  for the growth form in Eq. (94). The definition of the butterfly velocity is  $\lambda(v_B) = 0$ , and at large N one typically finds that  $\partial_v \lambda(v)|_{v_B} \neq 0$ . This corresponds to p = 0 in our discussion above. By contrast, the one-dimensional random circuit form has  $\partial_v \lambda(v)|_{v_B} = 0$ , which corresponds to p = 1. As we said, the latter form is conjectured to be generic in one dimension, but it is important to better understand this issue.

## **D.** Quantum chaos

Another motivation for studying OTOCs in the literature lies in their connection to quantum chaos. In the semiclassical regime, schematically, one can replace the commutator in the squared commutator with a Poisson bracket and obtain [25,122]

$$\mathcal{C} = \frac{1}{\mathrm{tr}I} \mathrm{tr}([x(t), p]^2) \approx \{x(t), p\}^2 = \left(\frac{\partial x(t)}{\partial x}\right)^2, \quad (144)$$

which measures the sensitivity of the final position to the initial position, and thus classical chaos. Then it seems natural to promote the squared commutator to be a diagnosis of quantum chaos [11]. However, it turns out to be quite subtle [123]. For instance, although C is expected to exhibit early-time Lyapunov growth in the semiclassical limit of many-body systems before saturation [122,124,125], not all quantum many-body systems have a semiclassical limit, some large-N limit in which quantum fluctuation is suppressed. The random quantum circuit introduced in Sec. VD does not have the exponential growth behavior expected to be a diagnosis of quantum many-body chaos. Furthermore, some integrable systems exhibit early-time exponential growth due to unstable dynamics [126–128]. Therefore, while the late-time value of the OTOC has a clear physical meaning, as discussed in Sec. V, the connection of the early-time growth of the squared commutator to previously proposed measures of quantum many-body chaos (see Refs. [129,130] for discussions), such as the spectral form factor from the random matrix behavior, needs to be further settled. A more general question is whether quantum scrambling and quantum many-body chaos measure the same property, and if not, when they differ from each other. This discussion also largely hinges on a precise definition of many-body chaos.

#### E. More on numerical methods

Ongoing experiments on scrambling and nonequilibrium quantum many-body dynamics in general are reaching system sizes beyond the capability of exact diagonalization, calling for new numerical tools. In Sec. VI, we discussed MPO- and MPS-based methods to calculate the tail of OTOCs for large system sizes in one dimension, utilizing the light-cone structure of operator spreading. One direction is to extend such a method to higher dimensions using other ansatzes for states or operators. For example, Wu *et al.* [131] calculated the OTOCs of the mixedfield Ising model in two dimensions using the restricted Boltzmann machine.

Moreover, it would be ideal to develop general-purpose numerical tools to directly simulate many-body teleportation for arbitrary unitary circuits with local structure and initial states for large system sizes of about 100 qubits. The conventional MPS- or MPO-based method does not work due to the rapid growth of the entanglement entropy that an MPS with a finite bond dimension cannot capture. It would be interesting to explore the effectiveness of other wavefunction ansatzes, such as various neural network states or the multiscale entanglement renormalization ansatz, in this context.

Understanding scrambling dynamics, i.e., how information flows from local to nonlocal degrees of freedom, is also useful for developing numerical methods to simulate conventional thermalization dynamics of local observables and calculate transport coefficients, which are usually given by time-ordered two-point correlation functions. Schematically, in a strongly interacting system, the equation of motion of single-qubit observables depends on the equation of motion of two-qubit observables, which depends on three-qubit observables, leading to an infinite series of equations involving arbitrary orders of correlation functions that becomes impractical to solve. Transport properties are captured by low-order correlations. It is tempting to assume that the dynamics of sufficiently high-order correlation functions do not feed back to the dynamics of simple correlation functions and to truncate the infinite series of equations or simplify the higher-order equations by approximation. These ideas have led to multiple new numerical algorithms [132–136]. There is still work to do to justify the assumptions and better understand the interplay between the scrambling dynamics and dynamics of local observables.

We believe that we are only just starting to explore this exciting field of quantum information scrambling. With the many connections discovered so far and the prospect of new large-scale experiments on the horizon, there are many exciting possibilities to explore, including the discovery of new maximally chaotic quantum systems, the laboratory simulation of holographic models of quantum gravity, a deeper understanding of quantum chaos, new insights into transport in strongly interacting systems, and much else. We therefore hope that the reader will consider getting into this field and bringing a new point of view.

# ACKNOWLEDGMENTS

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# APPENDIX A: USEFUL DEFINITIONS, OPERATOR IDENTITIES, AND ENTROPY INEQUALITIES

## 1. Matrix norm

We summarize various matrix norms used in this tutorial. Given an operator O acting on a Hilbert space with dimension d, we denote by  $\lambda_i$  the singular values of O. By definition, the  $\lambda_i$  are either positive or zero. We order the singular values so that  $\lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_d$ . Equivalently, the  $\lambda_i$  are the positive square roots of the eigenvalues of the Hermitian operator  $O^{\dagger}O$ . Throughout this tutorial, we use the following conventions for matrix norms,

$$\|O\|_1 = \sum_{i}^d \lambda_i = \operatorname{tr}(\sqrt{O^{\dagger}O}), \qquad (A1a)$$

$$\|O\|_2 = \sqrt{\sum_{i}^{d} \lambda_i^2} = \sqrt{\operatorname{tr}(O^{\dagger}O)}, \qquad (A1b)$$

$$\|O\|_{\infty} = \lambda_d. \tag{A1c}$$

Here  $||O||_1$  is the trace norm,  $||O||_2$  is the Frobenius norm, and  $||O||_{\infty}$  is the operator norm, which is the largest singular value of *O*. We have the inequalities

$$\frac{1}{d} \|O\|_2^2 \le \|O\|_\infty^2.$$
 (A2)

## 2. Bipartite entanglement entropy

#### a. State entanglement

In this section, we provide the definition of the bipartite entanglement entropy for a pure state. We consider a pure state  $|A\bar{A}\rangle$  of a quantum system consisting of a region A and its complement  $\bar{A}$ . One can perform Schmidt decomposition on the state,

$$|A\bar{A}\rangle = \sum_{n} \lambda_n |A\rangle_n \otimes |\bar{A}\rangle_n.$$
 (A3)

The states  $|A\rangle_n$  and  $|\bar{A}\rangle_n$  form an orthogonal basis in regions A and A, respectively, namely,  $\langle A_n | A_m \rangle = \delta_{mn}$  and  $\langle \bar{A}_n | \bar{A}_m \rangle = \delta_{mn}$ . The coefficients  $\lambda_n$ , called the singular values, are all positive. Normalization requires that

$$\sum_{n} \lambda_n^2 = 1. \tag{A4}$$

From the Schmidt decomposition, one can obtain the reduced density matrix of region *A*,

$$\rho(A) = \sum_{n} \langle \bar{A}_{n} | A\bar{A} \rangle \langle A\bar{A} | \bar{A}_{n} \rangle$$
$$= \sum_{n} \lambda_{n}^{2} | A \rangle_{n} \langle A |_{n}.$$
(A5)

Then the von Neumann bipartite entanglement entropy is

$$S(A) = -\operatorname{tr}(\rho(A)\log(\rho(A))) = -\sum_{n} \lambda_n^2 \log \lambda_n^2. \quad (A6)$$

One can generalize it to the Rényi entropies

$$S^{(\alpha)}(A) = \frac{1}{1-\alpha} \log(\mathrm{tr}\rho^{\alpha}) = \frac{1}{1-\alpha} \log \sum_{n} \lambda_{n}^{2\alpha} \quad (A7)$$

for  $\alpha > 0$ . Taking the limit  $\alpha \to 1$  yields the von Neumann entropy.

Crucially, since the entanglement entropies depend only on the singular values  $\lambda_n$  and not on the basis states, the entanglement entropies of region A are always the same as those of its complement  $\overline{A}$ ,

$$S^{(\alpha)}(A) = S^{(\alpha)}(\bar{A}). \tag{A8}$$

Note that this is only true for a pure state.

#### **b.** Operator entanglement

One can generalize the concept of the bipartite entanglement entropy to a many-body operator  $\hat{O}$ . The operator entanglement quantifies how well an operator can be approximated by the product of two operators,  $O_A$  and  $O_{\bar{A}}$ , in two complement regions.

The definition of the operator entanglement follows closely that of a pure state. The key idea is to map a manybody operator  $\hat{O}$  to a fictitious quantum state in a doubled Hilbert space  $|O\rangle$ . Then, the operator entanglement entropy of O is defined as the regular state entanglement entropy of  $|O\rangle$ , which is defined in the previous section.

Let us first consider a single-qudit operator to understand the map. In the computational basis, the operator takes the form

$$\hat{O} = \sum_{nm} |n\rangle O_{nm} \langle m|, \qquad (A9)$$

and the corresponding two-qudit state is

$$|O\rangle = \frac{1}{\sqrt{d}} \sum_{nm} O_{nm} |n\rangle |m\rangle.$$
 (A10)

Formally, the state is obtained by acting  $\hat{O}$  on the first qudit of an EPR pair, defined as

$$|I\rangle = \frac{1}{\sqrt{d}} \sum_{n} |n\rangle |n\rangle.$$
 (A11)

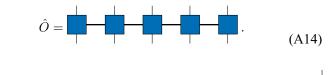
It is straightforward to show that  $\hat{O} |I\rangle = |O\rangle$ , and the state is normalized if the normalization of the operator is fixed to  $tr(OO^{\dagger}) = d$ . The following graph shows the map for the single-qudit operator:

$$\hat{O} = -$$
,  $|I\rangle = -$ ,  $|O\rangle = \hat{O}|I\rangle = -$ . (A12)

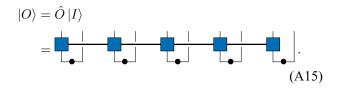
For N qudits, we define the state  $|I\rangle$  of 2N qudits as a product of N local EPR pairs,

$$|I\rangle = \frac{1}{d^{N/2}} \prod_{i}^{N} \sum_{n} |n_{i}\rangle |n_{i}\rangle$$
$$= \bigsqcup_{\bullet} \bigsqcup_{\bullet} \bigsqcup_{\bullet} \bigsqcup_{\bullet} \bigsqcup_{\bullet} \bigsqcup_{\bullet} \bigsqcup_{\bullet} \bigsqcup_{\bullet} \bigsqcup_{\bullet} (A13)$$

Now the operator  $\hat{O}$ , acting on N qudits in the original system, can be represented as the following tensor network (note that the bond dimension of the horizontal line can in general be exponentially large for complicated operators):



The operator state  $|O\rangle$  for an *N*-qudit operator  $\hat{O}$  is obtained by acting  $\hat{O}$  on the odd qudits in state *I*,



The graph illustrates the operator-to-state mapping for five qudits in the original system (equivalent to ten qudits in the doubled system). If  $\hat{O}$  is a product operator, for instance,  $I \otimes I \otimes X \otimes X \otimes I$ , then the corresponding operator state is a product state with zero entanglement. The bipartite entanglement of the state after the mapping, dubbed the operator entanglement, characterizes the entanglement structure of the operator. It quantifies how well a manybody operator can be approximated by the product of two operators in complement regions. The operator entanglement also determines the required bond dimension of the MPO to represent the operator faithfully and thus quantifies the computation resource to store the operator classically.

# 3. Entropy inequalities

The following is a list of useful entropy inequalities:

$$\begin{split} \langle \psi | \phi \rangle &\leq \sqrt{\langle \psi | \psi \rangle \langle \phi | \phi \rangle} \quad \text{(Cauchy-Schwarz inequality),} \\ S(A) + S(B) &\geq S(AB) \quad \text{(subadditivity),} \\ |S(A) - S(B)| &\leq S(AB) \quad \text{(triangle inequality),} \\ S(AB) + S(BC) &\geq S(ABC) + S(B) \quad \text{(strong subadditivity),} \\ \text{tr}\rho_1(\log_2 \rho_1 - \log_2 \rho_2) &\geq \frac{1}{2\ln 2} \|\rho_1 - \rho_2\|_1^2 \quad \text{(Pinsker's inequality)} \end{split}$$

We recommend the lecture notes by Preskill [137, Chap. 10] on quantum computation and the introductory article by Witten [138] for a brief description of entropy inequalities.

The ensemble average of the  $d \times d$  Haar random unitary matrix obeys

## 4. Operator identities

Now we summarize some operator identities that are repeatedly used throughout this article. For a quantum many-body system with dimension d, a complete operator basis S for a quantum system satisfies the conditions

$$\frac{1}{d}(\mathcal{S}^{\dagger}\mathcal{S}') = \delta_{\mathcal{S}\mathcal{S}'}, \qquad \frac{1}{d}\sum_{\mathcal{S}}\mathcal{S}_{ab}^{\dagger}\mathcal{S}_{cd} = \delta_{ad}\delta_{bc}.$$
 (A16)

$$\mathbb{E}\left(U_{a'a}U_{b'b}^{*}\right) = \frac{1}{d}\delta_{a'b'}\delta_{ab}, \qquad (A17a)$$

$$\mathbb{E}\left(U_{a'a}U_{b'b}^{*}U_{c'c}U_{d'd}^{*}\right)$$

$$= \frac{1}{d^{2}-1}\left(\delta_{a'b'}\delta_{c'd'}\delta_{ab}\delta_{cd} + \delta_{a'd'}\delta_{b'c'}\delta_{ad}\delta_{bc}\right)$$

$$- \frac{1}{d(d^{2}-1)}\left(\delta_{ab}\delta_{cd}\delta_{a'd'}\delta_{b'c'} + \delta_{a'b'}\delta_{c'd'}\delta_{ad}\delta_{bc}\right). \qquad (A17b)$$

# APPENDIX B: OPERATOR STRINGS AND THE OTOC IN QUDIT SYSTEMS AND MAJORANA SYSTEMS

# 1. OTOC in qudit systems

As a generalization of the Pauli matrix, a complete basis of operators in *q*-dimension Hilbert space can be defined as

$$\sigma^{mn} = \sum_{k=0}^{q-1} |k\rangle \langle k+m| \exp\left(i\frac{2\pi}{q}kn\right).$$
(B1)

This is a unitary but not a Hermitian basis. These operators obey

$$\frac{1}{q} \operatorname{tr}(\sigma^{kn^{\dagger}} \sigma^{k'n'}) = \delta_{kk'} \delta_{nn'}.$$
(B2)

They also satisfy the completeness relation

$$\frac{1}{q}\sigma_{ab}^{\dagger,mn}\sigma_{cd}^{mn} = \delta_{ad}\delta_{bc}.$$
(B3)

Similar to the qubit system, a Heisenberg operator in the system containing multiple qudits can be expanded in the operator string basis

$$W(t) = \sum_{\mathcal{S}} \alpha(\mathcal{S}), \tag{B4}$$

where S is a product of operators  $\sigma_r^{mn}$  acting on each qudit. Using the completeness relation, one can show that the averaged OTOC is

$$\frac{1}{q^2} \frac{1}{q^N} \sum_{mn} \operatorname{tr}(W^{\dagger}(t)\sigma_r^{mn,\dagger}W(t)\sigma_r^{mn})) = \sum_{\mathcal{S}_r=I} |\alpha(\mathcal{S})|^2.$$
(B5)

It measures the probability that the operator on the *r*th qubit in W(t) is the identity. From this relation, one can also show that the averaged squared commutator

$$\frac{1}{q^2 - 1} \sum_{mn} \frac{1}{q^N} \| [W(t), \sigma_r^{mn}] \|_2^2 = \frac{2q^2}{q^2 - 1} \sum_{\mathcal{S}(r) \neq I} |\alpha(\mathcal{S})|^2.$$
(B6)

In the scrambling limit, the operator string reaches a local equilibrium where each on-site operator is equally probable, and the averaged squared commutator reaches 2.

#### 2. OTOC in Majorana systems

Consider a system of N Majorana fermions, labeled  $\chi_{\alpha}$ , which obeys the commutation relation  $\{\chi_{\alpha}, \chi_{\beta}\} = \delta_{\alpha\beta}$ .

A conventional basis for operators in this system is the Majorana string

$$S = i^{m(m-1)/2} 2^{m/2} s_1 s_2 \cdots s_N, \tag{B7}$$

where  $s_r$  can be either  $\chi_r$  or  $I_r$  and m counts the number of Majorana operators in the string. The factor  $i^{m(m-1)/2}$  is to ensure the Hermicity of S and  $2^{m/2}$  is to ensure the normalization tr $S^2 = \text{tr}I$ . The Heisenberg operator  $\chi(t)$  can be expanded in this basis as

$$\chi(t) = \frac{1}{\sqrt{2}} \sum \alpha(\mathcal{S}, t) \mathcal{S}.$$
 (B8)

The expansion contains only an odd number of Majorana operators because the fermion parity is conserved. The unitary quantum dynamics ensures that  $\sum_{S} |\alpha(S)|^2 = 1$  for all time. Therefore,  $|\alpha(S)|^2$  can be interpreted as a probability distribution. Then the OTOC can be written as

$$F(r,t) = \frac{1}{\text{tr}I} \text{tr}(\chi(t)\chi_r\chi(t)\chi_r) = -\frac{1}{4} + \frac{1}{2}\sum_{s_r = \chi_r} |\alpha(\mathcal{S})|^2.$$
(B9)

It is directly related to the probability that the Majorana operator  $\chi_r$  appears in the operator string. In the scrambling limit, the Majorana operator and the identity operator have the same probability to appear and F(r, t) approaches 0. In the Majorana system, the counterpart of the squared commutator is the squared anticommutator

$$C(r,t) = \frac{4}{\text{tr}I} \text{tr}(\{\chi(t),\chi_r\}\{\chi(t),\chi_r\}) = 4\sum_{s_r = \chi_r} |\alpha(\mathcal{S})|^2.$$
(B10)

In the scrambling limit, C(r, t) approaches 2, the same value as that in the qudit system.

#### **APPENDIX C: OTHER TOY MODELS**

This section sketches how the OTOC is calculated in various models in addition to the random circuit model presented in the main text.

#### 1. Free fermions

We start with noninteracting systems. Although one should not expect these models to be generic, they provide important lessons on possible functional forms regarding the tail of the OTOC since the Lieb-Robinson bound applies to the noninteracting systems as well. Consider a noninteracting Majorana system described by the Hamil-tonian

$$H = \frac{1}{2} \sum_{a,b} h_{rr'} \chi_r \chi_{r'}, \qquad (C1)$$

where  $\chi$  is the Majorana operator that obeys the commutation relation  $\{\chi_r, \chi_{r'}\} = \delta_{rr'}$ . The Hermicity of Hrequires that  $h_{rr'} = -h_{rr'}^*$  and  $h_{rr'} = -h_{r'r}$ . Therefore, h is a purely imaginary antisymmetric matrix. In addition, we assume that H is translational symmetric and can thus be diagonalized by a Fourier transformation. Introduce  $\chi_k = \sum_r (1/\sqrt{N})e^{ikr}\chi_r$  satisfying  $\{\chi_k, \chi_{k'}\} = \delta_{k,-k'}$ . Then the Hamiltonian can be written as

$$H = \frac{1}{2} \sum_{k} \epsilon_k \chi_{-k} \chi_k.$$
 (C2)

The spectrum  $\epsilon_k$  is an odd function in k because matrix h is imaginary and antisymmetric. The goal here is to calculate the OTOC

$$F(r,t) = \frac{1}{\mathrm{tr}I} \mathrm{tr}(\chi_0(t)\chi_r\chi_0(t)\chi_r), \qquad (C3)$$

or, equivalently, the squared anticommutator defined in Appendix B 2,

$$C(r,t) = \frac{4}{\mathrm{tr}I} \mathrm{tr}(\{\chi_0(t),\chi_r\}\{\chi_0(t),\chi_r\}).$$
(C4)

Using the Hamiltonian in the momentum space, one can show that the Heisenberg operator  $\chi_0(t)$  is

$$\chi_0(t) = \sum_r g(r)\chi_r,\tag{C5}$$

where

$$g(r) = \sum_{k} \frac{1}{N} e^{i\epsilon_k t - ikr} = \frac{1}{2\pi} \int e^{i\epsilon_k t - ikr} dk.$$
 (C6)

Because  $\epsilon_k = -\epsilon_{-k}$ , the coefficient g(r, t) is real, as expected from expanding the Hermitian operator  $\chi_0(t)$ . The fact that only a single Majorana operator appears in the expansion is because of the noninteracting Hamiltonian. In general, Majorana strings of all lengths would appear. From g(r), one can obtain the OTOC as

$$C(r,t) = 4g^2(r,t).$$
 (C7)

Therefore, the behavior of C(r, t) is controlled by g(r, t), which can be analyzed using the saddle-point approximation. We expand the function around  $k_0$  as

$$g(r,t) \sim \int \frac{d\delta k}{2\pi} \exp\left\{i(v(k_0)t - x)\delta k + \frac{i}{2}\epsilon_{k_0}^{(2)}t\delta k^2 + \frac{i}{6}\epsilon_{k_0}^{(3)}t\delta k^3 + \cdots\right\},$$
 (C8)

where  $v(k_0)$  is the group velocity at  $k_0$ . One can always find suitable  $k_0$  that make the first derivative term vanish for  $|x| < \max(|v(k_0)|)t$ , while it is not possible to do so if  $|x| > \max(|v(k_0)|)t$ . This leads to a change of behavior of g(x, t) at  $|x| = \max(|v(k_0)|)t$ , indicating that the butterfly velocity  $v_B$  is the maximal velocity. In the region  $x > v_B t$ , the first-order derivative term is always nonzero. We keep up to the third-order term and obtain

$$g(r,t) \sim \frac{1}{2\pi} \int d\delta k e^{i(v_B t - r)\delta k + i\epsilon_{k_0}^{(3)} t\delta k^3/6}$$
$$\sim \frac{1}{t^{1/3}} \operatorname{Ai}\left(\frac{r - v_B t}{(-\epsilon_{k_0}^{(3)} t/2)^{1/3}}\right), \tag{C9}$$

where Ai(Z) is the Airy function. Note that here  $\epsilon_{k_0}^{(3)}$  is negative since the group velocity is maximal at  $k_0$ . In the limit that  $r - v_B t \gg |\partial_k^3 \varepsilon(k_0) t/2|^{1/3}$ , we can use the asymptotic form of the Airy function and obtain Eq. (93). In this case, the wavefront broadens subdiffusively with  $\delta r \sim t^{1/3}$ and the broadening exponent p = 1/2. On the other hand, in the long time limit  $t \gg r/v_B$ , one can find a  $k_0$  that makes the first derivative vanish and perform the Gaussian integral to get  $g(r, t) \sim 1/t^{1/2}$ . As a result, in the long time limit C(r, t) decays to 0 as 1/t in the noninteracting systems. This is in sharp contrast with a scrambling system where C(r, t) approaches 2 in the long time limit. These results from saddle-point analysis can be explicitly checked by exact calculation using the nearest-neighboring hopping model given by  $h_{rr'} = i\delta_{r,r'+1} - i\delta_{r,r'-1}$ . In this case,  $\epsilon_k = \sin k$  and  $g(r, t) = J_r(t)$ , a Bessel function.

## 2. Brownian models

We have seen that the random circuit model and the SYK model give rise to distinct functional forms of the OTOC. Both cases feature a ballistic light cone, but the wavefront broadens diffusively in the random circuit model, while the wavefront is sharp in the SYK model. The question then arises which, if any, of these characteristic shapes describes the generic case with a finite local Hilbert space dimension. Unfortunately, this question cannot be reliably answered using small-sized numerical simulations. These exhibit ballistic expansion with some broadened wavefront, but it is unclear if the broadening will vanish in a large system, tend to the diffusive limit, or have some other characteristic form. Noninteracting particles exhibit a ballistic expansion of C with yet another characteristic broadening of the wavefront (C does not saturate at late time in this model, instead falling back to zero). One should not expect the noninteracting limit to be generic, but the spectrum of multiple different universality classes is certainly raised.

To answer this question, we now introduce another class of models known as Brownian models [10,35-37,139-141], in which the interaction between the underlying degrees of freedom are stochastic variables. Here we consider a specific version of the model with spin-spin interactions, called the Brownian coupled cluster model. Like the random circuit model, it features a random time-dependent Hamiltonian, but unlike the random circuit model, it has a large-N limit. Using it, we develop a physical picture of why p = 1 is generic for one-dimensional scrambling systems. The model can be defined in any dimension, but here we continue to focus on d = 1. The degrees of freedom are arranged in clusters that are then connected in a one-dimensional array. Every cluster contains N spin-1/2 degrees of freedom, and there are L clusters. The Hamiltonian is time dependent and consists of two kinds of terms, within-cluster interactions and between-cluster interactions. In order to avoid mathematical complexities associated with stochastic calculus, it is simplest to present the model in discrete time.

The time evolution operator is

$$U(t) = \prod_{m=1}^{t/dt} \exp\left(-i\sum_{r} H_{r}^{(m)} - i\sum_{\langle rr' \rangle} H_{rr'}^{(m)}\right)$$
(C10)

with m a discrete time index. The within-cluster terms and the between-cluster terms are

$$H_r^{(m)} = J_{m,r,a,b}^{\alpha\beta} \sigma_{r,a}^{\alpha} \sigma_{r,b}^{\beta}, \tag{C11}$$

$$H_{rr'}^{(m)} = g \tilde{J}_{m,r,r',a,b}^{\alpha\beta} \sigma_{r,a}^{\alpha} \sigma_{r',b}^{\beta}, \qquad (C12)$$

where  $\alpha, \beta \in \{0, 1, 2, 3\}, a, b = 1, ..., N$  label spins within a cluster, r, r' label clusters (sometimes called sites), and  $\langle rr' \rangle$  means nearest neighbors. At each time step, the models contain two sets of uncorrelated random variables J and  $\tilde{J}$  with mean zero and variance dt/8(N-1) and dt/16N, respectively.

In the limit that  $dt \to 0$ , one can formulate a stochastic differential equation for the time evolution operator. From it, one can derive a master equation for the operator probabilities  $\overline{|c(S)|^2}$  averaged over circuit realizations, i.e., over realizations of the couplings J and  $\tilde{J}$ . We will not get into

the details of these equations here, but refer the reader to Ref. [36] for complete details. The only important property we need is that  $\overline{|c(S)|^2}$  depends only on the total number of nonidentity Pauli operators on each cluster. This is technically an approximation, but it holds after a short time even if the initial condition does not obey it because the circuit average erases any distinction between the different Pauli operators. The total number of nonidentity Pauli operators in  $P_S$  on cluster r is called the weight of the cluster and is denoted  $w_r(S)$ .

In this model, the operator-averaged squared commutator C has a cluster index as well as the index of the qubit within a cluster,

$$C^{a}(r,t) = \frac{1}{3 \text{tr}I} \sum_{S_{r,a}} \|W(t), S_{r,a}\|_{2}^{2}.$$
 (C13)

According to Eq. (82b),  $C^a(r, t) = \frac{8}{3} \sum_{S_{r,a} \neq I} |\alpha(S, t)|^2$ . It is convenient to analyze the  $C^a(r, t)$  averaged over all spins within a cluster r,

$$\phi(r,t) \equiv \frac{1}{N} \sum_{a} C^{a}(r,t) = \frac{8}{3} \sum_{S} w_{r}(S) |\alpha(S)|^{2}, \quad (C14)$$

which measures the averaged number of nonidentity operators that appears over the *N* spins, or the averaged operator weight  $\langle w_r \rangle$  within cluster *r*. At early time,  $\phi(r, t) \approx 0$ , while at late time it saturates to  $\phi(r, t) = 2$ , as the usual case.

Like the Haar random circuit, using random averaging of coupling, one can derive a master equation governing the dynamics of the operator probability  $|\alpha(S)|^2$ . When N is small, it can be shown that C(r, t) obeys a drift-diffusion equation, as in the random circuit model. This leads to a circuit averaged  $\phi(r, t)$  obeying the universal form in Eq. (94) with p = 1. Hence, the Brownian coupled cluster model recovers the result of the random circuit model at small N.

At infinite N, something very different occurs. It can be shown that  $\phi(r, t)$  obeys a so-called Fisher-Kolmogorov-Petrovksy-Piskunov- (FKPP) type equation of the form

$$\partial_t \phi = \frac{3}{2}(2-\phi) \left( \frac{g^2}{2} \partial_r^2 \phi + (1+g^2) \phi \right).$$
(C15)

Here g is the ratio of the strength of the between-cluster and within-cluster terms, and we have taken a continuum limit, which is qualitatively accurate. Although we will not explain its detailed derivation, one can see that this equation contains three essential pieces of physics: exponential growth in time, spreading in space, and saturation. The FKPP equation is very well known and describes a wide variety of physical processes, including the propagation of combustion waves, the dynamics of invasive species, and the physics of certain quantum chromodynamics processes.

The key physical property of the FKPP equation is that, starting from a localized source, it supports travelingwave solutions with  $C(r, t) = f(r - v_B t)$ , where  $v_B = \sqrt{18g^2(1+g^2)}$  is the butterfly velocity. Well ahead of the front at  $r = v_B t$ , the waveform is

$$\phi(r,t) \sim e^{\lambda(t-r/v_B)},\tag{C16}$$

which is Eq. (94) with p = 0. Hence, the Brownian coupled cluster model also recovers the physics of large-*N* and semiclassical results. The exponent  $\lambda = 6(1 + g^2)$  is an example of a quantum Lyapunov exponent.

Given the large- and small-N limits, the next question is how they are connected as N is varied. Physically, the infinite N limit functions to suppress quantum fluctuations, so that one may view the distribution  $\overline{|c(S)|^2}$  as being concentrated on a single weight configuration. At finite N, quantum fluctuations occur, meaning that the distribution  $|c(S)|^2$  now assigns nonvanishing probability to different operator weight configurations. It is important to understand that these fluctuations are proper quantum fluctuations. They are a consequence of the fact that W(t) is a superposition of many different Pauli strings of different weights. In particular, the randomness associated with the couplings J and  $\tilde{J}$  has already been averaged over and no longer enters the description. In essence, the circuit average serves to dephase the quantum operator amplitudes and convert the Heisenberg equation of motion for the operator amplitudes into a master equation for the operator probabilities.

Following Ref. [36], we call these quantum fluctuations "noise." In an abuse of notation where  $\phi(r, t)$  now represents a noisy field, we obtain a noisy FKPP equation,

$$\partial_t \phi = 3(1-\phi)f(\phi) + \sqrt{\frac{1}{N}(2-\phi/2)f(\phi)}\eta(r,t),$$
(C17)

where

$$f(\phi) = \left(\frac{g^2}{2}\partial_r^2 + (1+g^2)\right)\phi,$$

and  $\eta(r, t)$  is a white noise term representing quantum fluctuations. This noise term, while suppressed by 1/N, has a dramatic effect on the physics. Note also that the noise is multiplicative, vanishing when  $\phi = 0$ , so it respects the causal structure.

The main effect of the noise term is to make the front position noise dependent. This means that the front continues to move with velocity  $v_B$ , but it is also randomly buffeted forward and backward as in a random walk. Within a particular noise realization, the wavefront is sharp and exhibits p = 0. However, the physical quantity in the quantum problem is the noise-averaged value of  $\phi$ . Close enough to the physical front at  $r = v_B t$ , the random walk nature of the front position inevitably manifests and smears the sharp p = 0 front into a diffusive p = 1 front. Using the noisy FKPP literature [142], Xu and Swingle [36] showed that the corresponding diffusion constant was  $D \sim 1/\log^3 N$  at large N, a remarkably large value relative to standard 1/N corrections.

### 3. Coupled Sachdev-Ye-Kitaev model

The large-N or semiclassical setting is another class of typically interacting models for which scrambling dynamics and the OTOC is tractable. A representative model in this class is the Sachdev-Ye-Kitaev model [143,144] (see also the review [145]) describing a cluster of N interacting Majorana fermions,

$$H = \sum_{i_1 < i_2 < i_3 < i_4} j_{i_1 i_2 i_3 i_4} \chi_{i_1} \chi_{i_2} \chi_{i_3} \chi_{i_4}, \qquad (C18)$$

where  $\chi_i$  is the Majorana operator obeying the usual commutator  $\{\chi_i, \chi_j\} = \delta_{ij}$ . The coupling constants  $j_{i_1i_2i_3i_4}$  are uncorrelated Gaussian variables with zero mean and standard deviation  $\langle j^2 \rangle = 6J^2/N^3$ , where the factor of  $1/N^3$ is required to make sure that the energy of the system is extensive, i.e., scaling linearly with *N*.

One can generalize the SYK model to higher dimensions by considering M clusters of N fermions [73,146–148]. Each cluster only interacts with its nearest neighbors. The Hamiltonian contains both on-site terms and bond terms,

$$H = \sum_{r} H_r + \sum_{\langle rr' \rangle} H_{rr'}, \qquad (C19)$$

where  $H_r$  is the usual SYK Hamiltonian in Eq. (C18) and the bond term is given by

$$H_{rr'} = \sum_{\substack{1 \le j < k \le N \\ 1 \le l < m \le N}} J'_{jklm, rr'} \chi_{j, r} \chi_{k, r} \chi_{l, r'} \chi_{m, r'}.$$
(C20)

The OTOC of the generalized SYK model can be calculated using the same diagrammatic approach in the large-N limit. The details of the calculation can be found in Refs. [73,149], showing that the OTOC takes the simple exponential form in Eq. (91) up to higher-order corrections.

# APPENDIX D: PHENOMENOLOGICAL DESCRIPTION

Given these developments, a conjecture and a corresponding physical picture naturally present themselves. We claim that, due to the inevitable presence of quantum fluctuations, generic one-dimensional quantum systems always have squared commutators obeying the universal form in Eq. (94) with p = 1. There is an analogous claim in higher dimensions, where the value of p depends on the dimension and is related to a random surface growth problem (the Kardar-Parisi-Zhang universality class). Based on discussions in previous sections, the key pieces of evidence in favor of this claim are the random circuit model and the Brownian coupled cluster model. Interestingly, FKPP-like equations have also been obtained in various large-N and weak coupling calculations of squared commutators. These were all noiseless equations, but surely once quantum fluctuations are included, the dynamics will be governed by an FKPP-like equation with multiplicative noise and a corresponding broadened front.

Based on the previous discussion on the random circuit, SYK model, and Brownian coupled cluster, the universal features of the dynamics of the OTOC are

- (1) ballistic expansion,
- (2) late-time saturation,
- (3) local exponential (Lyapunov) growth,
- (4) random walk near the wavefront.

A phenomenological description that captures all the key features is

$$C(r,t) \sim \frac{1}{\sqrt{2\pi Dt}} \int d\Delta r \frac{e^{-\Delta r^2/2Dt}}{e^{\lambda((r+\Delta r)/v-t)} + 1},$$
 (D1)

which is a convolution between a Fermi-Dirac function and a Gaussian distribution, describing a traveling wave moving at velocity v. In the limit that  $\lambda \to \infty$ , the Fermi-Dirac distribution function becomes a step function, and Eq. (D1) recovers the behavior of the random circuit model. On the other hand, in the limit that  $D \to 0$ , the Gaussian distribution becomes a delta function, and Eq. (D1) reduces to the exponential form for the large-Nmodels. Let us work with the dimensionless variables

$$\tilde{r} = \frac{rv}{D}, \qquad \tilde{t} = \frac{tv^2}{D}, \qquad \xi = \frac{\lambda D}{v^2}.$$
 (D2)

The phenomenological description becomes

$$\mathcal{C}(r,t) = \frac{2}{\sqrt{2\pi}} \int d\Delta r \frac{e^{-\Delta r^2/2}}{e^{\xi \sqrt{t}(\Delta r + (r-t)/\sqrt{t})} + 1}, \qquad (\text{D3})$$

where we have dropped the tilde notation for simplicity, whose behavior is controlled by a single parameter  $\xi$ . The integration does not have a closed form, but we understand its behavior in different space-time regimes.

First, one can show that C(r, t) = 1 when r = t, setting the wavefront as expected. When  $(r - t)/\sqrt{t} \le \text{const}$ , the Fermi-Dirac function approaches a step function in the large-t limit, and

$$C(r,t) \approx \operatorname{erfc}\left(\frac{r-t}{\sqrt{2t}}\right).$$
 (D4)

On the other hand, when  $r - t > \xi t$ , the exponential term dominates the denominator and the integration leads to

$$\mathcal{C}(r,t) \approx 2\exp(\xi((1+\xi/2)t-r)), \qquad (D5)$$

which recovers the large-*N* form. Interestingly, both the Lyapunov exponent and the butterfly velocity increase by a factor  $(1 + \xi/2)$ . We perform the integration numerically and plot  $\log C$  as a function of r - t, fixing *t*, in

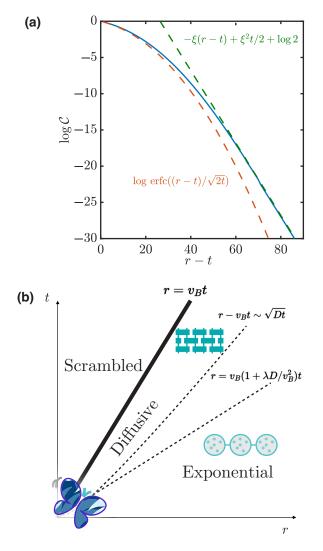


FIG. 10. (a) The squared commuter crossover from the diffusive regime to the exponential regime as r - t increases, i.e., moving away from the light cone. (b) A schematic illustration of the different space-time regions of the squared commutator. Near the light cone, its behavior is dominated by the diffusive region, displaying diffusive broadening.

Fig. 10(a), which clearly demonstrates that  $\log C$  interpolates between the two limiting cases, the error function and the exponential function.

Based on the discussion above, one arrives at the following picture of scrambling dynamics in an extended quantum many-body system with short-ranged interaction and finite local Hilbert space dimension. The information propagates ballistically with a butterfly speed  $v_B$ . Because of inevitable quantum fluctuation, in the spacetime region near  $r - v_B t \sim \sqrt{Dt}$ , which we refer to as the diffusive region, C(r, t) is characterized by a diffusive broadened wavefront. In this region, there is no well-defined Lyapunov exponent that is independent of position and velocity. However, far ahead of the wavefront  $r - v_B t > \lambda D/v_B t$ , C(r, t) grows exponentially with a well-defined Lyapunov exponent, and we thus refer to this region as the exponential region.

Still, it would be nice to check this claim that the wavefront broadens diffusely in a quantum spin chain with no randomness in space or time and generic interactions. More generally, the preceding discussion did not provide a method to calculate squared commutators for generic physical systems. One may wonder if there exist general methods or numerical algorithms to study the OTOC for these systems. This will be the topic of the next section. Before moving to the details of these methods, we would like to point out that numerical results strongly support diffusive broadening of the information wavefront in spin chains with large system sizes. Figure 11 shows data obtained from the matrix product operator approach introduced in Sec. VI for the mixed-field Ising model at infinite temperature. The simulation is performed for n = 201 spins out to quite a long time, several hundred 1/J. It has been

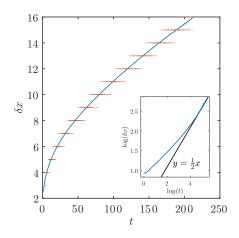


FIG. 11. Separation between two different contours of constant *C* as a function of time for the mixed-field Ising model. The inset shows a log-log plot of the same data. The asymptotic approach to a slope of 1/2, corresponding to p = 1, is clearly visible in the data. Taken from Fig. 4(b) of Ref. [36].

checked that the results are converged in the bond dimension with a bond dimension as low  $\chi = 32$ . What is plotted are the contours of the constant squared commutator. The inset shows the difference between different contours as a function of time on a log-log plot. Equation (95) predicts that the difference between contours should go like  $t^{p/(p+1)}$ , so on a log-log plot the data should approach a straight line of slope p/(p + 1). This is precisely what occurs with an asymptotic slope of 1/2, corresponding to p = 1. Hence we verify that, for a large, nonrandom interacting spin chain, the operator growth dynamics is ballistic with a diffusively broadened front, exactly as predicted.

The chaotic regime ahead of the front should also be present, but is difficult to observe. Based on the picture shown in Fig. 11(b), the size of the chaotic regime is determined by the ratio  $v_B^2/\lambda D$ . Therefore, one may observe the exponential growth in systems with large  $v_B$  but small  $\lambda$ , as shown in Ref. [150].

## **APPENDIX E: LIEB-ROBINSON BOUND**

This appendix reviews an elementary proof of a Lieb-Robinson bound for a simple one-dimensional spin to give a sense of how it works. The analysis follows a discussion of Osborne. Let us assume that the Hamiltonian can be written as a sum of terms  $h_r$  that act on sites r and r + 1. This can always be done coarse graining any finite-range interaction. Let the operator norm of  $h_r$  be J, which measures the local energy scale of the Hamiltonian.

Consider an operator W located at site  $r_0$ . The goal of the Lieb-Robinson bound is to upper bound how far from  $r_0$  this operator can spread after time t. The rough idea is to consider a series of approximations to W(t) that involve truncating more and more distant terms in the Hamiltonian. These truncations then converge, roughly speaking, to W(t), while also giving a bound on the spreading.

Denote the restriction of *H* to the interval  $[r - \ell, r + \ell]$  by  $H_{\ell}$ , which means keeping only terms from *H* that are fully supported on the interval. The restricted Hamiltonian reads

$$H_{\ell} = \sum_{r=r_0-\ell+1}^{r_0+\ell-1} h_r$$
(E1)

with the dependence on  $r_0$  suppressed. Using the  $H_\ell$ , define the sequence of Heisenberg operators  $W_\ell$  via

$$W_{\ell} = e^{iH_{\ell}t}We^{-iH_{\ell}t}.$$
 (E2)

To quantitatively estimate how these terms differ from each other, define the norms  $\alpha_{\ell}$  by

$$\alpha_{\ell} = \|W_{\ell} - W_{\ell-1}\|_{\infty}$$
(E3)

with  $\alpha_0 = \|W\|_{\infty}$ . In terms of these, it is possible to upper bound objects of the form  $\|W_{\ell} - W_{\ell'}\|_{\infty}$  as

$$\|W_{\ell} - W_{\ell'}\|_{\infty} \le \sum_{j=\ell'+1}^{\ell} \alpha_j,$$
 (E4)

by repeatedly adding and subtracting the  $W_{\ell''}$  and using the triangle inequality.

The  $\alpha_{\ell}$  are determined using a differential equation,

$$\frac{d}{dt}\alpha_{\ell} \le \left\|\frac{d(W_{\ell} - W_{\ell-1})}{dt}\right\|_{\infty}.$$
(E5)

Using the invariance of the norm under unitary transformations, the right-hand side can be equivalently written as

$$\left\|\frac{d}{dt}(U_{\ell+1}W_{\ell}U_{\ell+1}^{\dagger}-W)\right\|_{\infty},$$
(E6)

which is

$$\|[-iH_{\ell+1}, W_{\ell}] + [iH_{\ell}, W_{\ell}]\|_{\infty} = \|[H_{\ell+1} - H_{\ell}, W_{\ell}]\|_{\infty}.$$
(E7)

The next step identifies  $H_{\ell+1} - H_{\ell}$  with  $h_{r_0+\ell} + h_{r_0-\ell}$ , since these are the only new terms in  $H_{\ell+1}$  fully supported on  $[r_0 - \ell - 1, r_0 + \ell + 1]$  but not fully supported on  $[r_0 - \ell, r_0 + \ell]$ . We also use the fact that  $W_{\ell-1}$  has no nontrivial support on  $r_0 \pm \ell$  or  $r_0 \pm (\ell + 1)$  and hence commutes with  $H_{\ell+1} - H_{\ell}$ . Thus, the right-hand side of the  $\alpha_{\ell}$  differential equation can be taken to be

$$\|[H_{\ell+1} - H_{\ell}, W_{\ell} - W_{\ell-1}]\| \le 4J \|W_{\ell} - W_{\ell-1}\| = 4J\alpha_{\ell-1}.$$
(E8)

Using  $||AB|| \le ||A|| ||B||$  and the triangle inequality, one has

$$\frac{d\alpha_{\ell}}{dt} \le 4J\alpha_{\ell-1}.$$
(E9)

The factor of 4 is a crude upper bound that takes into account both  $h_{r_0+\ell}$  and  $h_{r_0-\ell}$ , which both appear twice due to the commutator.

Now we solve the upper limit of this system of differential equations with the initial condition that  $\alpha_0 = ||W||$  and  $\alpha_{\ell>0}(t=0) = 0$ . The result is

$$\alpha_{\ell}(t) \le \|W\|_{\infty} \frac{(4Jt)^{\ell}}{\ell!}.$$
(E10)

This result is almost the end of the calculation. The remaining thing to do is to estimate the difference between  $W_{\ell}$  and the true W(t). This is

$$\|W(t) - W_{\ell}\|_{\infty} \le \sum_{j=\ell+1}^{\infty} \alpha_j \le \sum_{j=\ell+1}^{\infty} \|W\|_{\infty} \frac{(4Jt)^j}{j!}.$$
 (E11)

There are various ways to treat this infinite sum. For  $\ell \gg 4Jt$ , the simplest estimate is to say that it cannot by much

larger than its first term, which is quite small. More precisely, the ratio of term *j* to term j + 1 is  $4Jt/(j + 1) \le 4Jt/(\ell + 2)$ , so making even a crude approximation using a geometric series in  $4Jt/(\ell + 2)$  converges to something order one times the first term. After using Stirling's approximation for large  $\ell$ , the first term is

$$\|W\|_{\infty} \frac{(4Jt)^{\ell+1}}{(\ell+1)!} \approx \|W\|_{\infty} \left(\frac{4eJt}{\ell+1}\right)^{\ell+1}.$$
 (E12)

This result corresponds to roughly the  $\ell$ th order in perturbation theory when expanding W(t) in a Taylor series. Physically, it will describe the commutator dynamics for sufficiently small *t* and large  $\ell$ .

Neglecting the difference between  $\ell$  and  $\ell + 1$ , the first term is order one when  $\ell = 4eJt$ . Setting  $\ell_0 = 4eJt$ , the first term can be written as

$$e^{\ell \log \left(\ell_0/\ell\right)}.\tag{E13}$$

Using  $-1 \ge -1 + \log (\ell_0/\tilde{\ell})$  (valid for  $\tilde{\ell} \ge \ell_0$ ) and integrating both sides from  $\ell_0$  to  $\ell$ , it follows that

$$-(\ell - \ell_0) \ge \ell \log \frac{\ell_0}{\ell}.$$
 (E14)

The left-hand side is the first-order expansion of the righthand side in  $\ell - \ell_0$ , so the inequality states that going beyond first order only decreases the value. Hence,

$$e^{\ell \log(\ell_0/\ell)} < e^{-(\ell-\ell_0)},$$
 (E15)

or, using  $\ell_0 = 4eJt$ ,

$$\|W(t) - W_{\ell}\|_{\infty} \le \|W\|_{\infty} f(t) e^{4eJt - \ell}.$$
 (E16)

Here f(t) is a polynomial prefactor that does not affect the basic exponential scaling. Note that the bound is definitely not tight at very large  $\ell$ , since  $1/\ell!$  decreases faster than  $e^{-\ell}$ . The bound is also trivial once  $\ell < \ell_0$  because the right-hand side is growing exponentially, while the left-hand side is bounded by 2||W||.

Having established that  $W_{\ell}$  is close to W(t) for  $\ell \gg Jt$ , it remains to upper bound the commutator. The idea is straightforward: if an operator V is a distance r from W then an upper bound on the commutator ||[W(t), V]|| is obtained by approximating W with  $W_{\ell=r-1}$  since  $W_{r-1}$  exactly commutes with V. First add and subtract  $W_{r-1}$  inside the norm to give

$$\|[W(t), V]\|_{\infty} = \|[W(t) - W_{r-1} + W_{r-1}, V]\|_{\infty}, \quad (E17)$$

and then use  $[W_{r-1}, V] = 0$  and the bound on  $||W(t) - W_{r-1}||$  to obtain

$$\|[W(t), V]\|_{\infty} \le 2 \|V\|_{\infty} \|W(t) - W_{r-1}\|_{\infty}.$$
 (E18)

Using the upper bound above, this is

$$\|[W(t), V]\|_{\infty} \le 2 \|V\|_{\infty} \|W\|_{\infty} f(t) e^{4eJt-r}, \qquad (E19)$$

which is Eq. (15) in the main text.

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